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MACHINE CODE GENERATION FOR SIMNON ON VAX-11

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Abstract The compiler in the interactive simulation program SIMNON generates pseudocode of the equations to be simulated. The pseudocode is interpreted at simulation time. This report describes a new stage in the compiler that generates machine code directly for the host computer from the pseudocode. The code generation for VAX-11 is described in detail.			
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1. INTRODUCTION

SIMNON is a FORTRAN-written interactive simulation program for non-linear systems described by ordinary differential equations and difference equations (Elmqvist, 1975).

SIMNON contains a compiler for the simulation language that produces a pseudocode which is interpreted and executed by a FORTRAN subroutine (CALCUL) at simulation time. This makes the execution of the code slow compared to code from a compiler that produces machine code. It is possible to add a new stage to SIMNON that takes the pseudocode and generates machine code for a specific computer. The code is then executed directly by the machine rather than interpreted by a program. The cost of generating machine code from the pseudocode is small compared to the cost of producing the pseudocode from the source files. (The term cost in this case means resources such as memory and CPU-time.) The gain in execution speed has been 5 - 8 times for the existing implementations.

Chapter 2 describes the output from the compiler to be used as input for the code generator. This information is mostly taken from (Elmqvist, 1978).

Chapter 3 describes the code generation in general and chapter 4 describes the code generated for VAX-11.

2. OUTPUT FROM SIMNON

This chapter describes the format of the output from the SIMNON compiler which will be the input to the code generator.

The pseudocode is stored in an integer array in common /PSCODE/. It is organized as linked lists. A node contains one or more equations or a call of a section in an external FORTRAN system (via SYSTS). The pseudocode area contains five different lists for different kinds of computations in the simulation part such as: initial computations, derivative computations or computation of discrete states. For the code generator it is irrelevant when each list is used. The important thing is where each list starts and this is stored in commonblock /ENTRYS/. Each variable in /ENTRYS/ points to the head of a list in /PSCODE/. See fig. 1 for description of a node head.

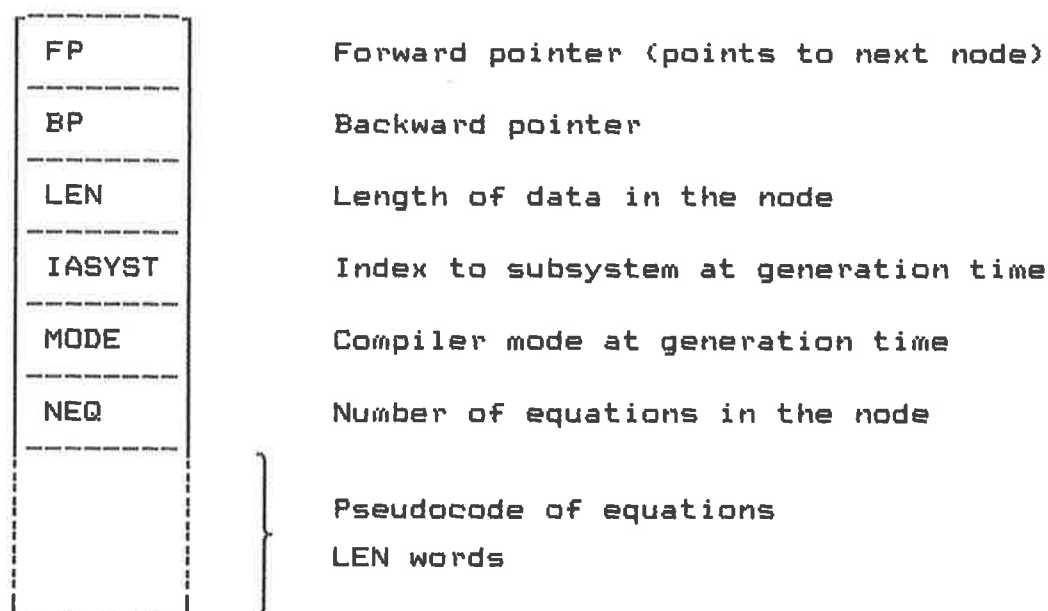


Fig. 1: Node organization

Each list in the pseudocode consists of a list head and zero or more nodes. A list head is an empty node (LEN=0).

The pseudocode consists of operators (integer 1 - 22) followed by zero or more integer operands. See fig. 2.

Output from SIMNON

OPERATION	MNEMONIC	CODE
Logical or	OR	1
Logical and	AND	2
Logical not	NOT	3
Test less than	TLT	4
Test greater than	TGT	5
Add	ADD	6
Subtract	SUB	7
Multiply	MUL	8
Divide	DIV	9
Negate	NEG	10
Raise	RAI	11
Jump if false	JMPF	12
Jump	JMP	13
Stack	FETCH	14
		pointer
Unstack	DEPOS	15
		pointer
Apply function	FUNC	16
		function number
Call FORTRAN system	CALL	20
		isyst
		ipart
Skip if not sampling	SCOND	21
		system number
No operation	NOP	22

Fig. 2: Operation set

The SIMNON compiler produces RPN (Reverse Polish Notation) code operating on a stack.

A detailed description of the operations follows.

P(n) is top stack element.

n is the stack pointer

k is the index in the pseudocode (=PC, program counter)
 Logical values are 0.0 (false) and 1.0 (true). A value is true if it is greater or equal 0.5

The pointer used in FETCH and DEPOS has the following meaning: if pnt > 10000 it points to a literal stored in common /VALUES/V(pnt-10000) otherwise it points to a variable whose address is stored in common /VARTB2/IPNTS(pnt).

OR $P(n-1) := (P(n-1) \geq 0.5) \text{ or } (P(n) \geq 0.5); n := n-1;$
 $k := k+1$

AND $P(n-1) := (P(n-1) \geq 0.5) \text{ and } (P(n) \geq 0.5); n := n-1;$
 $k := k+1$

Output from SIMNON

```

NOT      P(n):= not (P(n)≥0.5); k:=k+1

TLT      P(n-1):=P(n-1)<P(n); n:=n-1; k:=k+1

TGT      P(n-1):=P(n-1)>P(n); n:=n-1; k:=k+1

ADD      P(n-1):=P(n-1)+P(n); n:=n-1; k:=k+1

SUB      P(n-1):=P(n-1)-P(n); n:=n-1; k:=k+1

MUL      P(n-1):=P(n-1)*P(n); n:=n-1; k:=k+1

DIV      P(n-1):=P(n-1)/P(n); n:=n-1; k:=k+1

NEG      P(n):=-P(n); k:=k+1

RAI      P(n-1):=P(n-1)**P(n); n:=n-1; k:=k+1

JMPF
nr      k:=if P(n)<0.5 then k+nr+1 else k+2; n:=n-1
      (nr > 0)

JMP
nr      k:=k+nr+1 (nr > 0)

FETCH
pnt      n:=n+1; P(n):=var      ; k:=k+2
      pnt

DEPOS
pnt      var      :=P(n); n:=n-1; k:=k+2
      pnt

FUNC
nr      one-argument function:
      P(n):=func      (P(n)); k:=k+2
      nr

      two-argument function:
      P(n-1):=func      (P(n-1),P(n)); n:=n-1; k:=k+2
      nr

CALL
isyst
ipart    ISYST:=isyst; IPART:=ipart; call SYSTS;
      if ISTOP then EXIT from CALCUL; k:=k+3

SCOND
nr      if LCOND(nr) then k:= next node else
      k:=k+2

NOP      k:=k+1

```


Output from SIMNON

The following functions are used in FUNC:

NR	NAME	DESCRIPTION
1	SQRT(X)	square root of X, $X \geq 0$
2	EXP(X)	exponential function of X
3	LN(X)	natural logarithm of X, $X > 0$
4	LOG(X)	logarithm (base 10) of X, $X > 0$
5	SIN(X)	sine of X (X in radians)
6	COS(X)	cosine of X (X in radians)
7	TAN(X)	tangent of X (X in radians)
8	ATAN(X)	arctangent of X, result in radians in interval $[-\pi/2, \pi/2]$
9	ABS(X)	absolute value of X
10	SIGN(X)	the sign of X: +1.0 if $x > 0$ 0.0 if $x = 0$ -1.0 if $x < 0$
11	INT(X)	integer part of x
12	ATAN2(X,Y)	arctangent of x/y , result in radians in interval $[-\pi, \pi]$
13	MOD(X,Y)	X modulo Y $(X - \text{INT}(X/Y) * Y)$
14	MIN(X,Y)	minimum of X and Y
15	MAX(X,Y)	maximum of X and Y
16	ARCSIN(X)	arcsine of X in interval $[-1, 1]$
17	ARCCOS(X)	arccosine of X in interval $[-1, 1]$
18	SINH(X)	hyperbolic sine of X
19	COSH(X)	hyperbolic cosine of X
20	TANH(X)	hyperbolic tangent of X

A few examples of SIMNON pseudocode follows:

```

Y1 = (-A)/(3 + B)
FETCH      A
NEG
FETCH      3.00
FETCH      B
ADD
DIV
DEPOS      Y1

```

```

Y2 = A OR NOT B
FETCH      A
FETCH      B
NOT
OR
DEPOS      Y2

```

```

Y3 = (A<B) + 2
FETCH      A
FETCH      B
TLT
FETCH      2.00

```

Output from SIMNON

```
ADD
DEPOS      Y3
```

```
Y4 = IF A THEN 1 ELSE B↑2
FETCH      A
JMPF       5
FETCH      1.00
JMP        6
FETCH      B
FETCH      2.00
RAI
DEPOS      Y4
```

```
Y5 = SIN(ATAN2(A,B + 1))
FETCH      A
FETCH      B
FETCH      1.00
ADD
FUNC        ATAN2
FUNC        SIN
DEPOS      Y5
```

```
Y6 = IF A<B THEN (IF NOT B THEN 1 ELSE A + B) ELSE
      IF A + B THEN 2 ELSE B
FETCH      A
FETCH      B
TLT
JMPF       17
FETCH      B
NOT
JMPF       5
FETCH      1.00
JMP        6
FETCH      A
FETCH      B
ADD
JMP        14
FETCH      A
FETCH      B
ADD
JMPF       5
FETCH      2.00
JMP        3
FETCH      B
DEPOS      Y6
```

```
Y7 = -(A*(-B)/C - A/B + B*C - 2)
FETCH      A
FETCH      B
NEG
FETCH      C
DIV
MUL
FETCH      A
FETCH      B
```

Output from SIMNON

DIV	
SUB	
FETCH	B
FETCH	C
MUL	
ADD	
FETCH	2.00
SUB	
NEG	
DEPOS	Y7

3. CODE GENERATOR - GENERAL PRINCIPLES

The code generator scans the pseudocode in one pass and generates machine instructions directly for each pseudocode instruction. Since the pseudocode contains forward jumps it is necessary to go back and insert the jump address in the generated code when the target of the jump is processed.

The stack used in the RPN pseudocode corresponds to an operand address stack used by the code generator but there is no explicit stack or stack instructions in the generated code. The code generator makes use of the available general purpose fast hardware registers for storing intermediate results of computations and the allocation is made from a stack of free registers. Since the number of available registers generally is smaller than a worst-case expression in SIMNON the code generator automatically allocates temporary memory cells when the register stack is exhausted. The code generator will in fact work even when only one register is available.

The operand stack and register allocation

Consider the equation $Y=(A+2)*B$. It generates the following pseudocode:

```

1  FETCH  A
3  FETCH  2
5  ADD
6  FETCH  B
8  MUL
9  DEPOS  Y

```

The code generator performs the following actions:

```

1:  Push addr(A) on operand stack
3:  Push addr(2) on operand stack
5:  Pop the stack twice; if operands are registers, return
    the registers to the register stack; get a new register
    from the register stack; issue an instruction to add the
    two operands and store result in the register; push the
    register (or rather the address of the register) on the
    operand stack
6:  Push addr(B)
8:  Pop the stack twice; release registers; get a register;
    issue multiply instruction; push register
9:  Pop the stack; release register; issue instruction to
    store operand in Y

```

The following instructions could be generated:

```

ADD    A,2    →R1
MULT   R1,B   →R1
STORE  R1     →Y

```

If the code generator is to be used for a computer that only

Code generator - general principles

allows one of the operands to be a memory cell the code would be:

```

LOAD  A      →R1
ADD   R1,2    ; result is in register R1
MULT  R1,B
STORE R1      →Y

```

The logic would be: at step 5, 8 and 9 check that at least one of the popped operands is in a register; if not issue a load instruction

The operand address stack thus will contain a mixture of (addresses of) variables and registers. The registers represent results from already issued instructions. Note the difference between this stack and the stack used by the interpreter for the pseudocode that contains only values of variables or intermediate computations.

The register stack

Addresses of the free registers are stored in a stack that doesn't exist! The only variable actually associated with the stack is the stackpointer. The reason for this is that the registers are considered to be an ordered set and allocation/deallocation of the registers always made in that order. (Of course deallocation is made in the opposite order of allocation.) This makes it possible to compute the address of the register on top of the "stack" directly from the value of the stackpointer. This principle of allocating registers is important when generating code for conditional branches as shown later on.

Jumps

As mentioned earlier it is necessary to go back in the generated code and "patch" whenever a jump is encountered. The SIMNON compiler produces three branch or jump instructions, all of which are forward jumps. JMPF and JMP are results of IF-THEN-ELSE expressions and the displacement is given as a relative pseudocode address. SCOND is used for execution of equations only at specified conditions (i.e. time for sampling). The displacement is always to the end of the current node. The corresponding pseudocode address is easy to find using the information in the node head. When a jump instruction is encountered in the pseudocode the code generator issues a branch instruction with a void displacement and stores the following information in a special jump address table:

- a) the target pseudocode address
- b) the current abs. code address
- c) the type of jump instruction

For each pseudo instruction processed a check is made if the current pseudocode address is in column a) in the jump table. If it is found the current abs. code address is inserted in

Code generator - general principles

the abs. code at the address indicated by b) and the entry is removed from the table.

Conditional branches

The equation: $Y = \text{IF cond THEN expr1 ELSE expr2}$ generates the following pseudocode sequence:

```

cond-code
JMPF    L1
-----
expr1-code    block A
-----
JMP        L2
-----
L1: expr2-code    block B
-----
L2: DEPOS    Y

```

The interpreter evaluates the condition and then one of the blocks A or B. The code generator proceeds through both A and B and must generate correct code for both cases. Since both blocks in an IF-THEN-ELSE construction consists of expressions the result at the end of the block will be on top of the operand stack either in a register or in a variable (only if the expression is a simple variable). Furthermore the operand stack has increased with exactly one element from the beginning of the block to the end.

A necessary condition for correct code after the two branches is that at the end of both blocks the result is at the same address, i.e. the top element on the operand stack must be exactly the same. The code generator solves this the following way:

At the end of each block it checks if the top stack element is a variable, in which case a register is popped from the register stack and an instruction to load the variable into the register issued. The register operand is then pushed on the operand stack.

Furthermore it is necessary to restore both the operand stack and the register stack to the status of the beginning of the block at the end of block A to ensure that the status is the same at the end of both blocks. This is easily done by popping the operand stack and releasing the register.

The end of block A is always followed by a JMP instruction and the end of block B is followed by the target of a JMP instruction.

Code generator - general principles

Temporary variables

The temporary variables are allocated linearly from an array and are deallocated only at the end of an equation (DEPOS). There are two cases when temporary variables are needed. The first case is when a mathematical procedure is called and at least one argument is in a register. This is because of the parameter transfer mechanism that needs the address of an argument. Since registers are (usually) addressed differently from memory cells it would cause difficulties if an argument address was a register. Furthermore the mathematical procedure might use the register internally of course after saving the contents in a memory cell but the argument address would still point to the register.

The other case is when the register stack is exhausted. In this case the strategy is as follows: The operand stack is scanned from the top and when a register is found it is released and a load instruction to a temporary variable issued. The address of that variable will then replace the register address in the operand stack. The released register is of course always the last register in the register stack. This strategy will take care of any problem that could arise when generating code for conditional branches.

Interface SIMNON -- absolute code

To execute code in one of the five lists SIMNON takes one of the values from /ENTRY5/ and stores it in /ENTRY/ and then calls CALCUL (without arguments). The code generator replaces the pseudocode entryptoints in /ENTRY5/ with absolute code start addresses and thus the value in /ENTRY/ is the absolute start address in the generated code when CALCUL is called. CALCUL must be written in assembler but can be very simple: it makes a subroutine jump to the address in /ENTRY/ (there must be a possibility to return back to CALCUL after the code is executed).

Run-time tracing

The code generator generates one additional instruction for each new node that simply stores the node pointer for the pseudocode node in common /CSIERR/. This gives a possibility to give run-time error messages and print the source equation if the pseudocode is saved since there exists subroutines in SIMNON that will reconstruct and print out a source equation if a pointer to the pseudocode node is given. The kind of errors where this is used are arithmetic faults (floating-point overflow, divide with zero) or illegal arguments to a mathematical procedure, e.g. SQRT(-1). These errors will usually cause an interrupt or trap and most systems gives the user the possibility to write a routine that is automatically executed when such a condition occurs.

Code generator - general principles

Debugging

It is very valuable to be able to print out information concerning the code generation both for debugging purposes and checking of the generated code. The code generator will print information on four levels:

- a) the source equations
- b) the input pseudocode
- c) the generated code in symbolic assembler format
- d) the generated machine code in hexadecimal format

The printout is governed by flags in common /VXCLOG/ and the values are taken from global variables that can be changed by the LET command in SIMNON.

- a) is turned on if LOGSRC.VCODE is non-zero
- b) is turned on if LOGPSE.VCODE is non-zero
- c) is turned on if LOGINS.VCODE is non-zero
- d) is turned on if LOGHEX.VCODE is non-zero

The logical unit number for the output can be given in LUNLOG.VCODE. The information is printed incrementally, i.e. each time a new instruction is generated or next pseudoinstruction read.

4. CODE GENERATION FOR VAX-11

The architecture of the VAX includes variable-length instructions with a large number of addressing modes. Instructions can have up to three different operands which can either be registers or memory cells. There are 16 general purpose registers numbered R0 - R15 and R2 - R9 are used for the register stack in the code generator. The constants 0.5 and 1.0 are always in R10 and R11 since they are frequently used for logical testing and setting. All addresses refer to bytes. Detailed descriptions of the instructions and addressing modes are given in (VAX Architecture Handbook, 1979). The interface to the mathematical procedures is described in (VAX Run-Time Library, 1980)

The code generated for each pseudocode operation is described shortly. The VAX code is represented by symbolic assembler mnemonics somewhat simplified. The notation for the different addressing modes is not used. 0.5 and 1.0 are used as literals instead of the registers R10 and R11. Rx means a register from the register stack (R2 - R9).

Logical operators: OR, AND, NOT, ILT, IGT

The logical operators will give as result the values 0.0 (false) or 1.0 (true). A common case is when the value is used for testing in an IF-construction. The code generator tests this by looking at the next pseudo instruction to see if it is JMPF. The generated code will then be different.

```

A_OR_B
    CMPF    A,0.5
    BGEQ    L1
    CMPF    B,0.5
    BGEQ    L1
    CLRF    Rx
    BRB     L2
L1:  MOVF    1.0,Rx
L2:

```

```

IF_A_OR_B_THEN_...
    CMPF    A,0.5
    BGEQ    L1
    CMPF    B,0.5
    BGEQ    L1
    BRW     J1
L1:

```

```

A_AND_B
    CMPF    A,0.5
    BLSS    L1
    CMPF    B,0.5

```

Code generation for VAX-11

```

        BLSS    L1
        MOVF    1.0,Rx
        BRB     L2
L1:     CLRF     Rx
L2:

```

```

IF_A_AND_B_THEN...
        CMPF    A,0.5
        BLSS    L1
        CMPF    B,0.5
        BLSS    L1
        BRB     L2
L1:     BRW     J1
L2:

```

```

NOT_A
        CMPF    A,0.5
        BGEQ    L1
        MOVF    1.0,Rx
        BRB     L2
L1:     CLRF     Rx
L2:

```

```

IF_NOT_A_THEN...
        CMPF    A,0.5
        BLSS    L1
        BRW     J1
L1:

```

```

A<B
        CMPF    A,B
        BLSS    L1
        CLRF     Rx
        BRB     L2
L1:     MOVF    1.0,Rx
L2:

```

```

IF_A<B_THEN...
        CMPF    A,B
        BLSS    L1
        BRW     J1
L1:

```

The code for TGT () is the same as for TLT except that BLSS is changed to BGTR.

Arithmetic binary operations: ADD, MUL, SUB, DIV

Since ADD and MUL are symmetric (commutative) they are grouped together and the only difference in the code is the instruction code: ADDFx or MULFx. If at least one of the operands are in a register the 2-address form of the instruction is used, otherwise the 3-address form.

Code generation for VAX-11

A_+_B
ADDF3 B,A,Rx

A_+_Rx
ADDF2 A,Rx

Rx_+_B
ADDF2 B,Rx

SUB and DIV are also grouped together but here it's only if the first operand is a register that the 2-address form can be used.

A_-_B
SUBF3 B,A,Rx

Rx_-_A
SUBF2 A,Rx

A_-_Rx
SUBF3 Rx,A,Rx

NEG_-_Unitary_minus

=A
MNEGF A,Rx

RAI_-_Raise_a_number

The mathematical procedure from the Runtime Library for raising a real base to a real power is used. Note that the values of the arguments instead of the addresses are pushed on the stack in this case.

A_↑_B
PUSHL B
PUSHL A
CALLS 2,MTH\$POWRR
MOVF RO,Rx

Jump_instructions: JMPF, JMP

Since the displacement might exceed 256 bytes a BRW instruction is used instead of the shorter BRB. JMPF is the conditional branch of an IF-THEN-ELSE expression. If the condition is the result of a logical operation it is already taken care of (see above) otherwise the following code is generated:

IF_A_THEN_...
CMPF A,0.5
BGEQ L1

Code generation for VAX-11

```

        BRW    J1
L1:

```

Since JMP marks the end of the first branch block in an IF-THEN-ELSE an instruction to load the operand on top of the stack in a register is issued if it's in a variable.

FETCH_and_DEPOS

FETCH will not generate any code, DEPOS will generate an instruction to store the operand in a memory cell.

```

A=B
        MOVF    A,B

```

FUNC---library_function_call

Some of the functions are computed directly in inline code (ABS, SIGN, MIN, MAX) and the rest are calls to Runtime Library routines. In this case the argument addresses are pushed on the stack. If an argument is in a register it is first loaded in a temporary variable.

```

SIN(A)           one-argument function call
        PUSHAL  A
        CALLS   1,MTH$SIN
        MOVF    R0,Rx

```

```

ATAN2(Rx,A)      two-argument function call
        PUSHAL  A
        MOVF    Rx,TEMP$1
        PUSHAL  TEMP$1
        CALLS   2,MTH$ATAN2
        MOVF    R0,Rx

```

Inline coded functions:

```

MAX(A,B)         (BLSS is changed to BGTR in MIN)
        CMPF    B,A
        BLSS    L1
        MOVF    B,Rx
        BRB     L2
L1:      MOVF    A,Rx
L2:

```

```

SIGN(A)
        TSTF    A
        BGTR    L1
        BLSS    L2
        CLRf    Rx
        BRB     L3
L1:      MOVF    1.0,Rx
        BRB     L3

```

Code generation for VAX-11

```
L2:  MNEGF  1.0,Rx
L3:
```

```
ABS(A)
      BICL3  00008000,A,Rx
```

CALL -- calling external systems

Note that the values to be stored in ISYST and IPART are integers known at compile time.

```
      MOVZWL isyst,ISYST
      MOVZWL ipart,IPART
      CALLS  0,SYSTS
      BBC    0,ISTOP,L1
      RSB
```

```
L1:
```

SCOND -- conditional skip of a node

The pseudocode interpreter makes a call to a FORTRAN-written logical function called LCOND to evaluate the condition for SCOND. The code generator generates inline code instead.

```
      BBC    0,LSAMP,L1
      BBS    0,LSAMPS(i),L1
      BRW    0
```

```
L1:
```

Code examples:

The examples listed earlier will generate the following code. (The instruction to store the node number for each new node is omitted.)

```
; Y1 = (-A)/(3 + B)
      MNEGF  A,R2
      ADDF3  B,3.0,R3
      DIVF2  R3,R2
      MOVF   R2,Y1
```

```
; Y2 = A OR NOT B
      CMPF   B,0.5
      BGEQ   L1
      MOVF   1.0,R2
      BRB    L1
      CLRF   R2
L1:  CMPF   R2,0.5
      BGEQ   L2
      CMPF   A,0.5
      BGEQ   L2
      CLRF   R2
```

Code generation for VAX-11

```

        BRB      L3
L2:     MOVF      1.0,R2
L3:     MOVF      R2,Y2

; Y3 = (A < B) + 2
        CMPF      A,B
        BLSS      L1
        CLRF      R2
        BRB      L2
L1:     MOVF      1.0,R2
L2:     ADDF2     2.0,R2
        MOVF      R2,Y3

; Y4 = IF A THEN 1 ELSE B*2
        CMPF      A,0.5
        BGEQ      L1
        BRW      J1
L1:     MOVF      1.0,R2
        BRW      J2
J1:     PUSHL     2.0
        PUSHL     B
        CALLS     2,MTH$POWRR
        MOVF      R0,R2
J2:     MOVF      R2,Y4

; Y5 = SIN(ATAN2(A , B + 1))
        ADDF3     1.0,B,R2
        MOVF      R2,TEMP#1
        PUSHAL    TEMP#1
        PUSHAL    A
        CALLS     2,MTH$ATAN2
        MOVF      R0,R2
        MOVF      R2,TEMP#2
        PUSHAL    TEMP#2
        CALLS     1,MTH$SIN
        MOVF      R0,R2
        MOVF      R2,Y5

; Y6 = IF A < B THEN (IF NOT B THEN 1 ELSE A + B) ELSE
      IF A + B THEN 2 ELSE B
        CMPF      A,B
        BLSS      L1
        BRW      J3
L1:     CMPF      B,0.5
        BLSS      L2
        BRW      J1
L2:     MOVF      1.0,R2
        BRW      J2
J1:     ADDF3     B,A,R2
J2:     BRW      J5
J3:     ADDF3     B,A,R2
        CMPF      R2,0.5
        BGEQ      L3
        BRW      J4
L3:     MOVF      2.0,R2

```

Code generation for VAX-11

```
      BRW      J5
J4:   MOVF     B,R2
J5:   MOVF     R2,Y6
```

‡ Y7 = -(A*(-B)/C - A/B + B*C - 2)

```
      MNEGF    B,R2
      MULF2    A,R2
      DIVF2    C,R2
      DIVF3    B,A,R3
      SUBF2    R3,R2
      MULF3    C,B,R3
      ADDF2    R3,R2
      SUBF2    2.0,R2
      MNEGF    R2,R2
      MOVF     R2,Y7
```

5. REFERENCES

- Elmqvist, H (1975): SIMNON - An interactive simulation program for nonlinear systems, User's manual, TFRT 7502
- Elmqvist, H (1978): SIMNON - An interactive simulation program, Implementation, TFRT 7148
- VAX-11 Architecture Handbook (1979-80)
Digital Equipment Corporation
- VAX-11 Run-Time Library Reference Manual (1980)
Digital Equipment Corporation

6. Program listings

VCODE.FOR	- Main subroutine for code generation
CALCUL.MAR	- Interface routine between Simnon and machine code
GETREG.FOR	- Register assignment subroutine
NEXTOP.FOR	- Decodes pseudo code operation
UNSTK.FOR	- Handles operand stack
VBYTE.FOR	- Writes instruction bytes in code area
VINSTX.FOR	- Machine code instruction generator

SUBROUTINE VCODE

GENERATES MACHINE CODE FOR VAX-11 FOR SIMNON ,USING RPN
PSEUDO CODE IN /PSCODE/ AS INPUT.

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SUBROUTINES REQUIRED

IADDR
NEXTOP
UNSTK
GETREG
VINSTX

DATA BASE:

/PSCODE/, /ENTRYS/, /DESTIN/, /USER/ , /VARTB2/ , /VALUES/ ,
/SIMN/, /MESSS/, /DEVICE/, /LIMITS/ AND /COND/ : SEE SIMNON

/UPNTS/ CONTAINS VARIOUS POINTERS
L- CURRENT PSEUDO CODE INDEX
IHEAD- POINTER TO CURRENT LIST IN /PSCODE/
NODE- POINTER TO CURRENT NODE IN THE LIST
LEN- REMAINING LENGTH OF CURRENT NODE
IPCNT- REL. ADDR. FOR NEXT INSTRUCTION
IRGTOP- POINTER TO NEXT FREE REGISTER IN THE (NON-
EXISTENT) REGISTER STACK. THE ADDRESS OF THE
REGISTER IS COMPUTED FROM IRGTOP AS: -3 -IRGTOP
ITOP- STACK POINTER FOR OPERAND ADDRESS STACK

/UDATA/ CONTAINS TABLES AND STACKS
ISTACK(*)- OPERAND ADDRESS STACK
ICODE(*)- STORAGE FOR ABSOLUTE CODE

/UTEMPC/ CONTAINS TEMPORARY CELLS USED 1): IF THE REGISTER
STACK IS EMPTY AND 2): FOR STORING REGISTER
OPERANDS IN LIBRARY FUNCTION CALLS
ITMPNT - POINTER TO LAST USED TEMP. VARIABLE
MAXTMP - MAX VALUE ALLOWED OF ITMPNT
ITMPAD - ABS. ADDRESS TO ITEMP(1)
ITEMP(*) - TEMPORARY VARIABLES

/CSIERR/ INODE WILL CONTAIN THE NODE NUMBER OF THE EQUATION
CURRENTLY BEING EXECUTED IN A SIMU COMMAND. THIS GIVES
A POSSIBILITY TO PRINT ERROR MESSAGES AT RUN TIME
REFERENCING SOURCE LINE EQUATIONS.

/VXCLOG/ USED FOR DEBUGGING AND CHECKING OF CODE GENERATOR
LUNLOG - LOGICAL UNIT NUMBER FOR DEBUG OUTPUT
LOGSRC - IF NON-ZERO: PRINT SOURCE NODE EQUATION(S)
LOGPSE - IF NON-ZERO: PRINT PSEUDO-CODE INSTR.
LOGINS - IF NON-ZERO: PRINT ASSEMBLY INSTR.
LOGHEX - IF NON-ZERO: PRINT HEXADEC. MACHINE CODE

INTERNAL DATA FOR UCODE:

IFTAB(2,20) - FUNC TABLE
1,* - IF >0 :GOTO INDEX TO INTRINSIC CODE GENERATION

```

C          SEQUENCE ELSE
C          -(1,*) IS INDEX IN IFNADR(*) TO ABSOLUTE ENTRY
C          OF THE LIBRARY FUNCTION AND (2,*) IS NR. OF ARGS
C
C          IBASE- ABSOLUTE BASE ADDRESS FOR THE CODE
C          JMPPNT- POINTER TO LAST USED ENTRY IN JMP TABLE
C          JMPTAB(3,*)- FORWARD JUMP TABLE
C          1- JUMP-TO PSEUDO CODE ADDR.
C          2- JUMP-FROM REL. CODE ADDR (IPCNT ADDR)
C          3- TYPE OF JUMP:
C          1: JMPF
C          2: JMP
C          3: SCND
C
C          NOTES ON REGISTER USAGE:
C
C          THERE ARE 16 AVAILABLE REGISTERS NUMBERED FROM R0 TO R15
C          WITH THE CORRESPONDING ADDRESSES 0 TO 15. TO BE ABLE TO TELL
C          IF AN OPERAND ADDRESS IS A REGISTER ,THEIR NEGATIVE VALUES-1 ARE
C          USED IN THE PROGRAM.
C          R2 - R9 ARE USED IN THE REGISTER STACK, R10 AND R11 ARE USED
C          FOR CONSTANTS 0.5 AND 1.0
C
C          LOGICAL ISTOP,NXJPF,NOSYST,NOCODE,LSAMP,
C          1LSAMPS,EOL,LDUM1,LDUM2,LDUM3
C
C          INTEGER MOVF,CLRF,MNEGF,CMPF,ADDF2,SUBF2,MULF2,DIVF2,BGEQ,
C          *          ADDF3,SUBF3,MULF3,DIVF3,BRB,BRW,CALLS,BLSS,BBS,
C          *          BGTR,MOVZWL,BICL2,BICL3,TSTF,RSB,PUSHL,PUSHAL,BBC
C          INTEGER R0,R1,R2,RZP5,R1PZ,SIGNEG
C
C          DIMENSION IFTAB(2,20),IFNADR(16),JMPTAB(3,10),HSRC(2),
C          *          HLUN(2),HPSE(2),HINS(2),HHEX(2),HEXT(2),PSOP(23),
C          *          NPSOP(23)
C
C          COMMON/DEVICE/LKB,LTP,LLP,LDIS,LTO,LPLT,LXXX,LDK1,LDK2,LDK3,LDK4
C          COMMON/SIMN/NOSYST,LDUM1(4),NOCODE,LDUM2(4)
C          COMMON/LIMITS/MPSC,IDUM3(4)
C          COMMON/MESSS/MESS
C          COMMON/PSCODE/IPSEUD(100)
C          COMMON/VARTB2/IPNTS(100)
C          COMMON/VALUES/VALUE(100)
C          COMMON/ENTRYS/IENT(5)
C          COMMON/COND/LSAMP,LSAMPS(5)
C          COMMON/DESTIN/ISYST,IPART
C          COMMON/USER/ISTOP,LDUM3(6),IDUM2
C          COMMON/UPNTS/L,IHEAD,NODE,LEN,IPCNT,IRGTOP,ITOP
C          COMMON/UDATA/ISTACK(25),ICODE(7000)
C          COMMON/UTEMPC/ITMPNT,MAXTMP,ITMPAD,ITEMP(20)
C          COMMON/CSIERR/INODE,INDEX
C          COMMON/VXCLOG/LUNLOG,LOGSRC,LOGPSE,LOGINS,LOGHEX
C
C          EXTERNAL OTS$POWRR,MTH$SQRT,MTH$EXP,MTH$ALOG,MTH$ALOG10,
C          *          MTH$SIN,MTH$COS,MTH$TAN,MTH$ATAN,MTH$ASIN,
C          *          MTH$ACOS,MTH$SINH,MTH$COSH,MTH$TANH,
C          *          MTH$ATAN2,MTH$AMOD,SYSTS
C
C          DATA
C          *          MOVF/' 00000050' X/
C          *          ,MNEGF/' 00000052' X/
C          *          ,CLRF/' 000000D4' X/

```

```

*      ,CMPF/'00000051'X/
*      ,ADDF2/'00000040'X/
*      ,SUBF2/'00000042'X/
*      ,MULF2/'00000044'X/
*      ,DIVF2/'00000046'X/
*      ,ADDF3/'00000041'X/
*      ,SUBF3/'00000043'X/
*      ,MULF3/'00000045'X/
*      ,DIVF3/'00000047'X/
*      ,BGEQ/'00000018'X/
*      ,BLSS/'00000019'X/
*      ,BGTR/'00000014'X/
*      ,BRB/'00000011'X/
*      ,BRW/'00000031'X/
*      ,CALLS/'000000FB'X/
*      ,MOVZWL/'0000003C'X/
*      ,BICL2/'000000CA'X/
*      ,BICL3/'000000CB'X/
*      ,TSTF/'00000053'X/
*      ,RSB/'00000005'X/
*      ,PUSHL/'000000DD'X/
*      ,PUSHAL/'000000DF'X/
*      ,BBC/'000000E1'X/
*      ,BBS/'000000E0'X/

```

C

```

DATA IFTAB/
*      -1,1,
*      -2,1,
*      -3,1,
*      -4,1,
*      -5,1,
*      -6,1,
*      -7,1,
*      -8,1,
*      1,0,
*      2,0,
*      -9,1,
*      -10,2,
*      -11,2,
*      3,0,
*      4,0 ,
*      -12,1,
*      -13,1,
*      -14,1,
*      -15,1,
*      -16,1 /

```

C

```

DATA
*      RZP5/-11/
*      ,R1PZ/-12/
*      ,R0/-1/
*      ,R1/-2/
*      ,R2/-3/
*      ,SIGNED/'00008000'X/
DATA PSOP/4HOR ,4HAND ,4HNOT ,4HTLT ,4HTGT ,4HADD ,4HSUB ,
*      4HMUL ,4HDIV ,4HNEG ,4HRAI ,4HJMPF,4HJMP ,4HFETC,
*      4HDEPO,4HFUNC,4HJMS ,4HRET ,4HSTOP,4HCALL,4HSCON,
*      4HNOP ,4HNODE/
DATA NPSOP/11*0,6*1,0,0,2,1,0,0/
C
DATA MAXJMP/10/

```

```
DATA MXTEMP/76/  
DATA MAXCOD/27000/
```

C

```
DATA HLUN/4HLUNL,4HOG /  
* ,HSRC/4HLOGS,4HRC /  
* ,HPSE/4HLOGP,4HSE /  
* ,HINS/4HLOGI,4HNS /  
* ,HHEX/4HLOGH,4HEX /  
* ,HEXT/4HVCOD,4HE /
```

C

```
IF(NOSYST .OR. MESS.NE.0 .OR. NOCODE) RETURN
```

C

C

```
CHECK DEBUG OPTIONS  
CALL FINT(HLUN,HEXT,I1,IND)  
IF(IND .LE. 0) THEN  
    LUNLOG=I1  
ELSE  
    LUNLOG=LTO  
ENDIF
```

C

```
CALL FINT(HSRC,HEXT,I1,IND)  
IF(IND .LE. 0) THEN  
    LOGSRC=I1  
ELSE  
    LOGSRC=0  
ENDIF
```

C

```
CALL FINT(HPSE,HEXT,I1,IND)  
IF(IND .LE. 0) THEN  
    LOGPSE=I1  
ELSE  
    LOGPSE=0  
ENDIF
```

C

```
CALL FINT(HINS,HEXT,I1,IND)  
IF(IND .LE. 0) THEN  
    LOGINS=I1  
ELSE  
    LOGINS=0  
ENDIF
```

C

```
CALL FINT(HHEX,HEXT,I1,IND)  
IF(IND .LE. 0) THEN  
    LOGHEX=I1  
ELSE  
    LOGHEX=0  
ENDIF
```

C

C

```
COMPUTE VARIOUS ABS. ADDRESSES  
IBASE=IADDR(ICODE(1))-1  
IF(LOGHEX .NE. 0) WRITE(LUNLOG,4010)IBASE  
4010 FORMAT(' ***** BASE ADDRESS: ',Z8.8,' *****')  
INODAD=IADDR(INODE)  
ITMPAD=IADDR(ITEMP)  
IISYST=IADDR(ISYST)  
IIPART=IADDR(IPART)  
IISTOP=IADDR(ISTOP)  
ISIGNG=IADDR(SIGNEG)  
ILSAMP=IADDR(LSAMP)  
ILSMPS=IADDR(LSAMPS(1))-4  
ISYSTS=IADDR(SYSTS)
```

```

IPWRR=IADDR(OTS$PWRR)
IFNADR(1)=IADDR(MTH$SQRT)
IFNADR(2)=IADDR(MTH$EXP)
IFNADR(3)=IADDR(MTH$ALOG)
IFNADR(4)=IADDR(MTH$ALOG10)
IFNADR(5)=IADDR(MTH$SIN)
IFNADR(6)=IADDR(MTH$COS)
IFNADR(7)=IADDR(MTH$TAN)
IFNADR(8)=IADDR(MTH$ATAN)
IFNADR(9)=IADDR(MTH$AINT)
IFNADR(10)=IADDR(MTH$ATAN2)
IFNADR(11)=IADDR(MTH$AMOD)
IFNADR(12)=IADDR(MTH$ASIN)
IFNADR(13)=IADDR(MTH$ACOS)
IFNADR(14)=IADDR(MTH$SINH)
IFNADR(15)=IADDR(MTH$COSH)
IFNADR(16)=IADDR(MTH$TANH)

```

```

C
C      INITIALIZE POINTERS & FLAGS
C

```

```

IPCNT=1
ITOP=0
IRGTOP=0
JMPPNT=0
ITMPNT=0
MAXTMP=MXTEMP
NXJPF=.FALSE.

```

```

C
C      SCAN THE FIVE PSEUDO-CODE LISTS
C
DO 3000 ILIST=1,5
  IHEAD=IENT(ILIST)
  NODE=IHEAD
  LEN=0

```

```

C
C      REPLACE PSEUDO CODE ENTRY WITH ABS. CODE ENTRY ADDRESS
C
IENT(ILIST)=IPCNT+IBASE
IF(LOGPSE .NE. 0) WRITE(LUNLOG,4030)ILIST
IF(LOGHEX .NE. 0) WRITE(LUNLOG,4040)IENT(ILIST)
4030 FORMAT(1X/' * START OF LIST:',I2,' *')
4040 FORMAT(' ***** STARTADDRESS: ',Z8.8,' *****')
EOL=.FALSE.

```

```

C
C      GET NEXT PSEUDO INSTRUCTION WORD
C

```

```

1000 IOP=NEXTOP(EOL)
IF(EOL) GO TO 2500

```

```

C
1010 IF(IOP.LT.1 .OR. IOP.GT.23) GO TO 903
C

```

```

IF(LOGPSE .NE. 0) THEN
  IF(NPSOP(IOP) .EQ. 0) THEN
    WRITE(LUNLOG,4100)L,PSOP(IOP)
  ELSE
    J1=L+1
    J2=NPSOP(IOP)+L
    WRITE(LUNLOG,4100)L,PSOP(IOP),(IPSEUD(I),I=J1,J2)
  ENDIF
ENDIF

```

```

4100 FORMAT(1X,I4,10X,4(1H*),1X,A4,1X,I5,I4)

```

```

C
IF(IPCNT.GT.MAXCOD) GO TO 920

```

```

C
C      CHECK IF THIS IS A FORWARD JUMP ADDR
C
220  CONTINUE
      IF(JMPPNT.EQ.0) GO TO 250
      DO 230 I=1,JMPPNT
      IF(JMPTAB(1,I).EQ.L) GO TO 240
230  CONTINUE
      GO TO 250

C
C      MAKE SURE THAT TOP STACK ELEMENT IS A REGISTER IF IT IS AN
C      ITYP=2 JUMP (JMP) BEFORE INSERTING FORWARD JUMP ADDRESS IN CODE
C
240  IF(JMPTAB(3,I).EQ.2 .AND. ISTACK(ITOP).GE.0) THEN
      CALL UNSTK(IAD1)
      CALL GETREG(IREG)
      CALL VINST2(MOVF,IAD1,IREG)
      CALL STACK(IREG)
      ENDIF

C
C      INSERT CURRENT ADDR IN THE JUMP INSTR. TO THIS ADDRESS
C
      CALL VJUMP(JMPTAB(2,I))

C
C      REMOVE ADDRESS FROM JMP TABLE
C
      IF(I .LT. JMPPNT) THEN
          NR=JMPPNT-I
          DO 247 J=1,NR
              IJ=I+J
              DO 247 K=1,3
247          JMPTAB(K,IJ-1)=JMPTAB(K,IJ)
          ENDIF
          JMPPNT=JMPPNT-1

C
C      CHECK IF THERE IS MORE THAN 1 JUMP TO THIS ADDRESS
C
      GO TO 220

C
C      JUMP TO CODE GENERATION SEQUENCE
C
250  GO TO(2010,2020,2030,2040,2040,2060,2070,2060,2070,2100,2110,2120,
1 2130,2140,2150,2160,2170,2180,2190,2200,2210,1000,2230),IOP

C
C      OR
C
2010 CALL UNSTK2(IAD1,IAD2)
      IOP=NEXTOP(EOL)
      NXJPF=IOP.EQ.12
      IF(.NOT.NXJPF) CALL GETREG(IREG)
      CALL VINST2(CMPF,IAD1,RZP5)
      CALL VINST1(BGEQ,0)
      IFROM1=IPCNT
      CALL VINST2(CMPF,IAD2,RZP5)
      CALL VINST1(BGEQ,0)
      IFROM2=IPCNT
      IF(.NOT.NXJPF) THEN
          CALL VINST1(CLRF,IREG)
          CALL VINST1(BRB,3)
          CALL VBRANC(IFROM1,IPCNT)
          CALL VBRANC(IFROM2,IPCNT)

```

```

        CALL VINST2(MOVF,R1PZ,IREG)
        CALL STACK(IREG)
    ELSE
        CALL VBRANC(IFROM1,IPCNT+3)
        CALL VBRANC(IFROM2,IPCNT+3)
    ENDIF
    GO TO 1010
C
C    AND
C
2020 CALL UNSTK2(IAD1,IAD2)
    IOP=NEXTOP(EOL)
    NXJPF=IOP.EQ.12
    IF(.NOT.NXJPF) CALL GETREG(IREG)
    CALL VINST2(CMPF,IAD1,RZP5)
    CALL VINST1(BLSS,0)
    IFROM1=IPCNT
    CALL VINST2(CMPF,IAD2,RZP5)
    CALL VINST1(BLSS,0)
    IFROM2=IPCNT
    IF(NXJPF) THEN
        CALL VINST1(BRB,3)
        CALL VBRANC(IFROM1,IPCNT)
        CALL VBRANC(IFROM2,IPCNT)
    ELSE
        CALL VINST2(MOVF,R1PZ,IREG)
        CALL VINST1(BRB,2)
        CALL VBRANC(IFROM1,IPCNT)
        CALL VBRANC(IFROM2,IPCNT)
        CALL VINST1(CLRF,IREG)
        CALL STACK(IREG)
    ENDIF
    GO TO 1010
C
C    NOT
C
2030 CALL UNSTK(IAD1)
    IOP=NEXTOP(EOL)
    NXJPF=IOP.EQ.12
    IF(.NOT.NXJPF) CALL GETREG(IREG)
    CALL VINST2(CMPF,IAD1,RZP5)
C
C    SPECIAL CODE IF NEXT OP. IS JMPF
C    IF(NXJPF) THEN
        CALL VINST1(BLSS,3)
    ELSE
        CALL VINST1(BGEQ,5)
        CALL VINST2(MOVF,R1PZ,IREG)
        CALL VINST1(BRB,2)
        CALL VINST1(CLRF,IREG)
        CALL STACK(IREG)
    ENDIF
    GO TO 1010
C
C    < , >
C
2040 CALL UNSTK2(IAD1,IAD2)
    IF(IOP.EQ.4) THEN
        IOPER=BLSS
    ELSE
        IOPER=BGTR

```



```

ENDIF
IOP=NEXTOP(EOL)
NXJPF=IOP .EQ. 12
IF(.NOT.NXJPF) CALL GETREG(IREG)
CALL VINST2(CMPF,IAD2,IAD1)
C
C   SPECIAL CODE IF NEXT OP. IS JMPF
IF(NXJPF) THEN
    CALL VINST1(IOPER,3)
ELSE
    CALL VINST1(IOPER,4)
    CALL VINST1(CLRF,IREG)
    CALL VINST1(BRB,3)
    CALL VINST2(MOVF,R1PZ,IREG)
    CALL STACK(IREG)
ENDIF
GO TO 1010
C
C   + , *
C
2060 CALL UNSTK2(IAD1,IAD2)
IF(IAD2 .LT. 0) THEN
    IF(IOP .EQ. 6) THEN
        IOPER=ADDF2
    ELSE
        IOPER=MULF2
    ENDIF
    CALL VINST2(IOPER,IAD1,IAD2)
    CALL STACK(IAD2)
ELSE IF(IAD1 .LT. 0) THEN
    IF(IOP .EQ. 6) THEN
        IOPER=ADDF2
    ELSE
        IOPER=MULF2
    ENDIF
    CALL VINST2(IOPER,IAD2,IAD1)
    CALL STACK(IAD1)
ELSE
    CALL GETREG(IREG)
    IF(IOP .EQ. 6) THEN
        IOPER=ADDF3
    ELSE
        IOPER=MULF3
    ENDIF
    CALL VINST3(IOPER,IAD1,IAD2,IREG)
    CALL STACK(IREG)
ENDIF
GO TO 1000
C
C   - , /
C
2070 CALL UNSTK2(IAD1,IAD2)
IF(IAD2 .LT. 0) THEN
    IF(IOP .EQ. 7) THEN
        IOPER=SUBF2
    ELSE
        IOPER=DIVF2
    ENDIF
    CALL VINST2(IOPER,IAD1,IAD2)
    CALL STACK(IAD2)
ELSE

```

```

        IF(IOP .EQ. 7) THEN
            IOPER=SUBF3
        ELSE
            IOPER=DIVF3
        ENDIF
        IF(IAD1 .LT. 0) THEN
            CALL VINST3(IOPER,IAD1,IAD2,IAD1)
            CALL STACK(IAD1)
        ELSE
            CALL GETREG(IREG)
            CALL VINST3(IOPER,IAD1,IAD2,IREG)
            CALL STACK(IREG)
        ENDIF
    ENDIF
    GO TO 1000
C
C    UNITARY -
C
2100 CALL UNSTK(IAD1)
    CALL GETREG(IREG)
    CALL VINST2(MNEGF,IAD1,IREG)
    CALL STACK(IREG)
    GO TO 1000
C
C    & (A&B = A*B)
C
2110 CALL UNSTK2(IAD1,IAD2)
    CALL VINST1(PUSHL,IAD1)
    CALL VINST1(PUSHL,IAD2)
    CALL VINST2(CALLS,2,IPOWRR)
    CALL GETREG(IREG)
    CALL VINST2(MOVF,R0,IREG)
    CALL STACK(IREG)
    GO TO 1000
C
C    JMPF
C
2120 IF(NXJPF) GO TO 2130
    CALL UNSTK(IAD1)
    CALL VINST2(CMPF,IAD1,RZP5)
    CALL VINST1(BGEQ,3)
C
C    JMP
C
2130 JMPADR=NEXTOP(EOL)
C
C    STORE FORWARD JUMP ADDR IN JMP TABLE
C
    JMPPNT=JMPPNT+1
    IF(JMPPNT.GT.MAXJMP) GO TO 906
    JMPTAB(1,JMPPNT)=JMPADR+L
    JMPTAB(2,JMPPNT)=IPCNT
    IF(IOP.EQ.12) JMPTAB(3,JMPPNT)=1
    IF(IOP.EQ.13) JMPTAB(3,JMPPNT)=2
C
C    MAKE SURE THAT TOP STACK ENTRY IS A REGISTER AND
C    THEN RELEASE IT IF OPERATION IS JMP
C
    IF(IOP .EQ. 13) THEN
        CALL UNSTK(IAD1)
        IF(IAD1 .GE. 0) THEN

```

```

                CALL GETREG(IREG)
                CALL VINST2(MOVF,IAD1,IREG)
                JMPTAB(2,JMPPNT)=IPCNT
                IRGTOP=IRGTOP-1
            ENDIF
        ENDIF
C
        CALL VINST1(BRW,0)
        NXJPF=.FALSE.
        GO TO 1000
C
C        FETCH (STACK OPERAND ADDRESS)
C
2140 IAD1=NEXTOP(EOL)
        IF(IAD1 .LT. 10000) THEN
C
C            VARIABLE OPERAND
C            IAD1=IPNTS(IAD1)
        ELSE
C
C            LITTERAL OPERAND
C            IAD1=IADDR(VALUE(IAD1-10000))
        ENDIF
        CALL STACK(IAD1)
        GO TO 1000
C
C        DEPOS (=)
C
2150 CALL UNSTK(IAD1)
        IF(ITOP.NE.0) GO TO 907
        IF(IRGTOP.NE.0) GO TO 908
        IF(IAD1.LT.0 .AND. IAD1.NE.R2) GO TO 909
        IOPAND=NEXTOP(EOL)
        IF(IOPAND.GT.10000) GO TO 910
        IOPAND=IPNTS(IOPAND)
        ITMPNT=0
C
        CALL VINST2(MOVF,IAD1,IOPAND)
        GO TO 1000
C
C        FUNC
C
2160 IFUNC=NEXTOP(EOL)
C
2162 IFENT=IFTAB(1,IFUNC)
        IF(IFENT.LE.0) GO TO 2165
C
C        INTRINSIC CODED FUNCTIONS
C        GO TO(601,602,603,603),IFENT
C
C        ABS
C
601  CALL UNSTK(IAD1)
        IF(IAD1 .LT. 0) THEN
            IREG=IAD1
            CALL VINST2(BICL2,ISIGNG,IREG)
        ELSE
            CALL GETREG(IREG)
            CALL VINST3(BICL3,ISIGNG,IAD1,IREG)
        ENDIF
        CALL STACK(IREG)

```

```

        GO TO 1000
C
C      SIGN
C
602  CALL UNSTK(IAD1)
      CALL GETREG(IREG)
      CALL VINST1(TSTF,IAD1)
      CALL VINST1(BGTR,6)
      CALL VINST1(BLSS,9)
      CALL VINST1(CLR,IREG)
      CALL VINST1(BRB,8)
      CALL VINST2(MOVF,R1PZ,IREG)
      CALL VINST1(BRB,3)
      CALL VINST2(MNEGF,R1PZ,IREG)
      CALL STACK(IREG)
      GO TO 1000
C
C      MIN , MAX
C
603  CALL UNSTK2(IAD1,IAD2)
      IF(IFENT .EQ. 4) THEN
        IOPER=BLSS
      ELSE
        IOPER=BGTR
      ENDIF
      CALL GETREG(IREG)
      CALL VINST2(CMPF,IAD1,IAD2)
      CALL VINST1(IOPER,0)
      IFROM1=IPCNT
      CALL VINST2(MOVF,IAD1,IREG)
      CALL VINST1(BRB,0)
      IFROM2=IPCNT
      CALL VBRANC(IFROM1,IPCNT)
      CALL VINST2(MOVF,IAD2,IREG)
      CALL VBRANC(IFROM2,IPCNT)
      CALL STACK(IREG)
      GO TO 1000
C
C      LIBRARY FUNCTION
C
2165 NARGS=IFTAB(2,IFUNC)
      DO 2166 I=1,NARGS
        CALL UNSTK(IAD)
        IF(IAD .LT. 0) THEN
C
C          STORE REGISTER OPERAND IN TEMP. VARIABLE
          IF(ITMPNT .GT. MAXTMP) GO TO 911
          IAD1=ITMPAD+ITMPNT
          ITMPNT=ITMPNT+4
          CALL VINST2(MOVF,IAD,IAD1)
          IAD=IAD1
        ENDIF
2166 CALL VINST1(PUSHAL,IAD)
C
      CALL VINST2(CALLS,NARGS,IFNADR(-IFENT))
      CALL GETREG(IREG)
      CALL VINST2(MOVF,R0,IREG)
      CALL STACK(IREG)
      GO TO 1000
C
C      JMS

```

```

C
2170 CONTINUE
GO TO 903
C
C    RET
C
2180 CONTINUE
GO TO 903
C
C    STOP
C
2190 CONTINUE
GO TO 903
C
C    CALL
C
2200 JSYST=NEXTOP(EOL)
JPART=NEXTOP(EOL)
CALL VINST2(MOVZWL,JSYST,IISYST)
CALL VINST2(MOVZWL,JPART,IIPART)
CALL VINST2(CALLS,0,ISYSTS)
C
C    NO NEED TO TEST FOR ISTOP IF END OF LIST
IOP=NEXTOP(EOL)
IF(EOL) GO TO 2500
CALL VINST3(BBC,0,IISTOP,1)
CALL VINSTO(RSB)
GO TO 1010
C
C    SCOND
C
2210 ICOND=NEXTOP(EOL)
IF(ICOND.EQ.0) GO TO 1000
C
C    CHECK NEXT NODE BEFORE GENERATING SCOND INSTR.
NOD=NODE
2212 NOD=IPSEUD(NOD)
IF(NOD.EQ.IHEAD) GO TO 2214
C
C    INHIBIT SCOND IN NEXT NODE IF IT IS SAME SUB-SYSTEM
IF(IPSEUD(NOD+6).NE.IOP) GO TO 2218
IF(IPSEUD(NOD+7).NE.ICOND) GO TO 2218
C
C    REMOVE THIS SCOND
IPSEUD(NOD+7)=0
GO TO 2212
C
C    END OF LIST FOUND: SPECIAL JUMP ADDR
2214 NOD=-6
C
2218 CALL VINST3(BBC,0,ILSAMP,8)
CALL VINST3(BBS,0,ILSMPS+4*ICOND,3)
JMPPNT=JMPPNT+1
IF(JMPPNT.GT.20) GO TO 906
JMPTAB(1,JMPPNT)=NOD+6
JMPTAB(2,JMPPNT)=IPCNT
JMPTAB(3,JMPPNT)=3
CALL VINST1(BRW,0)
GO TO 1000
C
C    START OF NEW NODE

```

```

C
2230 CALL VINST2(MOVZWL,NODE,INODAD)
      IF(LOGSRC.NE.0) THEN
          ISYOLD=0
          WRITE(LUNLOG,4230)
          CALL PREQND(LUNLOG,NODE,ISYOLD)
      ENDIF
4230 FORMAT(1X)
      GO TO 1000
C
C      END OF LIST
2500 IF(JMPPNT.EQ.0) GO TO 2520
      IF(JMPPNT.NE.1) GO TO 901
      IF(JMPTAB(1,1).NE.0) GO TO 901
C
C      SCND JUMP TO LAST INSTR.
      CALL VJUMP(JMPTAB(2,1))
      JMPPNT=0
C
2520 CALL VINSTO(RSB)
C
3000 CONTINUE
      RETURN
C
C      ERRORS
C
C      JMP TABLE NOT EMPTY
901  IERR=1
      GO TO 999
C
C      BAD PSEUDO INSTRUCTION
903  IERR=3
      GO TO 999
C
C      JMP TABLE OVERFLOW
906  IERR=6
      GO TO 999
C
C      STACK NOT EMPTY
907  IERR=7
      GO TO 999
C
C      REGISTER STILL ASSIGNED
908  IERR=8
      GO TO 999
C
C      MORE THAN ONE REGISTER USED
909  IERR=9
      GO TO 999
C
C      OPERAND IS A LITTERAL
910  IERR=10
      GO TO 999
C
C      TEMP TABLE OVERFLOW
911  IERR=11
      GO TO 999
C
C      INTERNAL ERROR IN VCODE
912  IERR=12
      GO TO 999

```

```
C
C      NO MORE ROOM FOR ABS CODE
920  WRITE(LTO,1920)
1920 FORMAT(' NO MORE ROOM FOR ABS CODE')
      NOSYST=.TRUE.
      RETURN
C
999  WRITE(LTO,1999)IERR
1999 FORMAT(' **VCODE** : IERR=',I3)
      STOP
      END
```

```

      .TITLE CALCUL
;
; INTERFACE BETWEEN FORTRAN AND CODE GENERATED
; BY VCODE
;
;   AUTHOR: TOMMY ESSEBO 1980-11-11
;
      .PSECT ENTRY,PIC,OVR,REL,GBL,SHR,NOEXE,RD,WRT,LONG
ENTRY:  .LONG
      .PSECT $CODE,PIC,CON,REL,LCL,SHR,EXE,RD,NOWRT,LONG
CALCUL::
      .WORD 4M<IV,R2,R3,R4,R5,R6,R7,R8,R9,R10,R11>
      PUSHAL    SIMHDL          ;
      CALLS     #1,LIB$ESTABLISH ;
      MOVL      ENTRY,R0        ;
      MOVF      #0.5,R10        ;
      MOVF      #1.0,R11        ;
      JSB       (R0)            ;
      RET                          ;
      .END

```



```

SUBROUTINE GETREG(IREG)
C
C   ASSIGNS A NEW REGISTER FROM REGISTER STACK
C
C   AUTHOR TOMMY ESSEBO 1980-11-11
C
C   SUBROUTINES REQUIRED
C       VINSTX
C
C   INTEGER MOVF
C
COMMON/UPNTS/IDUM(5),IRGTOP,ITOP
COMMON/UTEMPC/ITMPNT,MAXTMP,ITMPAD,ITEMP(20)
COMMON/UDATA/ISTACK(25),IDUM1(5000)
COMMON/DEVICE/LKB,LTP,LLP,LDIS,LTO,IDUM2(6)
C
DATA MAXREG/7/
DATA MOVF/'00000050'X/
C
IF(IRGTOP.GE.MAXREG) GO TO 20
C
C   ASSIGN REGISTER
C
IREG=-3-IRGTOP
IRGTOP=IRGTOP+1
RETURN
C
C   NO REGISTER IN REGISTER STACK
C   LOOK FOR REGISTER IN OPERAND STACK
20 IF(ITOP .EQ. 0) GO TO 91
DO 30 I=1,ITOP
J=ITOP+1-I
IF(ISTACK(J) .LT. 0) GO TO 40
30 CONTINUE
GO TO 91
C
C   MOVE OPERAND FROM REGISTER TO TEMP. VARIABLE
40 IF(ITMPNT .GT. MAXTMP) GO TO 92
IAD1=ITMPAD+ITMPNT
ITMPNT=ITMPNT+4
IREG=ISTACK(J)
CALL VINST2(MOVF,IREG,IAD1)
ISTACK(J)=IAD1
C
C   CHECK THAT THIS IS LAST REGISTER
IF(IREG .NE. -2-MAXREG) GO TO 93
RETURN
C
91 WRITE(LTO,191)
191 FORMAT(' **VCODE** NO REGISTERS IN OPERAND STACK')
STOP
C
92 WRITE(LTO,192)
192 FORMAT(' **VCODE** TEMP TABLE OVERFLOW')
STOP
C
93 WRITE(LTO,193)
193 FORMAT(' **VCODE** REGISTER ALLOCATION ERROR')
STOP
END

```

```

      INTEGER FUNCTION NEXTOP(EOL)
C
C      RETURNS NEXT PSEUDO CODE OPERATION FROM /PSCODE/
C
C      EOL - RETURNED .TRUE. IF END OF LIST
C
C      AUTHOR TOMMY ESSEBO 1980-11-11
C
C      SUBROUTINES REQUIRED
C          NONE
C
C      LOGICAL EOL
C      COMMON/PSCODE/IPSEUD(100)
C      COMMON/UPNTS/L,IHEAD,NODE,LEN,IDUM(3)
C
C      IF(LEN.GT.0) GO TO 20
C
C          NEW NODE
10      NODE=IPSEUD(NODE)
          IF(NODE.EQ.IHEAD) GO TO 50
          L=NODE+5
          LEN=IPSEUD(NODE+2)
          IF(LEN.LE.0) GO TO 10
C
C          INDICATE NEW NODE FOR CODE GENERATOR
          NEXTOP=23
          RETURN
C
20      L=L+1
          LEN=LEN-1
          NEXTOP=IPSEUD(L)
          RETURN
C
C          END OF LIST
50      EOL=.TRUE.
          RETURN
          END

```

```

SUBROUTINE UNSTK(IOPAND)
C
C UNSTACKS TOP ITEM FROM OPERAND ADDRESS STACK AND FREES IT
C IF IT IS A REGISTER
C
C AUTHOR TOMMY ESSEBO 1980-11-11
C
C SUBROUTINES REQUIRED
C NONE
C
C DIMENSION IA(2)
C
COMMON/DEVICE/LKB,LTP,LLP,LDIS,LTO,LPLOT,LXXX,LDK1,LDK2,LDK3,LDK4
COMMON/UPNTS/IDUM(5),IRGTOP,ITOP
COMMON/UDATA/ISTACK(25),IDUM3(5000)
C
DATA MAXSTK/25/
C
N=1
GO TO 100
C
C ENTRY UNSTK2(IOPND1,IOPND2)
C
C SAME AS UNSTK BUT FOR 2 ITEMS
C (IOPND1 IS TOP ELEMENT)
C
N=2
C
100 DO 200 I=1,N
    IF(ITOP.LE.0) GO TO 91
    IA(I)=ISTACK(ITOP)
    ITOP=ITOP-1
    IF(IA(I).LT.0) THEN
        IR=-3-IA(I)
        IF(IR.NE. IRGTOP-1) GO TO 93
        IRGTOP=IR
    ENDIF
200 CONTINUE
    GO TO (210,220),N
C
210 IOPAND=IA(1)
    RETURN
C
220 IOPND1=IA(1)
    IOPND2=IA(2)
    RETURN
C
C ENTRY STACK(IOPND)
C
C STACKS OPERAND ADDRESS IOPND AND MAKES SURE THAT IF
C IT IS A REGISTER THE REGISTER STACK IS UPDATED
C
IF(ITOP.GE.MAXSTK) GO TO 92
ITOP=ITOP+1
ISTACK(ITOP)=IOPND
C
IF(IOPND.LT. 0) THEN
    IR=-3-IOPND
    IF(IR.NE.IRGTOP .AND. IR.NE.IRGTOP-1) GO TO 93

```

```

        IRGTOP=IR+1
    ENDIF
    RETURN
C
C
91  WRITE(LTO,191)
191 FORMAT(' **VCODE** UNSTK:OPERAND STACK UNDERFLOW')
    STOP
C
92  WRITE(LTO,192)
192 FORMAT(' **VCODE** STACK: OPERAND STACK OVERFLOW')
    STOP
C
93  WRITE(LTO,193)
193 FORMAT(' **VCODE** STACK: REGISTER ASSIGNMENT ERROR')
    STOP
C
    END

```

```

SUBROUTINE VBYTE(IBYTE)
C
C      WRITES BYTE IBYTE IN NEXT FREE BYTE IN ABS CODE
C
C      AUTHOR: TOMMY ESSEBO 1980-11-11
C
C      BYTE IBYTE,ICODE,BYTES(4)
C
C      COMMON/UDATA/IDUM1(25),ICODE(20000)
C      COMMON/UPNTS/IDUM2(4),IPCNT,IDUM3(2)
C      COMMON/VXCLOG/LUNLOG,LOGSRC,LOGPSE,LOGINS,LOGHEX
C
C      EQUIVALENCE (IBYTES,BYTES(1))
C
C      IF(LOGHEX .NE. 0) WRITE(LUNLOG,1000)IPCNT,IBYTE
1000  FORMAT(13X,Z4.4,2X,Z2.2)
C      ICODE(IPCNT)=IBYTE
C      IPCNT=IPCNT+1
C      RETURN
C
C      ENTRY VWORD(IWORD)
C
C      SAME FOR WORD (2 BYTES) IWORD
C
C      IF(LOGHEX .NE. 0) WRITE(LUNLOG,1100)IPCNT,IWORD
1100  FORMAT(13X,Z4.4,2X,Z4.4)
C      IBYTES=IWORD
C      DO 20 I=1,2
C      ICODE(IPCNT)=BYTES(I)
20    IPCNT=IPCNT+1
C      RETURN
C
C      ENTRY VLONGW(ILONGW)
C
C      SAME FOR LONGWORD (4 BYTES) ILONGW
C
C      IF(LOGHEX .NE. 0) WRITE(LUNLOG,1200)IPCNT,ILONGW
1200  FORMAT(13X,Z4.4,2X,Z8.8)
C      IBYTES=ILONGW
C      DO 40 I=1,4
C      ICODE(IPCNT)=BYTES(I)
40    IPCNT=IPCNT+1
C      RETURN
C
C      END

```

FILE: VINSTX.FOR

GENERATES MACHINE CODE FOR VAX-11

AUTHOR: TOMMY ESSEBO 1980-11-11

SUBROUTINES REQUIRED

VBYTE VWORD VLONGW

PSPACE PHOLL PINT

SUBROUTINE VINSTO(OP1)

BYTE BYTADR(4),ICODE

INTEGER OPTAB(5,27),BUFF(14),OP,OP1,OP2,OP3,OP4

DIMENSION INTOPA(3)

COMMON/UDATA/IDUM1(25),ICODE(20000)

COMMON/UPNTS/L,IDUM2(3),IPCNT,IDUM3(2)

COMMON/UTEMPC/IDUM4(2),ITMPAD

COMMON/VXCLOG/LUNLOG,LOGSRC,LOGPSE,LOGINS,LOGHEX

EQUIVALENCE (INTADR,BYTADR(1))

DATA OPTAB /

OPCODE	NR OF OPERANDS	CASE	MNEMONIC
--------	----------------	------	----------

* '00000050'X	, 2 , 1 ,	4HMOVF,4H	,
* '000000DD'X	, 1 , 1 ,	4HPUSH,4HL	,
* '000000D4'X	, 1 , 1 ,	4HCLRF,4H	,
* '00000052'X	, 2 , 1 ,	4HMNEG,4HF	,
* '0000003C'X	, 2 , 7 ,	4HMOVZ,4HWL	,
* '00000051'X	, 2 , 1 ,	4HCMPF,4H	,
* '00000053'X	, 1 , 1 ,	4HTSTF,4H	,
* '00000040'X	, 2 , 1 ,	4HADDF,4H2	,
* '00000041'X	, 3 , 1 ,	4HADDF,4H3	,
* '00000042'X	, 2 , 1 ,	4HSUBF,4H2	,
* '00000043'X	, 3 , 1 ,	4HSUBF,4H3	,
* '00000044'X	, 2 , 1 ,	4HMULF,4H2	,
* '00000045'X	, 3 , 1 ,	4HMULF,4H3	,
* '00000046'X	, 2 , 1 ,	4HDIVF,4H2	,
* '00000047'X	, 3 , 1 ,	4HDIVF,4H3	,
* '000000CA'X	, 2 , 1 ,	4HBICL,4H2	,
* '000000CB'X	, 3 , 1 ,	4HBICL,4H3	,
* '000000DF'X	, 1 , 6 ,	4HPUSH,4HAL	,
* '00000014'X	, 1 , 2 ,	4HBGTR,4H	,
* '00000018'X	, 1 , 2 ,	4HBGEQ,4H	,
* '00000019'X	, 1 , 2 ,	4HBLSS,4H	,
* '00000011'X	, 1 , 2 ,	4HBRB ,4H	,
* '00000031'X	, 1 , 3 ,	4HBRW ,4H	,
* '000000E1'X	, 3 , 8 ,	4HBBC ,4H	,
* '000000E0'X	, 3 , 8 ,	4HBBS ,4H	,
* '00000005'X	, 0 , 4 ,	4HRSB ,4H	,
* '000000FB'X	, 2 , 5 ,	4HCALL,4HS	/

DATA NROP/27/

NADR=0

OP=OP1

```

GO TO 100
C
ENTRY VINST1(OP2,IAD1)
NADR=1
INTOPA(1)=IAD1
OP=OP2
GO TO 100
C
ENTRY VINST2(OP3,IAD2,IAD3)
NADR=2
INTOPA(1)=IAD2
INTOPA(2)=IAD3
OP=OP3
GO TO 100
C
ENTRY VINST3(OP4,IAD4,IAD5,IAD6)
NADR=3
INTOPA(1)=IAD4
INTOPA(2)=IAD5
INTOPA(3)=IAD6
OP=OP4
GO TO 100
C
ENTRY VJUMP(JADR)
C
C     INSERTS CURRENT ADDR AT ADDRESS JADR IN ABS. CODE
C     (BRW INSTRUCTION JUMPS ONLY)
C
INTADR=IPCNT-JADR-3
IF(LOGINS .NE. 0) WRITE(LUNLOG,1000)IPCNT,JADR
1000 FORMAT(6X,I5,8X,' JUMP FROM:',I5)
IF(LOGHEX .NE. 0) THEN
    JADR1=JADR+1
    WRITE(LUNLOG,1100)JADR1,INTADR
1100 FORMAT(13X,Z4.4,2X,Z4.4)
ENDIF
DO 40 I=1,2
40  ICODE(JADR+I)=BYTADR(I)
RETURN
C
ENTRY VBRANC(IFROM,ITO)
C
C     INSERTS BYTE ADDRESS IN CODE
C
C     ITO-IFROM IS INSERTED AT BYTE IFROM-1 IN ICODE
C
INTADR=ITO-IFROM
ICODE(IFROM-1)=BYTADR(1)
C
IF(LOGINS .NE. 0) THEN
    IBR=IFROM-2
    WRITE(LUNLOG,1200)ITO,IBR
ENDIF
1200 FORMAT(6X,I5,8X,' BRANCH FROM:',I5)
IF(LOGHEX .NE. 0) THEN
    IBR=IFROM-1
    WRITE(LUNLOG,1300)IBR,INTADR
ENDIF
1300 FORMAT(13X,Z4.4,2X,Z2.2)
RETURN
C

```

```

C      -----
C
C      FIND OPERATION
100  DO 120 IOP=1,NROP
      IF(OP .EQ. OPTAB(1,IOP)) GO TO 140
120  CONTINUE
      GO TO 901
C
C      CHECK NR OF OPERANDS
140  IF(NADR .NE. OPTAB(2,IOP)) GO TO 902
      IF(LOGINS .NE. 0) THEN
          IP=1
          CALL PSPACE(IP,BUFF,56)
          IP=7
          CALL PINT(IP,BUFF,IPCNT)
          IP=19
          CALL PHOLL(IP,BUFF,OPTAB(4,IOP),8)
          IF(NADR .EQ. 0) GO TO 190
          DO 180 IOPAND=1,NADR
              IP=18+9*IOPAND
              IAD=INTOPA(IOPAND)
              IF(IAD .LT. 0) THEN
C
C                  REGISTER OPERAND
                  CALL PHOLL(IP,BUFF,4HR,4)
                  CALL PINT(IP,BUFF,-1-IAD)
              ELSE
                  DO 150 I=1,20
                      IF(ITMPAD+(I-1)*4 .EQ. IAD) GO TO 160
150             CONTINUE
C
                  CALL PINT(IP,BUFF,IAD)
                  GO TO 170
C
C                  TEMP. VAR OPERAND
160             CALL PHOLL(IP,BUFF,4HTMP$,4)
                  CALL PINT(IP,BUFF,I)
C
170             CONTINUE
                  ENDIF
180             CONTINUE
190             WRITE(LUNLOG,1400)BUFF
1400  FORMAT(1X,14A4)
                  ENDIF
C
C      WRITE OPCODE
      CALL VBYTE(OPTAB(1,IOP))
C
C      EVALUATE OPERANDS
      GO TO(200,250,300,350,400,450,500,550),OPTAB(3,IOP)
C
C      CASE 1: ABSOLUTE OR REGISTER MODE OPERANDS
200  DO 220 I=1,NADR
      IF(INTOPA(I) .LT. 0) THEN
C
C          REGISTER ADDRESS
          CALL VBYTE(79-INTOPA(I))
      ELSE
C
C          ABS. ADDRESS
          CALL VBYTE(159)

```



```

        CALL VLONGW(INTOPA(1))
    ENDIF
220  CONTINUE
    RETURN

C
C      CASE 2: BYTE BRANCH
250  CALL VBYTE(INTOPA(1))
    RETURN

C
C      CASE 3: WORD BRANCH
300  IF(INTOPA(1) .NE. 0) GO TO 903
    CALL VWORD(INTOPA(1))
    RETURN

C
C      CASE 4: NO OPERAND
350  RETURN

C
C      CASE 5: CALLS
400  CALL VBYTE(INTOPA(1))
    CALL VBYTE(159)
    CALL VLONGW(INTOPA(2))
    RETURN

C
C      CASE 6: PUSHAL
450  CALL VBYTE(159)
    CALL VLONGW(INTOPA(1))
    RETURN

C
C      CASE 7: MOVZWL
500  CALL VBYTE(143)
    CALL VWORD(INTOPA(1))
    CALL VBYTE(159)
    CALL VLONGW(INTOPA(2))
    RETURN

C
C      CASE 8: BBC,BBS
550  CALL VBYTE(0)
    CALL VBYTE(159)
    CALL VLONGW(INTOPA(2))
    CALL VBYTE(INTOPA(3))
    RETURN

C
C      ERRORS
C
C      ILLEGAL OPCODE
901  IERR=1
    GO TO 999

C
C      ILLEGAL NR OF OPERANDS
902  IERR=2
    GO TO 999

C
C      ILLEGAL BRANCH ADDRESS
903  IERR=3
    GO TO 999

C
999  WRITE(6,1999),IERR
1999 FORMAT(' VINSTX ERROR:',I2)
    RETURN

C
    END

```