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Service Restoration In Electric Power Distribution Networks

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SERVICE RESTORATION
IN
ELECTRIC DISTRIBUTION
NETWORKS

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ABSTRACT

Automatic service restoration in distributional networks has been studied. After a thorough examination of the literature in this field, one algorithm was selected and implemented as a Fortran program. This implementation was then proved to be effective. Further improvements to the algorithm was introduced and in three cases these proved to be successful. For example: An operational aid indicating roughly the order of the switching operations was added, the number of switching operations decreased in some cases and the the priority of the zones was made more important.

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1. BACKGROUND

1.1 TASK

The task has been to study the existing methods of automatic service restoration in distribution networks. Furthermore one algorithm should be selected, and implemented in Fortran.

1.2 HISTORY

Automatic service restoration is a new area in electric power distribution. Nothing was published in this area until 1980. In 1980 Ross (1) defined the service restoration problem as the minimization of the weighted sum of unserved customer hours and unserved energy subject to line capacity constraints, transformer constraints and voltage drop constraints. Since then different methods have been developed. Today two different approaches stand out, one mathematical and one based on expert system technique.

1.3 DISTRIBUTION NETWORKS

The distribution network has a radial structure, looking similar to that of a tree. In the feeding substation, there is a transformer. The transformer is connected to a feeder via a bar and a feeder breaker. Branchpoints give the treelooking structure. Several switches (cut switches or tie switches) are included, creating zones (or nodes or sections) in the network. These sectionalizing switches makes it possible to isolate a fault so that a repair can be done. Because of the radiality of the network most of the faults not only causes the faulted zone to be deenergized but also all the nonfaulted zones further down viewed from the transformer. These zones can often be supplied from other transformers. Therefore service restoration is needed.

Distribution networks are often large. The normal size is about 100-1000 feeders, each one consisting of 10-500 nodes (varying much between different power companies). The voltage in these networks is in the range from 4-43.5 kV.

1.4 RESTORATION

When a fault has occurred a protective relay trips the circuit breaker. Next the fault has to be identified. One way to do this is by first opening all the sectionalizing switches and then closing the circuit breaker. The switches are then reclosed one by one in until the circuit breaker is tripped again. The faulted zone is identified and has to be isolated. At this point the operator must decide a sequence of switching operations that have to be performed

in order to restore the nonfaulted deenergized zones, without causing any violation. This is done by hand calculation based on a loads forecast or by the operators' experience. Often is the sequence thus produced not optimal. A better solution would certainly be reached if a computer would produce it.

1.5 OPTIMIZATION ASPECTS

Studies of the service restoration problem show that it is a large scale combinatorial optimization problem. The only known method for solving combinatorial optimization problems is to test all possible solutions. Because of the size of the problem this solution often cannot be reached in the lifetime of a computer. In a network containing 1000 switches, there are 10^{300} possible solutions.

The most common method for solving large scale combinatorial optimization problems is to use a heuristic search method. This means that knowledge about the problem is used to guide the search. In this case a possible rule might be that all the circuit breakers are left closed unless the fault has occurred in a transformer or a bus. By applying rules like this the number of possible solutions reduce and the execution time becomes reasonable. By rejecting possible solutions that seem to lead to a non-optimal solution one may lose the true optimal solution. Instead an approximate optimal solution is reached. This solution is often so good that it is unnecessary to find the true optimal solution. Such is the case in service restoration. Heuristic search methods are suitable both for ordinary programming languages such as fortran and symbol manipulating languages such as lisp or prolog.

2. INFORMATION RETRIEVAL

During my information retrieval I found two main approaches, which will be dealt with separately.

2.1 EXPERT SYSTEM APPROACH

The most common approach to service restoration is an expert system although the other methods are faster. This is due to the 'higher' programming level in lisp or prolog. It is easier to make changes in your algorithm using these languages.

Morelato (2) used the expert system approach providing a framework for a whole family of algorithms concerning distribution network automation. He proposed a true expert system search approach, using a decision tree to solve different problems such as service restoration, the minimization of losses and costs, load balancing or combinations of the above. In fact, service restoration is a special case of load balancing. If the deenergized zones are connected to a support feeder creating a violation and then applies the load balancing algorithm, a service restoration algorithm would have been used. Any violations that may exist have to be cut off in the end. In an example he demonstrated how his service restoration algorithm works. This approach didn't seem right for fortran programming and besides there were some other weaknesses. For one thing voltage drop limits were discarded.

Lee and others(3) propose an algorithm that seems to be one of the best. They have certainly the highest level of ambition. Aspects like the service crews availability, traffic situations and such things are mentioned but not included in the algorithm. The algorithm is an extension of an earlier work (4) by the same authors. In their latest paper the authors criticize the mathematical approach for not considering such important factors as priority of zones (A hospital ought to have higher priority than a sports arena) and availability of crews. The algorithm is not a search algorithm. Instead it is something of trial and error. In short, the algorithm is as follows:

- * Put all the deenergized non faulted zones in a list
- * Determine priority of zones
- * Try to restore the zone with highest priority first

This is done until no more zones can be restored. The restoration is done in three steps.

- * First a whole group of feeders is restored. If the restoration causes a violation, the restoration is abandoned and the second step is tried.
- * If the first step fails, only single zones are restored. If the restoration causes a violation, the restoration is abandoned and a third step is tried.
- * The third step first tries to transfer load from the feeder that will take up the restored load so that when doing so, it will not be violated. Then the restored load will be transferred.

The last two steps seem to be quite simple but they are built by 250 rules. The proposed system may be used in a partly automated network (Some switches can be remote-controlled). This causes some switches to be preferred when it comes to switching operations and therefore the remotely controlled switches get a higher priority. The system not only produces a sequence of switching operations but also indicates the order of them. The execution time seems to be satisfactory (less than 5 seconds).

2.2 MATHEMATICAL APPROACH

This approach is the one used in the few existing physical test systems.

In (5) Castro and others show the results from a simulation of a proposed system. The simulation is a part of a project discarding economics to create a more effective distributional network. Communication links are thought to be installed in order to guide switches by remote control. He gives only hints of what the system contains and hardly anything of how it works. A major part of the paper describes the advantage of using a switchtable (an interface between pure algorithms and the real world). The system is thought to have two functions for network control.

First a load balancing algorithm which iterates a solution. Only one switch operation is suggested at a time until the network is balanced. If that is the case, the risk of overload in the network is reduced.

The second function is service restoration in case of emergency. The basic difference between these two algorithms according to Castro is that the service restoration produces the whole sequence of switching operations at one time and does not wait for the first operation to be completed.

This paper is perhaps more interesting for those who deal with the problem of installing a commercial system.

In (6) Kato and others also handles the case of a real system. Kato works for Tokyo Electric Power Co. and in the paper he describes their installed test system in Ginza. (It was installed in 1981) Much is said about communication etc. but a paragraph about service restoration is found.

Also here only hints are given about the algorithm. It works in two steps.

First all the loads that can be transferred directly without causing any violations are transferred. Then loads in those feeders who will be violated by further restoration is transferred to other feeders before they are able to take up new restored load.

Test runs (The system was put into trial application in 1984) show that non-faulted zones could be restored in five minutes, within about a tenth of the time for manual restoration.

Aoki and others (7),(8) have produced several papers in the field. Obviously he has studied Lee (3) because (7) is an improved algorithm compared with (8). He has introduced priority for the zones thereby showing that it is as easy to use priority in the mathematical approach as in the expert system approach. There are more improvements. Voltage drop constraints are included as well as the powerful effective gradient method (9),(10). The algorithm works in four steps.

1. After fault isolation the deenergized zones will be restored to the feeder with the largest violation margin, discarding any violations.

2. If violations occurred these are tried to be removed by transferring loads to adjacent feeders until no violations exist or no more load transfers can be made without creating new violations.
3. If violations still exist loads will be disconnected until no violations exist.
4. In the fourth step the previous deenergized loads (in step 3) will be reenergized one by one and step 2 will be tried once again among former violation feeders. If it fails, the load will be disconnected again.

The effective gradient method is used to select in which order the loads are tried to be transferred or cut off. This is the only algorithm which has been tested on a large network (1188 nodes) and it produced a solution in less than a second.

2.3 SELECTION

There were two main candidates for implementation, Aoki (7) and Lee(3). Lee seemed to be the best in some respects and Aoki in others. The greatest advantage with Lee's algorithm is that the priority of zones will be determined by the operator during execution and interpretations of the priority-measurements are postponed until this moment. (No file containing predecided priority exists). On the other hand seemed this advantage possible but not suitable for implementation in Fortran.

Aoki's algorithm on the other hand was obviously possible to implement in Fortran and since it had no real disadvantages in any respect, this algorithm was selected.

3. IMPLEMENTATION

3.1 ALGORITHM

As it was previously stated the algorithm solves the problem in four steps. Below follows a flowchart of the algorithm

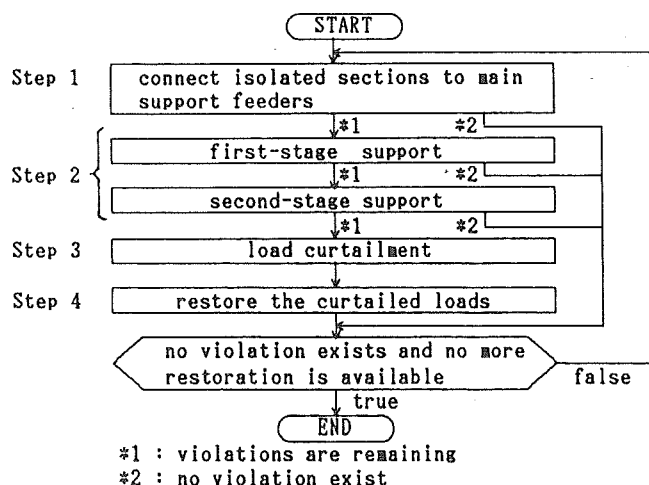


figure 1 (by Aoki). Flowchart of algorithm.

Step 1

First, the fault is isolated. This creates groups of deenergized sections. If a fault occurs in a branch- point there will be more than one group. For each of these groups all possible connections are examined. The best connection, that is the one to a feeder with the biggest margin of getting violated, is chosen and the group is reenergized.

If the connection causes a violation, the magnitude of these violations is stored in a violation vector. For each feeder both voltage drop violations and line capacity violations are stored. The violation vector is later used in order to decide between which feeder pair a load transfer should be performed.

If no violations occurred the violation vector equals zero and a feasible solution is found. The algorithm stops.

Step 2

Step 2 is only used when at least one feeder is violated. It has two principal parts, first and second stage support. Both these parts try to eliminate violations by transferring load from a violated feeder to a nonviolated.

In the first stage support, load is tried to be transferred from a violated feeder to a support feeder if it is possible. One cut switch has to be chosen for load transferring and the algorithm first chooses the one that will reduce the violation the most. This is done by solving the following optimization problem.

$$\begin{aligned} & \text{Max} \sum_{j \in I_b} \frac{h_j}{\alpha_j(H_j + \beta)} y_j \\ & \text{Sub.to} \sum_j A_{ij} y_j \leq b_i (i = 1, 2, \dots) \end{aligned}$$

The variable y_j is a 0-1 variable. It indicates the switch status. I is the set of possible cut switches that may be used for load transfer. H_j stands for remaining violation while h_j stands for the effective withdrawal of the violation. A is the capacity of the transformer, α_j the node priority and β is only included for avoiding zero divide. This problem is solved by using the dual effective gradient method (9),(10) which states that the order in which the load transfers are to be tried, is in descending order of the measure

$$\frac{h_j}{\alpha_j(H_j + \beta)}$$

This procedure is repeated until no more violations exists or until no more transfers can be performed without causing any violations in the support feeder. A flowchart of the first stage support is given in figure 2.

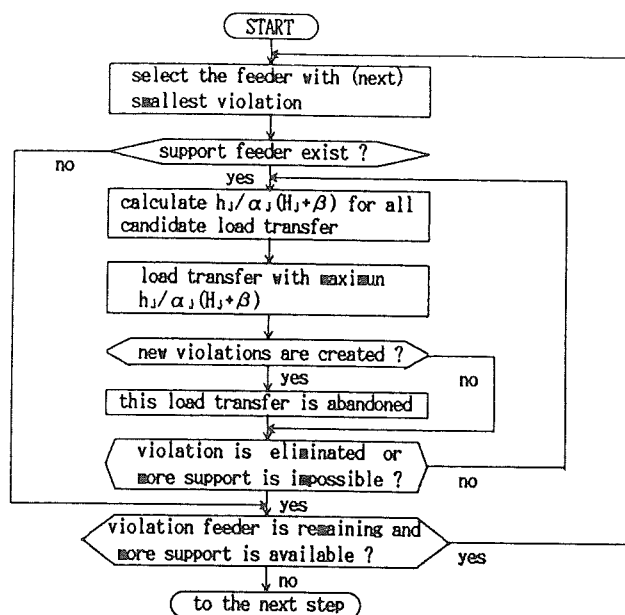


figure 2 (by Aoki). Flowchart of first stage support

The second stage support is only tried when the first stage support fails to transfer load. It works similarly to the first stage support. The second stage support begins with a modified first stage support allowing violations in the support feeder. Then this violation is eliminated by performing a first stage support on the violated support feeder. Aoki suggested a slightly different measure for the first stage support of the violated feeder (Not the violated support feeder). He suggests

$$\frac{h_j}{\alpha_j(a_j + \beta)}$$

Here a_j is the load magnitude of the transferred node. The second stage support performed until violations in the feeder is eliminated or until no secondary support feeders is able to take up any more load. In figure 3 a flowchart of the second stage support is found.

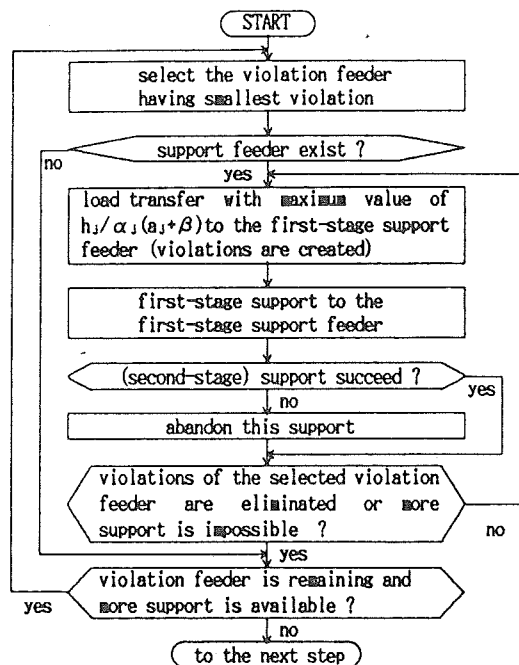


figure 3 (by Aoki) flowchart of second stage support

These two procedures are closely associated with each other and because of the similarity, a way of implementing them that could be hard to understand was selected. Some comments are needed.

First the candidate for a load transfer is to be selected. A file containing the pair of switches that is involved in the load transfer, the transferred load and an estimate of the reduction of the violations is stored in a file by a procedure called stage. This file is read by the procedure 'stage1' which calculates the different measures and selects a pair of switches able to perform a load transfer.

The two feeders involved in the load transfer is identified and all the different currents, node voltages and other data that could be changed due to a load exchange is first stored in a file called history. Then the actual load exchange is performed and new physical data is computed. If no new violations occur, all the 'backup'-files are deleted.

If new violations is created the former 'state' is read back in the net and a new trial is made until success or no more possible exchanges may be performed. Then 'stage1' terminates with ok=false, keeping all the 'backup'-files.

'Stage2' then calls 'stage1' with ok=false . This causes 'stage1' to perform a load exchange discarding violations and as an outparameter you get 'outfee' which is the violated feeder. Then an ordinary stage1 is performed for this feeder.

The 'stack of backupfiles' thus created is difficult to handle and if someone ever needs to make any changes in these procedures they are advised to think at least twice, before making any changes in order to fully understand them. If any bad instruction occurs in these procedures, the algorithm will produce undesired solutions, and the fault might be hard to detect.

Step 3

If violations still remain after step 2 then step 3 is executed. In step 3 a load curtailment is performed. The dual effective gradient method is also this time used to produce a measurement helpful in selecting the end section that will be disconnected. The measurement will be

$$a_j(H_j + \beta)$$

For every turn of step 3 the section with the smallest value of the above will be selected. A small load causing a small remaining violation is chosen. When no more violations remain, step 4 is executed

step 4

In step 4 the previously curtailed loads are tried to be restored. This can work if more than one feeder were violated. Then there is a possibility of restoring loads between two former violation feeders. The dual effective gradient method produces this time the following measure

$$\frac{a_j}{(|R_j| + \beta)}$$

R_j is the violation vector. When the cutswitch j is closed, one feeder will be violated.

The curtailed sections will one by one be reenergized in decreasing order of the measurement above and a first stage support will be performed. If the first stage support fails the load will once again be disconnected. In figure 4, a flowchart of step 3 and 4 is found

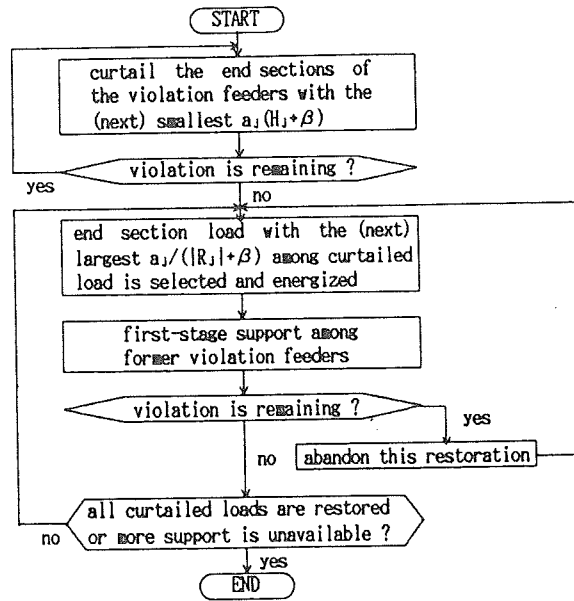


figure 4 (by Aoki) Flowchart of step 3 and 4

3.2 DATASTRUCTURE

In the datastructure six different types are found. These are

- 1: transformer
- 2: feeder
- 3: cable
- 4: branchpoint,
- 6: switch

Every type is represented by matrices containing pointers to the closest neighbours and physical data. More details about the data structure is found in the appendix.

3.3 NETWORK REPRESENTATION

There are a few simple rules that must be followed when representing a network in order to make the algorithm work.

- The feeder is only used to provide a useful entrance to the network and to attach the violation vector.
- There must be a switch between the feeder and the first node in the network.
- A cable is used to connect elements of the other types and can therefore not be attached to another cable.
- The switch also represents the circuit breaker in the algorithm in order to reduce the size of the code.

3.4 COMPUTATIONS

The following data is presupposed to be known:

1. The network's topology
2. Cable impedances
3. Node impedances based on a load forecast
4. Transformer voltages

The transmission net was presumed to be strong so that the transformer voltage was kept fairly constant even if the load increased. Series impedances were used as a model of impedances. Knowing this and the fact that the network has a radial structure, the voltage and the current in an arbitrary point in the network could be computed. In figure 5 a radial structure is seen.

First the total impedance must be computed. This is done by traversing the 'tree' of the transformer. When a 'leaf' is found the computation begins. On the way up the computations proceed storing parallel impedances in matrix BRAPH, when a branch point is reached. The search of a new leaf starts.

When the impedance is known the current is easily computed.

Then by applying the laws of Kirchhoff, all currents and voltages are easily computed by once again traversing the tree.

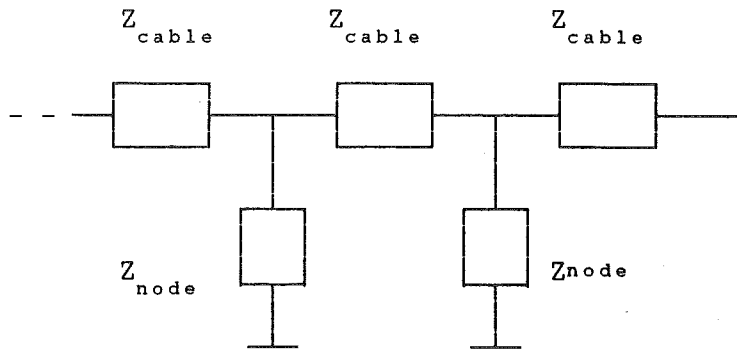


figure 5. Radial structure

This simple move around method seemed to be the best in this case. More sophisticated methods taught in circuit theory seemed to be more difficult to handle in this case because the number of impedances in the local network (the one that is connected to a single transformer) is continuously changing. Too much work would be spent on setting up matrices.

There are some problems concerning numerics. After a leaf node the current often differs from zero but the fault is not big enough for considering double precision (about 0.1 percent of transformer current). It is necessary to remember that the load forecast probably contains some error.

3.5 IMPROVEMENTS

During the implementation of the algorithm, some possible improvements were discovered.

The multi stage support may be generalised. A third stage support can be performed tolerating violations in the secondary support feeder and a first stage performed on this violated feeder. Aoki mentions this but consider it unnecessary because "the distribution systems are often designed so that all de-energized loads in an island can be restored without creating any serious violations" (Island means a group of de-energized sections). This is certainly the case when the transformer carries a normal load but faults also occur in case of a peak load. A third stage support was implemented.

The measurement given in both the first and second stage supports could be improved. The suggested improvement was

$$\frac{h_j}{\alpha_j(H_j + \beta)} * \frac{1}{a_j}$$

This causes a small load causing small remaining violation to be selected and this seemed to be an even better measurement than the one by Aoki.

In the load curtailment the priority is discarded by Aoki. This seemed to be wrong. A high priority zone should not be selected for load curtailment if

another candidate with lower priority exists. Therefore a possible improvement might be

$$\frac{a_j(H_j + \beta)}{\alpha_j}$$

In this algorithm a small numerical value of alfa indicates a high priority zone

The Aoki algorithm only indicates the switches that must be changed. It never tells you in which order to perform the switching operations. This may be fatal because it can create serious violations due to the fact that first violations are created and then these violations are tried to be reduced. Therefore some kind of indication of order is needed. It is only the switching operations that first reenergizes the isolated zones that can cause a violation, and therefore these must be performed in the end. An improvement of the algorithm would be that switching operations due to fault isolation would be indicated by the value 1. These operations must be performed first. Then the switching operations due to step 3 and 4 are performed so that no violations occur. Then those switching operations caused by the first and second stage support can be performed and eventually the switching operation caused by the reenergizing of isolated nodes.

3.6 TESTRUNS

In order to investigate if the suggested improvements really improved the result, a series of test runs was made on a small network taken from (7). This network is shown in figure 6.

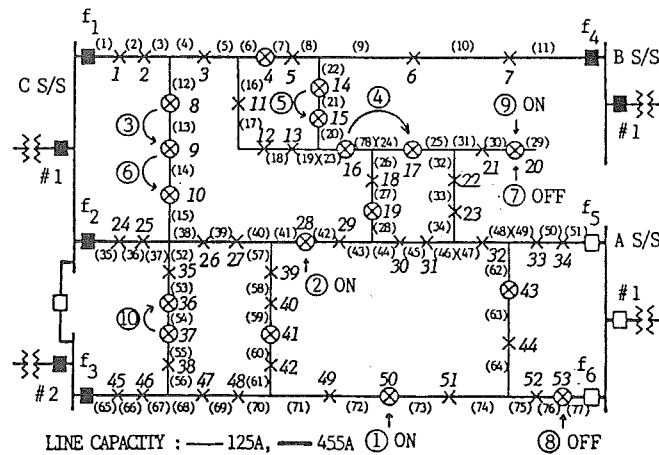


figure 6. Topology of test system (by Aoki)

A number of different qualities were demanded for a good solution.

- The execution time has to be rather short, less than 10 seconds is preferred.

- The number of switching operations need be as low as possible. This is due to the fact that today when almost every switch is manually controlled, it takes some time to perform the switching operations. (mostly due to the distances between the switches in the network)
- The restored load ought to be as big as possible. It may be difficult to measure the restored load because it depends on the node voltage and this can be changed during execution. Different feeders have different loads and therefore different voltage drops.

The plain Aoki algorithm was tested. The test showed that this implementation was equally effective as the Aoki implementation (Execution time ranging from 5 ms to 1 s) The number of switching operations and the amount of restored load seemed to be equivalent. This implementation was then used as a reference.

One sequence of test runs of specially selected faults was run. The faults were severe causing major violations in order to see the differences more clearly.

The third stage support was tested. It was obvious that the execution time rised dramatically to about 5 seconds. This however was not so serious.

It was soon noticed that the third stage support had little effect on the amount of restored load unless the network was unbalanced with all primary and secondary support feeders carrying major loads while other feeders only carried minor to normal load. This was hardly a surprising result.

More important was the number of switching operations. It increased by a factor of approximately 1.5 (In average). When the second stage support has failed nothing is tried in order to reduce the violation until the load curtailment take place.

In the load curtailment process only one single switch is changed. But if the third stage support is used successfully, then for every successful turn of this step six switches are changed. (Two switches connected to the primary support feeder, two to the secondary and two to the tertiary support feeder) The switches changed due to a successful third stage support can be seen in the figure 7.

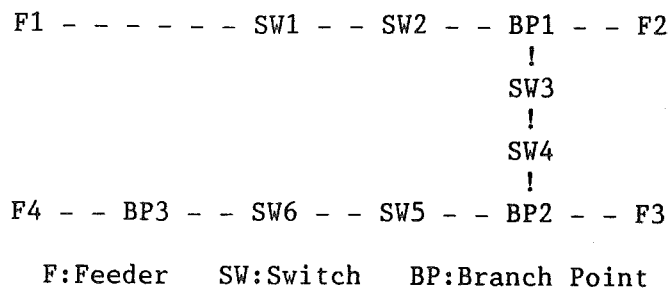


Figure 7. Topology typical for a third stage support.

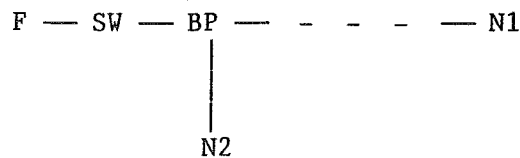
Between switches and branch points an arbitrary number of sections may be found. The third support is applied to feeder F1 transferring load to the primary support feeder F2(SW1,SW2), causing a violation. The violation is

eliminated by a further load transfer to feeder F3 (SW3,SW4), which in its turn has to transfer load to F4 (SW5,SW6) before it can take up any more load. The load transfers in this case are carried out by opening the odd switches and closing the even.

This increased number of switching operations made the suggested improvement less desired because the number of switching operations was already large (about 10-15). The 'optimal' solution must have a chance of being reached before the repair of the faulted zone is finished.

During test runs of the third stage support, the suggested improvement of the measurement in the second and first stage support was detected. When a fault caused serious violations, often the most severe line capacity violation was in the 'root' cable. This could be reduced by transferring any load in the tree, and sometimes a node which had not been deenergized was selected. This cannot be desired. A load transfer should always be made among former deenergized loads in order to best reduce the violation.

It is a known fact that in the beginning (before the reenergizing of isolated sections) the violation feeder was non violated. Therefore it is desired to perform the load transfer in the reenergized sections because these created the violation. The reason for violation must have a greater attention than the size of load and this is done by dividing the proposed measurement (Aoki) by H_j . An example is shown in figure 8.



F stands for feeder, SW for switch, BP for branch point and N for node

Figure 8. Typical topology

The node N1 has been among the reenergized nodes. When it was reenergized, it caused a violation in feeder F. In this case the voltage for N1 is 20 V below the voltage drop limit but N2 is not violated. In the root cable, that is the cable from the feeder to the branchpoint, the current exceeds the line capacity with 40 A. In this case the violation vector would be (20,40) Only these two candidates (N1,N2) for load transfer exist. Both nodes have the same priority. The load of N1 is 60 A and the load of N2 is 40 A. Computing the measurement (by Aoki) for these nodes gives:

$$N1: h_j = \sqrt{(2000)}, H_j = 0 \text{ and measurement} = 0.75$$

$$N2: h_j = 40, H_j = 4,7 \text{ and measurement} = 1.0$$

In this case N2 would have been selected. If the proposed measurement was to be used N1 would be selected instead. The new measurements would be:

N1: 74.0

N2: 0.21

It is obvious that the best choice for reducing the violations is transferring N1. Then both a voltage drop violation and a line capacity violation would be avoided. This was later proved to be an improvement, but not a major one. It reduced the number of switching operations with, in some cases, one or two. (The number of switching operations never increased but for the most cases it was the same.) The amount of restored load was unaffected (in average). The third suggested improvement to include the priority in the load curtailment was also an improvement in the sense that according to priority rules some previously wrong decisions were made right. An example of this is seen in figure 9.

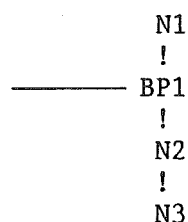


Figure 9. Typical configuration in case of load curtailment

Here both nodes N1 and N3 are violated with a voltage equally large, 20 V below the voltage drop limit. N1 is a high priority zone ($\alpha_{N1}=1$) while N2 is a low priority zone ($\alpha_{N2}=10$). The loads are: N1=10 A , N2=20 A.

Because of the equality in violation both the zones cause equal reductions in the violation vector. In this case the measurement proposed by Aoki will only depend on the size of the load. Therefore the high priority zone will be selected.

This could cause many damages in the society. An important operation in a hospital may have to be interrupted or postponed.

According to the proposed measurement(priority included) the low priority zone would have been selected for load curtailment. Then it is up to the operator to decide the different priorities. In this case the number of switching operations and amount of restored load must be of second interest. Anyway no major changes emerged. This is due to the fact that the service restoration problem is a combinatorial optimization problem. A decision that perhaps is seen as a 'bad' decision in the sence that it restores less load could in fact be better than the decision suggested by the heuristic search approach. It is the combination of decisions taken that lead to the optimal solution.

4. RESULTS

- * As a result of an information retrieval two major approaches to the service restoration problem was found (Expert System Approach and Mathematical Approach)
- * The mathematical approach is the only one which has been tested in a physical system
- * The result from this test system was promising.
- * The mathematical approach was selected due to the facts that code was to be written in fortran and the succesful tests that others had done in the field with this approach.
- * An algorithm was chosen for implementation, which fulfilled all possible requirements for a good performance such as a short execution time and a small number of switching operations.
- * This algorithm was then implemented in fortran.
- * The implementation was proven to be as effective as the original implementation.
- * Four possible improvements were suggested
- * These suggested improvements caused a minor improvement in performance of the algorithm
- * An operator aid was added. Not only the switching operations were presented but also a rough indication of the order of these operations was given.
- * In a physical system the use of a service restoration algorithm depends on the accuracy of the load forecast.

5. FUTURE

As it has been previously hinted, the service restoration algorithm is only one part of a group of algorithms thought to be used in an automatized distribution network in the future. The reason for the focus on the service restoration algorithm at the moment depends on the possibility to use it with locally controlled switches. It would be too expensive having service crews continuously driving around changing switches in order to improve the performance of the network. In case of an emergency however, it is essential to transfer load. In the future these other algorithms (the load balancing algorithm, the minimum loss and the minimum cost algorithm) as well as the service restoration algorithm will probably all be found in an ordinary automatized distribution network. Due to the high cost of installing communication links, hardly all sectionalizing switches will be remotely controlled. Instead, only those that often will be used for a load transfer will be supplied with remote control.

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7. APPENDIX

7.1 USER'S MANUAL

All the useful files are stored on node SNEV02 at ABB Network Control. The files will be found in directory EXARB ROOT: MBERGSTRAND .

In the file CONSTANTS.FOR are all the necessary constants declared and this file is included in all files in order to facilitate any changes in the structure.

This was the same reason for having the file COMMON.FOR . In this file all global variabls needed were declared . This file is also included whenever it was necessary.

The constants MAXNRT, MAXNRF, MAXNRC, MAXNRB, MAXNRN and MAXNRS have to be changed before the program is used on a new network.(nr of transformers, feeders,cables,branchpoints,nodes and switches) Another constant, BRANCH, may have to be changed. It indicates the maximum allowed branches for any branchpoint in the network.

When this is done all the files RESTORE, OUTFI, MOVING, TESTS, COMP, STEP1, STEP2, FAULT, FETCH and CHG must be recompiled and linked together. The main program is in file RESTORE.FOR.

Then an input file must be made, containing the physical data and the topology of the network.

All the different components (transformers,switches feeders,cables,branch points and nodes) must be numbered.

The program uses formatted input. The procedures of interest with appropriate format is located in the file OUTFI.FOR. The procedures are called WRITEF and INFI. Studies of the structure chart found later in the appendix is essential understanding WRITEF and INFI. The structure chart tells what the different elements in the matrices used stand for.

In the file the different types is in the order

- o Transformers
- o Feeders
- o Cables
- o Branch Points
- o Nodes
- o switches.

The format used is as follows:

- All real values are written using exponents (0.100E+01).
- First two records are needed to define the transformers (common to all the different transformers).
- The first contains the constant MAXNRT (MAXimum Number of Transformers ,integer) written in position 3-8
- The second contains the maximum voltage in the network (real) written in position 3-14

Then one record is used to represent every single transformer. This record is repeated MAXNRT times. It contains the following data:

- Maximum value of P in pos. 3-14 (real)
- Maximum value of Q in pos. 17-28 (real)
- A pointer to a cable in pos. 31-36 (integer)

The transformers, as well as any other types, have to be in order (Trafo 1, trafo 2...).

FEEDERS come after TRANSFORMERS

- This is followed by one record containing the constant MAXNRF written in pos. 3-14 (integer)

Then one record containing the topology for each single feeder is repeated MAXNRF times. It contains of:

- A pointer to the cable that leads towards the transformer in pos. 3-8
- A pointer to the cable that leads 'down' towards the nodes in pos. 11-16

Both are integer values

The next type in the data file is CABLE.

- First one record containing the constant MAXNRC, written in pos 3-8 must be made. (integer)

That is followed by seven records, containing the physical data and the topology for each single cable. These seven records are repeated MAXNRC times. The records are shown below:

Record 1. R pos. 3-14 real

Record 2. X pos. 3-14 real

Record 3. I_{max} pos. 3-14 real

Record 4. Type 1 pos. 3-8 integer

Record 5. Nr 1 pos. 3-8 integer

Record 6. Type 2 pos. 3-8 integer

Record 7. Nr 2 pos. 3-8 integer

Type and Nr are referring to the element, which the cable is connected to. BRANCH POINTS come after CABLES.

- One record containing the constant MAXNRB in position 3-8 (integer)

Then for each branch point:

- First the actual number of branches, less than the constant BRANCH, is written in one record in pos 3-8
- Then one record containing one pointer to a cable is written in position 3-8 for every branch.

This is repeated MAXNRB times.

An example:

In one network there are 2 branch points. The first branch point has 4 branches and the second has 3. The first branch point is connected to cables 6,7,8 and 9. The second is connected to cables 2,3 and 4. This is how it appears in the data file.

2
4
6
7
8
9
3
2
3
4

NODES come after BRANCH POINTS.

- One record containing the constant MAXNRN in pos. 3-8
- Then the voltage drop limit is written in one record in pos. 3-14. The voltage drop limit is an actual voltage and not a percentage of the transformer voltage.

Then three records is used to represent one node.

For each node:

Record 1.	Pointer to cable	pos. 3-8 (int)
	Pointer to cable	pos. 11-16 (int)
Record 2.	R (Resistance)	pos. 3-14 (real)
	X (Inductance)	pos. 17-28 (real)
Record 3.	Priority	pos. 3-14 (real)

As an example, the physical data and the topology for one node, connected to the cables 3 and 4 is seen below as it must be written in the data file. The node's impedance is (543,j310) and the priority is 1.0.

	3	4
0.54300E+03		0.31000E+03
0.10000E+01		

The above is repeated for each node.
SWITCHES come after NODES.

- One record contains the constant MAXNRS in pos. 3-8.
The topology come before the status.
One record contains the topology.
- A pointer to a cable in pos. 3-8
- A pointer to a cable in pos. 11-16
This is repeated for each switch.
Then follows the status for each switch in one record for each switch.
- Record: Switch status (0,1) in pos. 3-8

This is repeated MAXNRS times.

An example:

One network has two switches. The first switch is connected to cable 10 and 12, and it is opened (Status=0). The second switch is connected to cable 8 and 9, and it is closed (Status=1).

This is how it looks in the data file:

```
2
10
12
8
9
0
1
```

Following these instructions in order, a correct data file is produced.

All files of the type '.DAT' contains a network, with the name indicating whether the load is minor, medium or major.

During execution of the program all commands and answers must be written in CAPITAL LETTERS.

When the program starts its execution, first a check of the topology is made. Some faults can be detected and messages will be written on the screen telling the user what to do. The faults that will be detected is of the kind: "if you can 'go' to a neighbour but you can not go back to the starting point" it must be wrong in the network topology. Most of the faults are detected by this check but fault can also 'cooperate' so that two faults make it look like no fault. If errors of this kind occur in the network these faults most certainly are detected later causing messages like this: WRONG IN STEPDO.NODE, FORTRAN STOP.

If messages of this type appear, the program provides a way of easily finding these errors. When the first test of the topology is made, the program asks if the operator desires a visible track. If the answer is 'Y' the program will print out all the steps taken in the subroutine init. Then it is easy to follow the program until an error is found. Then the latest point must be remembered. In this point the error is found. Edit the data file and try again. If the topology is correct the text INIT PASSED will be written on the screen

Later in this section transcripts of possible test runs are showed with comments added afterwards.

The program will then ask if any changes are to be made in the network (concerning load). If you wish to do so the answer is Y. Then the program will ask if you wish to change the total load. If not, you have the possibility to change a single node impedance. If you choose to change the total load the program will ask for a factor with which every single node impedance in the network will be multiplied. As a guideline you can use 1.1 if you wish to reduce the load with 10 percent ($P = \frac{U^2}{R}$) but the system is not linear. In case of a change in a single node the old value of the impedance is written on the screen and then the new value is read. Then the program asks for a filename. The new load condition is saved for later use. The filename can be seven characters long. All names except HISTORY, STAGE and CURTAI are allowed. Files with these names are used internally. See also in the transcripts.

Then the program will ask for a point in the network where the fault has occurred. It can be any point in the network. A point is indicated by its type and number (both integer values). Available types are

- 1: Transformers
- 2: Feeders
- 3: Cables
- 4: Branch Points
- 5: Nodes
- 6: Switches

Then the algorithm starts working, continuously writing data on the screen. A careful study of decisions taken can be done. In the end a menu is presented containing execution time, the switches that have changed status, the order of these switching operations, deenergized load and restored load. Then the program asks if some more faults (A new test case) is desired and if so the program restarts.

7.2 EXAMPLES

Four different examples are shown in order to demonstrate how the program works.

TRAFOS IS OK
FEEDER IS OK

FAULTED NETWORK
BRANCH-POINT NR 1 IS CONNECTED TO CABLE NR 30

CABLE NR 30 IS CONNECTED TO
TYPE= 5 NR= 6
AND

TYPE= 6 NR= 4
FAULT IN BRANCH-POINT(4) OR IN CABLE(3) ?
THE FOLLOWING BRANCHES

- 1
- 2
- 30

THE FAULTING ONE IS (NOT) CONNECTED TO 30
THE BRANCH IS CONNECTED TO CABLE NR ?
BRANCHPOINTS OK

NODES OK
SWITCH OK
SWITCH OK

TRAFOS IS OK
FEEDER IS OK
BRANCHPOINTS OK
NODES OK
SWITCH OK
SWITCH OK

THE NETWORK IS CORRECT
DO YOU WISH TO HAVE A VISIBLE TRACK? Y/N
DO YOU WISH TO MAKE ANY CHANGES IN THE NETWORK?Y/N
INIT PASSED

FAULT IN

- 1: TRANSFORMER
- 2: FEEDER (MAIN SUPPORT)
- 3: CABLE
- 4: BRANCHINGPOINT
- 5: NODE

WHICH TYPE OF COMPONENT?

WHICH NODE?

ENTERS STAGE 1 2
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT, FEEDER 1.614994 6
IMPROVEMENT, FEEDER 0.2358537 1
IMPROVEMENT, FEEDER 0.2358537 5
SAVENE*****
2 6
SAVENE*****

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ABB Network Control
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T

ENTERS STAGE 1 2
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT, FEEDER 5.272868 6
IMPROVEMENT, FEEDER 0.2615825 1
IMPROVEMENT, FEEDER 0.2615825 5
SAVENE*****
2 6

SAVENE*****
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T
ENTERS STAGE 1 2
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT, FEEDER 6271.001 6
IMPROVEMENT, FEEDER 0.000000E+00 1
IMPROVEMENT, FEEDER 0.000000E+00 5
IMPROVEMENT, FEEDER 0.000000E+00 6
SAVENE*****
2 6

SAVENE*****
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T
ENTERS STAGE 1 2
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT, FEEDER 2805.759 6
IMPROVEMENT, FEEDER 0.000000E+00 1
IMPROVEMENT, FEEDER 0.000000E+00 5
IMPROVEMENT, FEEDER 0.000000E+00 6
SAVENE*****
2 6

SAVENE*****
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T
TOTAL COMPUTATION TIME OF ALGORITM 275 ms

THE FOLLOWING SWITCHES HAS CHANGED STATUS

37	4
38	3
45	1
46	1
50	3

TOTAL DEENERGIZED LOAD DUE TO FAULT
0.29190E+07

TOTAL RESTORED LOAD
0.29190E+07

ANOTHER FAULT? Y/N

FORTRAN STOP

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COMMENTS

In the previous example the possibility to discover a single fault was demonstrated.

The error in the topology was identified and then the data concerning the topology for the two possible faulted elements were written on the screen. The operator then identifies the faulted element. (In this case the branch point) All connections for this erroneous branch point are written in the screen and the faulted connection is marked by the text "THE FAULTING ONE IS NOT CONNECTED TO". The program then asks for the number of the cable (the right one) that it is connected to. Then the whole test is done again so that no more faults occurred due to this 'new' topology.

ABB Network Control

WHICH FILE IS TO BE READ?

TRAFOS IS OK
 FEEDER IS OK
 BRANCHPOINTS OK
 NODES OK
 SWITCH OK
 SWITCH OK

THE NETWORK IS CORRECT
 DO YOU WISH TO HAVE A VISIBLE TRACK? Y/N

DO YOU WISH TO MAKE ANY CHANGES IN THE NETWORK?Y/N

POINT IN IS 1 1
 POINT OUT FROM STEP DOWN IS 6 57

INDATA NEWDIR (POINT,DIR,BRA)
 6 57 0 4

OUTDATA NEWDIR (POINT,DIR,BRA)
 6 57 0 4

INDATA TKESTP (POINT,OLDP,DIR,BRA)
 6 57 1 1 0 4
 POINT IN IS 6 57
 POINT OUT FROM STEP DOWN IS 4 1

OUTDATA TKESTP
 4 1 6 57 0 4

INDATA NEWDIR (POINT,DIR,BRA)
 4 1 0 4

OUTDATA NEWDIR (POINT,DIR,BRA)
 4 1 0 4

INDATA TKESTP (POINT,OLDP,DIR,BRA)
 4 1 6 57 0 4
 POINT IN IS 4 1
 POINT OUT FROM STEP DOWN IS 2 1

OUTDATA TKESTP
 2 1 4 1 0 4

INDATA NEWDIR (POINT,DIR,BRA)
 2 1 0 4

OUTDATA NEWDIR (POINT,DIR,BRA)
 2 1 0 4

INDATA TKESTP (POINT,OLDP,DIR,BRA)
 2 1 4 1 0 4
 POINT IN IS 2 1
 POINT OUT FROM STEP DOWN IS 6 54

OUTDATA TKESTP
 6 54 2 1 0 4

INDATA NEWDIR (POINT,DIR,BRA)
 6 54 0 4

OUTDATA NEWDIR (POINT,DIR,BRA)

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ABB Network Control

```

INDATA NEWDIR (POINT, DIR, BRA)
    5          1          1          5

OUTDATA NEWDIR (POINT, DIR, BRA)
    5          1          1          5

INDATA TKESTP (POINT, OLDP, DIR, BRA)
    5          1          6          1          1          5
POINT IN IS          5          1
CABUP IN NODE IS          2
CABLE UP IS          3
POINT OUT FROM STEPUP IS          6          54

OUTDATA TKESTP
    6          54          5          1          1          5

INDATA NEWDIR (POINT, DIR, BRA)
    6          54          1          5

OUTDATA NEWDIR (POINT, DIR, BRA)
    6          54          1          5

INDATA TKESTP (POINT, OLDP, DIR, BRA)
    6          54          5          1          1          5
POINT IN IS          6          54
POINT OUT FROM STEPUP IS          2          1

OUTDATA TKESTP
    2          1          6          54          1          5

INDATA NEWDIR (POINT, DIR, BRA)
    2          1          1          5

OUTDATA NEWDIR (POINT, DIR, BRA)
    2          1          1          5

INDATA TKESTP (POINT, OLDP, DIR, BRA)
    2          1          6          54          1          5
POINT IN IS          2          1
POINT OUT FROM STEPUP IS          4          1

OUTDATA TKESTP
    4          1          2          1          1          4

INDATA NEWDIR (POINT, DIR, BRA)
    4          1          1          4

OUTDATA NEWDIR (POINT, DIR, BRA)
    4          1          0          5

INDATA TKESTP (POINT, OLDP, DIR, BRA)
    4          1          2          1          0          5
POINT IN IS          4          1
WRONG IN STEPDO.BRA

```

FORTRAN STOP

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COMMENTS

In this example the possibility to discover cooperating faults was demonstrated.

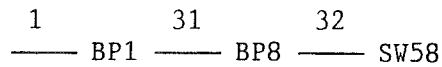


Figure 10. One part of the topology.

In branch point 1 (BP1) had a pointer to cable 32 been stored instead of a pointer to cable 31. This error was not detected by the first test. By demanding the program to show how it 'moves' (Answer Y to the question 'DO YOU WISH TO HAVE A VISIBLE TRACK') the last point before the program stops is shown in the screen. Then it is easy to make the proper change in the data file.

In this case the program stopped at point (4,1). That is branch point 1. Several pages in the middle were left out.

ABB Network Control

WHICH FILE IS TO BE READ?

TRAFOS IS OK
FEEDER IS OK
BRANCHPOINTS OK
NODES OK
SWITCH OK
SWITCH OK

THE NETWORK IS CORRECT
DO YOU WISH TO HAVE A VISIBLE TRACK? Y/N

DO YOU WISH TO MAKE ANY CHANGES IN THE NETWORK?Y/N

CHANGE TOTAL LOAD? Y/N

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CHANGE IN TYPE NR?

R IS 1009.800

X IS 408.0000

NEW R IS

NEW X IS

ON WHICH FILE DO YOU WISH TO STORE IT
INIT PASSED

FAULT IN

- 1: TRANSFORMER
- 2: FEEDER (MAIN SUPPORT)
- 3: CABLE
- 4: BRANCHINGPOINT
- 5: NODE

WHICH TYPE OF COMPONENT?

WHICH NODE?

TOTAL COMPUTATION TIME OF ALGORITHM 14 ms

THE FOLLOWING SWITCHES HAS CHANGED STATUS

41	4
42	1
48	1
49	1
50	4

TOTAL DEENERGIZED LOAD DUE TO FAULT
0.21720E+06

TOTAL RESTORED LOAD
0.21720E+06

ANOTHER FAULT? Y/N

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COMMENTS

In the third example the possibility to change a single node impedance was shown. Then a fault causing no violation at all was chosen to show how fast a solution can be reached.

ABB Network Control

TRAFOS IS OK
FEEDER IS OK
BRANCHPOINTS OK
NODES OK
SWITCH OK
SWITCH OK

THE NETWORK IS CORRECT
DO YOU WISH TO HAVE A VISIBLE TRACK? Y/N

DO YOU WISH TO MAKE ANY CHANGES IN THE NETWORK?Y/N

CHANGE TOTAL LOAD? Y/N
IMPEDANCES WILL BE MULTPLIED WITH?

ON WHICH FILE DO YOU WISH TO STORE IT
INIT PASSED

FAULT IN

- 1: TRANSFORMER
- 2: FEEDER (MAIN SUPPORT)
- 3: CABLE
- 4: BRANCHINGPOINT
- 5: NODE

WHICH TYPE OF COMPONENT?

WHICH TRANSFORMER?

```
ENTERS STAGE 1          3
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT,FEEDER  0.000000E+00          6
OK=. FALSE.
ENTERS STAGE 2
STAGE2 CALLS STAGE1 *****
ENTERS STAGE 1          3
OK IS FALSE
BUSY READING
STAGE1 WITH OK=F HAD AN UNSUCCESSFULL COMPLETION
EXIT STAGE 2 OK= F
ENTERS STAGE 1          5
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT,FEEDER  112015.5          6
IMPROVEMENT,FEEDER  13.42432          4
IMPROVEMENT,FEEDER  0.1634084          4
SAVENE*****
      5          6
SAVENE*****
CALLS RESNET*****
CALLS RESNET*****
IMPROVEMENT,FEEDER  112015.5          6
IMPROVEMENT,FEEDER  13.42432          4
IMPROVEMENT,FEEDER  0.1634084          4
SAVENE*****
      5          4
SAVENE*****
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T
ENTERS STEP 3
```

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NC 20000-8 40000 89-07 ABBELLS




```

ENTERS STEP 4
ENTERS STAGE 1          3
ENTERS STAGE*****
READS IN STAGE1 *****
IMPROVEMENT, FEEDER    4902.289          4
IMPROVEMENT, FEEDER    4902.289          5
IMPROVEMENT, FEEDER    0.000000E+00      6
SAVENE*****
      3          4
SAVENE*****
LEAVES STAGE1 AFTER A SUCCESFULL SESSION,OK=T
TOTAL COMPUTATION TIME OF ALGORITM 295 ms

```

THE FOLLOWING SWITCHES HAS CHANGED STATUS

```

      4      3
      8      3
      9      3
     11      3
     16      4
     37      4
     54      1
     57      1
     58      1

```

TOTAL DEENERGIZED LOAD DUE TO FAULT
0.99335E+07

TOTAL RESTORED LOAD
0.99335E+07

ANOTHER FAULT? Y/N

FORTRAN STOP

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COMMENTS

In the fourth example the total load was first changed providing the right answers in the proper places (Y to all but 'DO YOU WISH TO HAVE A VISIBLE TRACK?'). Then a serious fault was chosen (minor load). A fault in transformer 1 occurred. The topology of the test system is once again showed in figure 10. A closed switch is marked by 'x' and an open switch is surrounded by a circle

The fault is in this case isolated by closing the circuit breakers (switch 54,57,58 or #1,f1,f2 in the figure) Isolated trees are created with roots in nodes 1 and 35. These are marked (1), (35) in figure 11.

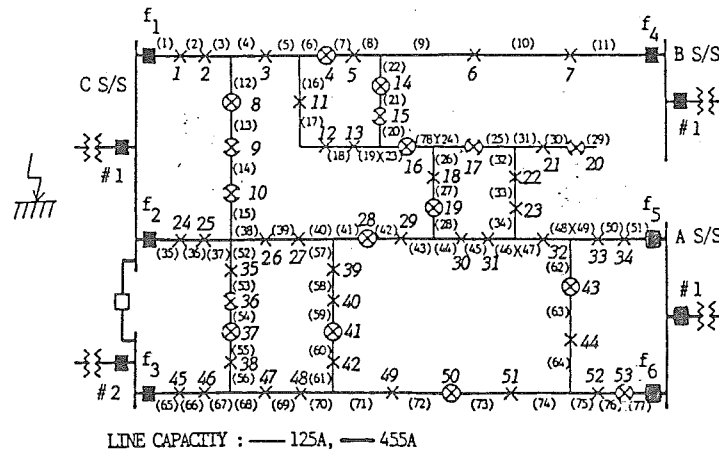


figure 11. Topology

Then for each tree all possible connections are examined. Possible switches for the tree with node 1 as the root are 4,14,16. Note that switch 8 is not a possible connection because this switch leads to another isolated section. Switch 16 is in this case the most suitable switch for reenergizing the 'tree' because feeder 5 has a bigger margin of getting violated than feeder 4. Switch 16 is then opened. The same examination is done for the tree with node 35 as root and in this case switch 37 is opened.

These two switching operations caused violations in feeders 3 and 5. First the violation in feeder 3 is tried to be decreased. The only possible load transfer is transferring node 72 from feeder 3 to feeder 6 (the switches 8,28 and 41 leads all to a violated feeder) but since this load transfer does not reduce the violation at all it is abandoned. There is a voltage drop violation in node 13. This violation would be eliminated by a load transfer of this node, but it is impossible to transfer this node since its neighbour's feeder also is violated. (The neighbour feeder is now feeder 5)

A first stage is then tried on feeder 5. The transcript shows that there are three possible switches for load transfer. These are switch nr: 43, 14, 4. Switch 43 cause the biggest reduction in the violation vector because Switch 43 is opened and switch 33 is closed but then feeder 6 would be violated. Instead the possibility to transfer load by opening switch 11 and closing switch 4 is

examined. This load transfer turned out to be successful. Feeder 5 is not violated any more.

The next step in the algorithm is the load curtailment. Feeder 3 is still violated and therefore node 13 is deenergized by opening switch 9. Then no violations at all exist and the fourth step is performed. Node 13 is reenergized by closing switch 9 and a first stage support is tried. The load transfer succeeds and node 13 is transferred to feeder 4. (Node 12 is energized by feeder 4 since the previous first stage support of feeder 5) Then all loads are restored and the algorithm stops.

The switching operations that must be performed is compared with the NON-FAULTED system. So in this case the circuit breakers are already changed. They were tripped by the fault.

Realizing this the operator can then command the switching operations to be executed in order.

7.3 OPTIMIZATION

PROBLEM FORMULATION

The service restoration is formulated by Aoki as follows:

$$\begin{aligned} & \text{Maximize } \sum_i \sum_{j \in J_{i0}} a_{ij} x_{ij} \\ & \text{Subject to} \\ & \text{(line capacity constraint)} \\ & \quad \sum_{j \in J_{ik}} a_{ij} x_{ij} \leq b_{ik} \\ & \text{(transformer capacity constraint)} \\ & \quad \sum_{j \in J_t} a_{ij} x_{ij} \leq b_t \\ & \text{(voltage drop constraint)} \\ & \quad \sum_{l \in T_e} \left(\sum_{q \in J_{il}} s_{iq} x_{ij} \right) z_{il} \leq V_{ie} \end{aligned}$$

Where,

- x_{ij} : such 0-1 variable as 0 if section j at feeder i is de-energized and 1 if it is energized. x_{ij} cannot arbitrarily be set to 1. $x_{ij}=1$ only if the adjacent section is energized and the sectionalizing switch connected to it is closed. x_{ij} is a function of the sectionalizing switch status.
- a_{ij} : load magnitude of section j at feeder i
- b_{ik} : line capacity at k -th point of feeder i
- b_t : transformer capacity of transformer t
- z_{il} : impedance of section l at feeder i
- s_{ik} : $s_{ik} = a_{ik}$ (if $k \neq 1$), $s_{ik} = a_{ik}/2$ (if $k=1$) as uniformly distributed load is assumed
- v_{ie} : voltage drop limit at the e -th end of feeder i
- J_{i0} : index set of de-energized load sections on feeder i
- J_{ik} : index set of load sections connected to the leaf side of point k of feeder i
- J_{il} : index set of load sections which exist at the leaf side of section l (included) of the feeder i
- J_t : index set of load sections ij connected to transformer t
- T_e : index set of load sections which exist at the trunk of the tree between bus and section e

EFFECTIVE GRADIENT METHOD

The effective gradient method is used by some economists dealing with problems such as: Which orders will not be served due to a limited resource. The effective gradient method is very simple to understand. It says that if you have to exit from a restricted area with an unlimited number of steps, but wish to stay as close to the border as possible and the distance to the permitted area is not known you should take the shortest step possible in that direction.

7.4 STRUCTURE CHART

The network is represented by the following matrixes where the first index points out the number of the component in the network. The following map shows where the data is stored for every component.

TRAFOS

P_{max}	Q_{max}	P	Q	V	$R_{network}$	$X_{network}$
-----------	-----------	---	---	---	---------------	---------------

MSFEED

Trafo Identity	Cable up	Cable down
----------------	----------	------------

CABTOP

First neighbour Type	Number	Second neighbour Type	Number
-------------------------	--------	--------------------------	--------

CABPHY

R	X	I_{max}	I
---	---	-----------	---

BRAPOI

Actual nr of branches	Branch up	Branch 1	Branch 2	...
--------------------------	--------------	-------------	-------------	-----

NODETO

Up	Cable 1	Cable 2	Feeder identity
----	---------	---------	-----------------

NODEPH

P	Q	V	R	X
---	---	---	---	---

SWITCH

Status	Up	Cable 1	Cable 2	Allowed change
--------	----	---------	---------	----------------

Some comments are required to understand the structure charts.

In up (cable up, branch up) the index pointing out the cable leading in the direction of the transformer is stored.

For open switches, up will be assigned a new value every time the switch will be 'visited'.

Allowed change is used to prevent a load transfer to the faulted zone.

In the program another matrix with three indexes, BRAPH is used to store the impedance one sees if one look down a branch. This matrix have matched the two first indexes with BRAPOI.

CC

RESTORE.FOR

CC

PROGRAM MAIN

CC

AUTHOUR MAGNUS BERGSTRAND

VERSION 1

DATE 1989-10-07

PURPOSE

WHEN A FAULT HAS OCCURRED IN AN ELECTRICAL DISTRIBUTION NETWORK THIS FAULT MUST BE ISOLATED. THE NON FAULTED ZONES CAN OFTEN BE RESTORED WITHOUT CAUSING ANY VIOLATIONS. THE PROGRAM SUGGEST SUCH A POSSIBLE RESTORATION

METHOD

THE METHOD IS DISCRIBED IN 'SERVICE RESTORATION IN ELECTRIC DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND
THE METHOD WAS FIRST DISCRIBED BY AOKI IN
'A NEW ALGORITHM FOR SERVICE RESTORATION IN DISTRIBUTION SYSTEMS'
IEEE/PES WINTER MEETING, NEW YORK 1989 PAPER 89 WM 035-2 PWRD

DESCRIPTION

FIRST IS THE FAULTED ZONE ISOLATED AND THE NON FAULTED ZONES CONNECTED TO AN ADJACENT FEEDER. IF THIS CONNECTION DOES NOT CAUSE ANY VIOLATION THE ALGORITHM STOPS AND A SOLUTION IS FOUND. IF VIOLATIONS OCCURED THESE ARE TRIED TO BE DECREASED IN STEP2. STEP2 HAS TWO PRINCIPAL PARTS THE FIRST AND SECOND STAGE SUPPORT. THE FIRST STAGE SUPPORT TRANSFERS LOAD TO AN ADJACENT FEEDER FROM A VIOLATED FEEDER NOT ALLOWING ANY NEW VIOLATIONS. THE SECOND STAGE SUPPORT TRANSFERS FIRST LOAD TO THE PRIMARY SUPPORT FEEDER ALLOWING VIOLATIONS AND THEN TRANSFERS LOAD FROM THE PRIMARY SUPPORT FEEDER TO A SECONDARY SUPPORT FEEDER NOT ALLOWING VIOLATIONS. IF VIOLATIONS STILL EXIST THESE ARE REMOVED BY THE THIRD STEP (THE LOAD CURTAILMENT)
AFTER THE LOAD CURTAILMENT THERE IS A POSSIBILITY TO RESTORE LOAD AMONG FORMER VIOLATION FEEDERS AND THIS IS DONE IN THE FOURTH STEP.

SEE ALSO

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
OUTFI	INFI	READS THE NETWORK TOPOLOGY AND PHYSICAL DATA FROM DATA FILE
OUTFI	WRITEF	SAVES THE NETWORK TOPOLOGY AND PHYSICAL DATA ON A DATA FILE
CHG	INC	CHANGES TOTAL LOAD
CHG	CHANGE	CHANGES PHYSICAL DATA IN A SINGLE ELEMENT
TESTS	NETTES	TESTS THE TOPOLOGY OF THE NETWORK
TESTS	INIT	INITIATES THE NETWORK
COMP	FILPOW	COMPUTES POWER
COMP	COMPII	COMPUTES IMPEDANCE


```

C   COMP      FILVOL      COMPUTES AND FILLS IN NODE VOLTAGES AND CABLE CURRENTS
C   FAULT     FLTID       READS THE FAULTED SECTION
C   RESTORE   RESULT      PRODUCES THE RESULT MENUE
C   RESTORE   GARBAG      'GARBAGE COLLECTOR' CLEANS UP AMONG SWITCHING
C                               OPERATIONS

```

```

C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER FAULT(2)
INTEGER COU,PA
CHARACTER*7 CH2
LOGICAL VIO,READY,WR
INTEGER TIME,MSCPU,I,A
CHARACTER DATA

```

```

EXTERNAL MSCPU

```

```

20  WRITE(*,*)'          WHICH FILE IS TO BE READ?'
    READ(*,300)CH2
    CALL INFI(CH2)
    CALL NETTES(CH2)
10  WRITE(*,*)'          DO YOU WISH TO HAVE A VISIBLE TRACK? Y/N'
    READ(*,300)DATA
    IF (DATA.EQ.'Y') THEN
        WR=.TRUE.
    ELSEIF (DATA.EQ.'N') THEN
        WR=.FALSE.
    ELSE
        GOTO 10
    ENDIF
    READY=.FALSE.
40  IF (.NOT. READY) THEN
    WRITE(*,100)'DO YOU WISH TO MAKE ANY CHANGES IN THE NETWORK?Y/N'
    READ(*,300)DATA
    IF (DATA.EQ.'Y' .OR. DATA.EQ.'y') THEN
        WRITE(*,100)'CHANGE TOTAL LOAD? Y/N'
        READ(*,300)DATA
        IF (DATA.EQ.'Y' .OR. DATA.EQ.'y') THEN
            CALL INC
        ELSE
            CALL CHANGE
        ENDIF
        WRITE(*,100)'ON WHICH FILE DO YOU WISH TO STORE IT'
        READ(*,300) CH2
        CALL WRITEF(CH2)
    ELSEIF (DATA.NE.'N') THEN
        GOTO 40
    ELSE
        READY=.TRUE.
        GOTO 40
    ENDIF
ENDIF
CALL INIT(WR)
DO 50 COU=1,MAXNRT
    CALL FILPOW(COU)
50  CONTINUE
DO 60 COU=1,MAXNRF
    CALL FILVOL(COU,VIO)
60  CONTINUE

C   THE NETWORK IS NOW INITIATED
C   IDENTIFY FAULT AND START CLOCK

```

```

CALL FLTID(FAULT)
TIME=MSCPU()

C   RUN ALGORITHM

CALL STEP1(FAULT)
CALL STEP2(VIO)
IF (VIO) THEN
  CALL STEP3
ENDIF

C   STOP CLOCK

TIME=MSCPU()-TIME

C   CLEAN UP AMONG SWITCHING OPERATIONS AND PRESENT RESULT

CALL GARBAG
CALL RESULT(TIME)
WRITE(*,100)'ANOTHER FAULT? Y/N'
READ(*,300) DATA
READY=DATA.EQ.'N' .OR. DATA.EQ.'n'
IF (.NOT. READY) THEN
  GOTO 20
ENDIF
100  FORMAT('D',TR6,A)
200  FORMAT(' ',T10,A)
300  FORMAT(' ',T1,A)
STOP
END

SUBROUTINE RESULT(TIME)
C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C   AUTHOUR MAGNUS BERGSTRAND
C   VERSION 1
C   DATE 1989-10-07
C   PURPOSE
C   PRESENTS THE RESULT
C   VARIABLES
C   TIME IS EXECUTION TIME IN ms
C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER TIME,I
REAL SQRT

INTRINSIC SQRT

WRITE(*,100)'TOTAL COMPUTATION TIME OF ALGORITM',TIME,' ms'
WRITE(*,*)
WRITE(*,200)'THE FOLLOWING SWITCHES HAS CHANGED STATUS'
WRITE(*,*)
DO 10 I=1,MAXNRS
  IF (CHGSWI(I)) THEN
    WRITE(*,300) I,SWINR(I)

```

```

        ENDIF
10    CONTINUE
    WRITE(*,*)
    WRITE(*,200)'TOTAL DEENERGIZED LOAD DUE TO FAULT'
    WRITE(*,500) SQRT(ISOLAT(1)**2+ISOLAT(2)**2)
    WRITE(*,*)
    WRITE(*,200)'TOTAL RESTORED LOAD'
    WRITE(*,500) SQRT(RESTOR(1)**2+RESTOR(2)**2)
    WRITE(*,*)
100   FORMAT(' ',TR6,A,IS,A)
200   FORMAT(' ',TR6,A)
300   FORMAT(' ',TR15,IS,TR3,I2)
500   FORMAT(' ',TR6,E12.5)
    RETURN
    END

```

SUBROUTINE GARBAG

```

C     CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C     AUTHOUR MAGNUS BERGSTRAND
C
C     VERSION 1
C
C     DATE 1989-10-07
C
C     PURPOSE
C     DECREASE THE NUMBER OF NECESSARY SWITCHING OPERATIONS
C
C     DESCRIPTION
C     WHEN FOR INSTANCE TWO SECTIONS HAS BEEN DEENERGIZED IN STEP3
C     TWO SWITCHES HAVE CHANGED STATUS. THE SAME RESULT COULD MAYBE
C     BEEN GIVEN IF ONLY THE SWITCH NEAREST THE FEEDER CHANGED STATUS
C     ALL SWITCHES ARE EXAMINED
C     WHEN AN OPEN SWITCH IS FOUND THE SUBROUTINE 'MOVES' TOWARDS
C     THE FEEDER.
C     IF AN OPEN SWITCH IS FOUND ON THE WAY IT IS POSSIBLE TO REDUCE THE NUMBER
C     OF SWITCHING OPERATIONS WITH 1.
C     THE SWITCHING OPERATION PERFORMED ON THE STARTING SWITCH WAS UNNECESSARY
C
C     SEE ALSO
C     FILES   ROUTINES   COMMENTS
C     STEP2  STEP3      THE LOAD CURTAILMENT
C     MOVING STEPUP     TAKES ONE STEP UP
C
C     CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C     INCLUDE 'CONSTANTS.FOR'
C     INCLUDE 'COMMON.FOR'
C
C     INTEGER P(2)
C     INTEGER I
C     LOGICAL READY
C
C     DO 10 I=1,MAXNRS
C       IF (SWITCH(I,1).EQ.OPEN .AND. CHGSWI(I)) THEN
C         P(1)=SWITYP
C         P(2)=I
C         READY=.FALSE.
20      IF (.NOT. READY) THEN
C          CALL STEPUP(P,.FALSE.)
C          READY=(P(1).EQ.SWITYP .OR. P(1).EQ.FEETYP)
C          READY=(READY .OR. P(1).EQ.TRATYP)
C          GOTO 20
        ENDIF

```

```
      IF (P(1).EQ.SWITYP) THEN
        IF (SWITCH(P(2),1).EQ.OPEN) THEN
          CHGSWI(I)=.NOT. CHGSWI(I)
          SWITCH(I,1)=CLOSED
        ENDIF
      ENDIF
    ENDIF
10  CONTINUE
    RETURN
  END
```

C CCC

C CONSTANTS.FOR

C CCC

C IN THIS FILE ALL CONSTANTS USED ARE DECLARED

C READ MAXNR_x AS MAX NUMBER OF x

C BRANCH IS THE MAXIMUM NUMBER OF BRANCHES IN ANY BRANCH POINT IN THE SYSTEM

C

PARAMETER (MAXNRT=4,MAXNRF=6,MAXNRN=78,MAXNRB=18)

PARAMETER (MAXNRC=177,BRANCH=4,MAXNRS=64)

PARAMETER (TRATYP=1,FEETYP=2,CABTYP=3)

PARAMETER (BRATYP=4,NODTYP=5,SWITYP=6)

PARAMETER (OPEN=0,CLOSED=1,UP=1,DOWN=0)

PARAMETER (NET=10)

PARAMETER (BETA=0.00001)

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      COMMON.FOR
C
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-09-30
C
C      READ ALSO 'SERVICE RESTORATION IN ELECTRICAL DISTRIBUTION NETWORKS'
C      BY MAGNUS BERGSTRAND (APPENDIX SECTION STRUCTURE CHART)
C
C      DESCRIPTION:
C      IN THIS FILE ALL COMMON VARIABLES IS DECLARED.
C      BELOW FOLLOWS A SHORT DESCRIPTION OF EACH VARIABLE.
C
C      ALFA IS THE NODE PRIORITY
C
C      IN 'SWINR' THE ORDER OF THE SWITCHING OPERATIONS IS STORED
C
C      IN ISOLAT AND RESTORE ARE THE AMOUNT OF ISOLATED LOAD AND RESTORED LOAD
C      STORED (W,VAR)
C
C      'TRAFOS' CONTAINS ALL INFORMATION ABOUT THE DIFFERENT TRANSFORMERS IN THE
C      NETWORK EXCEPT THE CABLE TO THE NETWORK WHICH IS STORED IN 'TRANEI'
C
C      MSFEED CONTAINS ALL DATA FOR THE FEEDERS
C
C      CABTOP CONTAINS TOPOLOGY FOR CABLES
C
C      CABPHY=PHYSICAL DATA FOR CABLES. CURRENTS IN AMPERE AND IMPEDANCES
C      IN OHM.
C
C      BRAPOI CONTAINS THE TOPOLOGY FOR BRANCH POINTS
C
C      NODETO CONTAINS THE TOPOLOGY FOR THE NODES
C
C      NODEPH CONTAINS THE PHYSICAL DATA FOR NODES
C      VOLTAGE IN VOLT,IMPEDANCE IN OHM AND POWER IN W
C
C      SWITCH CONTAINS ALL USEFUL INFORMATION ABOUT THE SWITCHES
C
C      IN BRAPH THE IMPEDANCE SEEN LOOKING DOWN A BRANCH IS STORED (OHM)
C
C      A CHART OF WHERE THE DATA IS STORED IN THE FOLLOWING MATRICES:
C      TRAFOS,MSFEED,CABTOP,CABPHY,BRAPOI,NODETO,NODEPH,SWITCH IS FOUND
C      IN THE REPORT 'SERVICE RESTORATION IN ELECTRIC DISTRIBUTION NETWORKS'
C
C      TRAMAX IS THE MAXIMUM VOLTAGE FOR ANY TRANSFORMER IN THE NETWORK
C
C      VIOVEC IS THE VIOLATION VECTOR.
C      THE FIRST INDEX INDICATES THE (VIOLATED) FEEDER
C
C      THE LOGICAL VECTOR CHGSWI INDICATES IF A SWITCH HAS CHANGED STATUS
C      COMPARED WITH THE INITIAL STATE
C
C      IMPLICIT LOGICAL (A-Z)
C      REAL ALFA(MAXNRN)
C      INTEGER SWINR(MAXNRS)
C      REAL ISOLAT(2),RESTOR(2)
C      REAL TRAFOS(1:MAXNRT,1:7),TRAMAX
C      INTEGER TRANEI(1:MAXNRT)

```

```
INTEGER MSFEED(1:MAXNRF,1:3)
INTEGER CABTOP(1:MAXNRC,1:4)
REAL CABPHY(1:MAXNRC,1:4)
INTEGER BRAPOI(1:MAXNRB,1:BRANCH+3)
INTEGER NODETO(1:MAXNRN,1:4)
REAL NODEPH(1:MAXNRN,1:5),VOLLIM
INTEGER SWITCH(1:MAXNRS,1:5)
REAL BRAPH(1:MAXNRB,3:BRANCH+2,1:2)
REAL VIOVEC(MAXNRF,2)
LOGICAL CHGSWI(MAXNRS)
COMMON TRAFOS,CABPHY,NODEPH
COMMON TRANEI,MSFEED,CARTOP
COMMON BRAPOI,NODETO,SWITCH
COMMON TRAMAX,VOLLIM,RESTOR
COMMON BRAPH,VIOVEC,ISOLAT
COMMON CHGSWI,SWINR,ALFA
```

C CCC

C CHG.FOR

C CCC

SUBROUTINE CHANGE

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
C THIS PROCEDURE HANDLES THE CHANGE OF PHYSICAL
C DATA IN A SINGLE ELEMENT

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS ARE DECLARED
COMMON		ALL COMMON VARIABLES ARE DECLARED
CHG	INC	CHANGES TOTAL LOAD

C VARIABLES:

C VOLTAGE IN VOLT
C CURRENT IN AMPERE
C IMPEDANCE IN OHM

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

EXTERNAL WRITEF

INTEGER P(2),I

WRITE(*,100)
WRITE(*,*) ' CHANGE IN (TYPE , NR) ?'
READ(*,*) P(1),P(2)

C THE POINT WHICH WILL BE CHANGED WAS READ

```

10 IF (P(1).EQ. TRATYP) THEN
    WRITE(*,*) ' NEW PMAX IS'
    READ(*,*) TRAFOS(P(2),1)
    WRITE(*,*) ' NEW QMAX IS'
    READ(*,*) TRAFOS(P(2),2)
ELSEIF (P(1).EQ. CABTYP) THEN
    WRITE(*,*) ' NEW R IS'
    READ(*,*) CABPHY(P(2),1)
    WRITE(*,*) ' NEW X IS'
    READ(*,*) CABPHY(P(2),2)
    WRITE(*,*) ' NEW IMAX IS'
    READ(*,*) CABPHY(P(2),3)
ELSEIF (P(1).EQ. NODTYP) THEN
    WRITE(*,*) ' R IS',NODEPH(P(2),4)

```



```

WRITE(*,*)'      X IS',NODEPH(P(2),5)
WRITE(*,*)'      NEW R IS'
READ(*,*) NODEPH(P(2),4)
WRITE(*,*)'      NEW X IS'
READ(*,*) NODEPH(P(2),5)
ELSE
  GOTO 10
ENDIF
100 FORMAT('1')
RETURN
END

```

SUBROUTINE INC

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
C READS A FACTOR P WHICH ALL THE NODE IMPEDANCES
C IN THE NETWORK WILL BE MULTIPLIED WITH CAUSING
C A CHANGE IN THE NETWORK LOAD

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS ARE DECLARED
COMMON		ALL COMMON VARIABLES ARE DECLARED
CHG	CHANGE	CHANGES PHYSICAL DATA OF ANY COMPONENT IN THE NETWORK

C VARIABLES:

C VOLTAGE IN VOLT
C CURRENT IN AMPERE
C IMPEDANCE IN OHM

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

REAL P
INTEGER I,J

```

WRITE(*,*)'      IMPEDANCES WILL BE MULTIPLIED WITH?'
READ(*,*) P
DO 10 I=1,MAXNRN
  NODEPH(I,4)=NODEPH(I,4)*P
  NODEPH(I,5)=NODEPH(I,5)*P
10 CONTINUE
RETURN
END

```

C CCC

C FAULT.FOR

C CCC

SUBROUTINE FLTID(FAULT)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE
C READS THE FAULTED POINT

C SEE ALSO

C FILES ROUTINES COMMENTS
C RESTORE MAIN MAIN PROGRAM

C VARIABLES
C FAULT (OUT) IS THE FAULTED POINT (TYPE,NR)

C CCC

INTEGER FAULT(1:2)

```
10 WRITE(*,*) '          FAULT IN'  
WRITE(*,*)  
WRITE(*,*) '          1: TRANSFORMER'  
WRITE(*,*) '          2: FEEDER (MAIN SUPPORT)'  
WRITE(*,*) '          3: CABLE'  
WRITE(*,*) '          4: BRANCHINGPOINT'  
WRITE(*,*) '          5: NODE'  
WRITE(*,*)  
WRITE(*,*) '          WHICH TYPE OF COMPONENT?'  
READ(*,*) FAULT(1)  
IF ((FAULT(1).LE.0).OR.(FAULT(1).GE.6)) GOTO 10  
  WRITE(*,*)  
  WRITE(*,*)  
  IF (FAULT(1).EQ.1) THEN  
    WRITE(*,*) '          WHICH TRANSFORMER?'  
  ELSEIF (FAULT(1).EQ.2) THEN  
    WRITE(*,*) '          WHICH FEEDER?'  
  ELSEIF (FAULT(1).EQ.3) THEN  
    WRITE(*,*) '          WHICH CABLE?'  
  ELSEIF (FAULT(1).EQ.4) THEN  
    WRITE(*,*) '          WHICH BRANCHINGPOINT?'  
  ELSEIF (FAULT(1).EQ.5) THEN  
    WRITE(*,*) '          WHICH NODE?'  
  ENDIF  
READ(*,*) FAULT(2)  
RETURN  
END
```

C CCC

C OUTFI.FOR

C CCC

SUBROUTINE WRITEF(N)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-08

C PURPOSE

C SAVES THE NETWORK TOPOLOGY AND PHYSICAL DATA ON A FILE

C DESCRIPTION

C USES FORMATTED OUTPUT

C READ ALSO USER'S GUIDE IN 'SERVICE RESRORATION IN ELECTRICAL
C DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND WHERE EXPLANATIONS AND
C EXAMPLES ARE FOUND.

C SEE ALSO

C FILES ROUTINES COMMENTS
C OUTFI INFI READS A FILE CONTAINING THE NETWORK TOPOLOGY
C AND PHYSICAL DATA FOR A NETWORK USING THE SAME
C FORMAT

C VARIABLES

C THE STRING N CONTAINS THE NAME OF THE FILE

C CCC

INCLUDE 'CONSTANTS.FOR'

INCLUDE 'COMMON.FOR'

INTEGER I,J
CHARACTER*7 N

OPEN(1,FILE=N,STATUS='NEW',ACCESS='SEQUENTIAL',FORM='FORMATTED')

WRITE(1,600) MAXNRT

WRITE(1,200) TRAMAX

DO 10 I=1,MAXNRT

 WRITE(1,300) TRAFOS(I,1),TRAFOS(I,2),TRANEI(I)

10 CONTINUE

WRITE(1,600) MAXNRF

DO 20 I=1,MAXNRF

 WRITE(1,400)MSFEED(I,2),MSFEED(I,3)

20 CONTINUE

WRITE(1,600) MAXNRC

DO 40 I=1,MAXNRC

 DO 50 J=1,3

 WRITE(1,200) CABPHY(I,J)

50 CONTINUE

 DO 60 J=1,4

 WRITE(1,600) CABTOP(I,J)

60 CONTINUE

40 CONTINUE

WRITE(1,600) MAXNRB

DO 70 I=1,MAXNRB

 WRITE(1,600) BRAP01(I,1)

```

        DO 80 J=3,(BRAPOI(I,1)+2)
            WRITE(1,600) BRAPOI(I,J)
80      CONTINUE
70      CONTINUE
WRITE(1,600) MAXNRN
WRITE(1,200) VOLLIM
DO 90 I=1,MAXNRN
    WRITE(1,400) NODETO(I,2),NODETO(I,3)
    WRITE(1,500) NODEPH(I,4),NODEPH(I,5)
    WRITE(1,200) ALFA(I)
90      CONTINUE
WRITE(1,600) MAXNRS
DO 100 I=1,MAXNRS
    WRITE(1,400) SWITCH(I,3),SWITCH(I,4)
100     CONTINUE
DO 110 I=1,MAXNRS
    WRITE(1,600) SWITCH(I,4)
110     CONTINUE
CLOSE(1,STATUS='KEEP')
200    FORMAT(' ',TR2,E12.5)
300    FORMAT(' ',TR2,2(E12.5,TR2),I6)
400    FORMAT(' ',2(TR2,I6))
500    FORMAT(' ',2(TR2,E12.5))
600    FORMAT(' ',TR2,I6)
RETURN
END

```

```

SUBROUTINE INFI(N)
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-08
C
C      PURPOSE
C      READS A FILE CONTAINING THE NETWORK TOPOLOGY AND PHYSICAL DATA
C      FOR A NETWORK.
C
C      DESCRIPTION
C      FORMATTED INPUT IS USED.
C      READ ALSO USER'S GUIDE IN 'SERVICE RESBORATION IN ELECTRICAL
C      DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND WHERE EXPLANATIONS AND
C      EXAMPLES ARE FOUND.
C
C      SEE ALSO
C      FILES      ROUTINES      COMMENTS
C      OUTFI      WRITEF      SAVES A NETWORK ON A FILE USING THE SAME FORMAT
C
C      VARIABLES
C      THE STRING N IS THE NAME OF THE FILE
C
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      INCLUDE 'CONSTANTS.FOR'
C      INCLUDE 'COMMON.FOR'
C
C      INTEGER I,J,K
C      CHARACTER*7 N
C
C      OPEN(1,FILE=N,STATUS='OLD',ACCESS='SEQUENTIAL',FORM='FORMATTED')
C      REWIND 1

```

```

READ(1,600) I
IF (I.EQ.MAXNRT) THEN
  READ(1,200) TRAMAX
  DO 10 J=1,I
    READ(1,300) TRAFOS(J,1),TRAFOS(J,2),TRANEI(J)
10  CONTINUE
ELSE
  WRITE(*,*)'      WRONG NR OF TRAFOS',I,MAXNRT
  STOP
ENDIF
READ(1,*) I
IF (I.EQ.MAXNRF) THEN
  DO 20 J=1,I
    READ(1,400) MSFEED(J,2),MSFEED(J,3)
20  CONTINUE
ELSE
  WRITE(*,*)'      WRONG NR OF FEEDERS',I,MAXNRF
  STOP
ENDIF
READ(1,600) I
IF (I.EQ.MAXNRC) THEN
  DO 30 J=1,I
    DO 40 K=1,3
      READ(1,200) CABPHY(J,K)
40  CONTINUE
    DO 50 K=1,4
      READ(1,600) CABTOP(J,K)
50  CONTINUE
30  CONTINUE
ELSE
  WRITE(*,*)'      WRONG NR OF CABLES',I,MAXNRC
  STOP
ENDIF
READ(1,600) I
IF (I.EQ.MAXNRB) THEN
  DO 60 J=1,I
    READ(1,600) BRAPOI(J,1)
    DO 70 K=3,(2+BRAPOI(J,1))
      READ(1,600) BRAPOI(J,K)
70  CONTINUE
60  CONTINUE
ELSE
  WRITE(*,*)'      WRONG NR BRANCHPOINTS',I,MAXNRB
  STOP
ENDIF
READ(1,600) I
IF (I.EQ.MAXNRN) THEN
  READ(1,200) VOLLIM
  DO 80 J=1,I
    READ(1,400) NODETO(J,2),NODETO(J,3)
    READ(1,500) NODEPH(J,4),NODEPH(J,5)
    READ(1,200) ALFA(J)
80  CONTINUE
ELSE
  WRITE(*,*)'      WRONG NR OF NODES',I,MAXNRN
  STOP
ENDIF
READ(1,600) I
IF (I.EQ.MAXNRS) THEN
  DO 90 J=1,I
    READ(1,400) SWITCH(J,3),SWITCH(J,4)
90  CONTINUE
    DO 100 J=1,I
      READ(1,600) SWITCH(J,1)
100 CONTINUE
ELSE

```

```

        WRITE(*,*)'WRONG NR OF SWITCHES',I,MAXNRS
        STOP
    ENDIF
200  FORMAT(' ',TR2,E12.5)
300  FORMAT(' ',TR2,2(E12.5,TR2),I6)
400  FORMAT(' ',2(TR2,I6))
500  FORMAT(' ',2(TR2,E12.5))
600  FORMAT(' ',TR2,I6)
    CLOSE(1,STATUS='KEEP')
    RETURN
END

```

SUBROUTINE LINES(NR)

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-08
C
C      PURPOSE
C      PRODUCES BLANK ROWS USED IN LAYOUT
C
C      VARIABLES
C      NR IS THE NUMBER OF BLANK ROWS
C
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      INTEGER NR,I
C      DO 10 I=1,NR
C          WRITE(*,*)
10  CONTINUE
C      RETURN
C      END

```

SUBROUTINE TYPTAB

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-08
C
C      PURPOSE
C      PRODUCES A MENUE ON THE SCREEN
C
C      SEE ALSO
C      FILES      ROUTINES      COMMENTS
C      TESTS      NOTEST        THIS ROUTINE FOR INSTANCE CALLS TYPTAB
C
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      WRITE(*,*)
C      WRITE(*,*) '          TYPES IN THE NETWORK'
C      WRITE(*,*)
C      WRITE(*,*) '          TRAF0          :1'
C      WRITE(*,*) '          FEEDER         :2'
C      WRITE(*,*) '          CABLE           :3'

```

```
WRITE(*,*) '          BRANCHPOINT :4'  
WRITE(*,*) '          NODE :5'  
WRITE(*,*) '          SWITCH :6'  
WRITE(*,*) '  
WRITE(*,*) '          CABLE CONNECT TO ...?'  
RETURN  
END
```

C CCC

C FETCH.FOR

C CCC

SUBROUTINE GETPOI(SWI,P)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE

C FETCHES THE POINT (TYPE,NR) ON THE OTHER SIDE OF AN OPEN SWITCH

C SEE ALSO

C	FILES	ROUTINES	COMMENTS
C	FETCH	GETVOL	GETS THE VOLTAGE ON THE OTHER SIDE OF AN OPEN SWITCH
C	FETCH	GETFEE	GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN SWITCH
C	FETCH	REALFE	GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN SWITCH AND TESTS THAT IT IS AN UNBROKEN CONNECTION (REAL POWER)
C	FETCH	GETPLI	COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED (REACTIVE POWER)
C	FETCH	GETQLI	COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED (REACTIVE POWER)
C	CONSTANTS		ALL CONSTANTS DECLARED
C	COMMON		ALL COMMON VARIABLES DECLARED

C VARIABLES

C SWI IS THE NUMBER OF THE OPEN SWITCH
C P (OUT) IS THE POINT (TYPE,NR)

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER SWI,P(2)

```

IF (SWITCH(SWI,2).EQ.3) THEN
  I=SWITCH(SWI,4)
ELSE
  I=SWITCH(SWI,3)
ENDIF
IF (CABTOP(I,1).EQ.SWITYP .AND. CABTOP(I,2).EQ.SWI) THEN
  P(1)=CABTOP(I,3)
  P(2)=CABTOP(I,4)
ELSE
  P(1)=CABTOP(I,1)
  P(2)=CABTOP(I,2)
ENDIF
RETURN
END

```

REAL FUNCTION GETVOL(SWI)

C CCC

C AUTHOUR MAGNUS BERGSTRAND


```

C   AUTHOUR MAGNUS BERGSTRAND
C
C   VERSION 1
C
C   DATE 1989-10-07
C
C   PURPOSE
C   GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN SWITCH
C
C   SEE ALSO
C   FILES      ROUTINES  COMMENTS
C   FETCH     GETPOI    GETS THE POINT ON THE OTHER SIDE OF AN OPEN SWITCH
C   FETCH     GETVOL    GETS THE VOLTAGE ON THE OTHER SIDE OF AN OPEN SWITCH
C   FETCH     REALFE    GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN
C                       SWITCH AND TESTS THAT IT IS AN UNBROKEN CONNECTION
C   FETCH     GETPLI    COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED
C                       (REAL POWER)
C   FETCH     GETQLI    COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED
C                       (REACTIVE POWER)
C   MOVING    STEPUP    TAKES ONE STEP UP
C   CONSTANTS ALL CONSTANTS DECLARED
C   COMMON    ALL COMMON VARIABLES DECLARED
C
C   VARIABLES
C   SWI IS THE NUMBER OF THE OPEN SWITCH
C
C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C   INCLUDE 'CONSTANTS.FOR'
C   INCLUDE 'COMMON.FOR'
C
C   INTEGER SWI,P(2)
C
C   IF (SWITCH(SWI,1).EQ.CLOSED) THEN
C     WRITE(*,*)'      WRONG IN GETFFE'
C     STOP
C   ELSE
10  CALL GETPOI(SWI,P)
C     IF (P(1).EQ.NODTYP) THEN
C       GETFEE=NODETO(P(2),4)
C     ELSEIF (P(1).EQ.BRATYP) THEN
C       CALL STEPUP(P,.FALSE.)
C       GOTO 10
C     ELSEIF (P(1).EQ.FEETYP) THEN
C       GETFEE=P(2)
C     ELSEIF (P(1).EQ.SWITYP) THEN
C       CALL STEPUP(P,.FALSE.)
C       GOTO 10
C     ELSE
C       WRITE(*,*)'      DEGENERATED NET OR WRONG USE OF GETFEE'
C       STOP
C     ENDIF
C   ENDIF
C   RETURN
C   END
C
C   SUBROUTINE REALFE(SWI,FEE,FOUND)
C
C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C   AUTHOUR MAGNUS BERGSTRAND
C
C   VERSION 1

```

C DATE 1989-10-07

C PURPOSE

C GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN SWITCH
C AND TESTS IF IT IS AN UNBROKEN CONNECTION

C DESCRIPTION

C FIRST A CALL OF GETFEE IS MADE
C THEN STARTING FROM THE POINT ON THE OTHER SIDE OF THE SWITCH
C IT MOVES UPWARDS UNTIL A FEEDER OR AN OPEN SWITCH IS FOUND.
C IF A FEEDER IS FOUND THEN 'FOUND'=TRUE

C SEE ALSO

C	FILES	ROUTINES	COMMENTS
C	CONSTANTS		ALL CONSTANTS DECLARED
C	COMMON		ALL COMMON VARIABLES DECLARED
C	FETCH	GETPOI	GETS THE POINT ON THE OTHER SIDE OF AN OPEN SWITCH
C	FETCH	GETVOL	GETS THE VOLTAGE ON THE OTHER SIDE OF AN OPEN SWITCH
C	FETCH	GETFEE	GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN SWITCH
C	FETCH	GETPLI	COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED (REAL POWER)
C	FETCH	GETQLI	COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED (REACTIVE POWER)
C	MOVING	STEPUP	TAKES ONE STEP UP

C VARIABLES

C SWI IS THE NUMBER OF THE OPEN SWITCH
C FEE (OUT) IS THE NUMBER OF THE FEEDER
C FOUND (OUT) INDICATES WHETHER IT WAS AN UNBROKEN CONNECTION OR NOT

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER SWI,FEE
INTEGER P(2)
LOGICAL FOUND,READY
INTEGER GETFEE

EXTERNAL GETFEE

```
FEE=GETFEE(SWI)
CALL GETPOI(SWI,P)
READY=.FALSE.
10 IF (.NOT. READY) THEN
    CALL STEPUP(P,.FALSE.)
    IF (P(1).EQ. TRATYP) THEN
        WRITE(*,*)'      DEGENERATED NET.WRONG IN REALFE'
        STOP
    ELSEIF (P(1).EQ. FEETYP) THEN
        READY=.TRUE.
    ELSEIF (P(1).EQ. SWITYP) THEN
        IF (SWITCH(P(2),1).EQ.OPEN) THEN
            READY=.TRUE.
        ENDIF
    ENDIF
    GOTO 10
ENDIF
FOUND=(FEE.EQ.P(2) .AND. P(1).EQ.FEETYP)
RETURN
END
```

REAL FUNCTION GETPLI(SWI)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE
C GETS THE MARGIN OF A TRANSFORMER GETTING VIOLATED (REAL POWER)

C SEE ALSO
C FILES ROUTINES COMMENTS
C CONSTANTS ALL CONSTANTS DECLARED
C COMMON ALL COMMON VARIABLES DECLARED
C FETCH GETPO1 GETS THE POINT ON THE OTHER SIDE OF AN OPEN SWITCH
C FETCH GETVOL GETS THE VOLTAGE ON THE OTHER SIDE OF AN OPEN SWITCH
C FETCH GETFEE GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN
C SWITCH
C FETCH REALFE GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN
C SWITCH AND TESTS THAT IT IS AN UNBROKEN CONNECTION
C FETCH GETQLI COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED
C (REACTIVE POWER)

C VARIABLES
C THE MARGIN IN W.
C SWI IS THE NUMBER OF THE OPEN SWITCH

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER SWI,P(2)
INTEGER FEE,TRA

IF (SWITCH(SWI,1).EQ.CLOSED) THEN
 WRITE(*,*)' MISUSED GETPLI'
 STOP
ELSE
 FEE=GETFEE(SWI)
 TRA=MSFEED(FEE,1)
 GETPLI=TRAFOS(TRA,1)-TRAFOS(TRA,3)
ENDIF
RETURN
END

REAL FUNCTION GETQLI(SWI)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE
C GETS THE MARGIN OF GETTING VIOLATED FOR A TRANSFORMER ON THE OTHER SIDE
C OF AN OPEN SWITCH (REACTIVE POWER)

C SEE ALSO
C FILES ROUTINES COMMENTS
C CONSTANTS ALL CONSTANTS DECLARED

```

C      COMMON          ALL COMMON VARIABLES DECLARED
C      FETCH          GETPOI      GETS THE POINT ON THE OTHER SIDE OF AN OPEN SWITCH
C      FETCH          GETPOI      GETS THE VOLTAGE ON THE OTHER SIDE OF AN OPEN SWITCH
C      FETCH          GETPOI      GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN
C                                  SWITCH
C      FETCH          REALFE      GETS THE FEEDER IDENTITY ON THE OTHER SIDE OF AN OPEN
C                                  SWITCH AND TESTS THAT IT IS AN UNBROKEN CONNECTION
C      FETCH          GETPLI      COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED
C                                  (REAL POWER)
C      FETCH          GETQLI      COMPUTES THE MARGIN OF A TRANSFORMER GETTING VIOLATED
C                                  (REACTIVE POWER)

```

```

C      VARIABLES
C      THE MARGIN IN VA.
C      SWI IS THE NUMBER OF THE OPEN SWITCH

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER SWI,P(2)
INTEGER FEE,TRA

```

```

IF (SWITCH(SWI,1).EQ.CLOSED) THEN
  WRITE(*,*)'      MISUSED GETQLI'
  STOP
ELSE
  FEE=GETFEE(SWI)
  TRA=MSFEED(FEE,1)
  GETQLI=TRAFOS(TRA,2)-TRAFOS(TRA,4)
ENDIF
RETURN
END

```

C CCC

C STEP1.FOR

C CCC

SUBROUTINE STEP1(FAULT)

C CCC

AUTHOUR MAGNUS BERGSTRAND

VERSION 1

DATE 1989-10-07

PURPOSE

ISOLATES A FAULTED ZONE AND CONNECTS DEENERGIZED NON FAULTED ZONES TO ADJACENT FEEDERS.

METHOD

THE METHOD IS DESCRIBED IN 'SERVICE RESTORATION IN ELECTRIC DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND THE METHOD WAS FIRST DESCRIBED BY AOKI IN 'A NEW ALGORITHM FOR SERVICE RESTORATION IN DISTRIBUTION SYSTEMS' IEEE/PES WINTER MEETING, NEW YORK 1989 PAPER 89 WM 085-2 PWRD

DESCRIPTION

FAULT IS THE FAULTED POINT (TYPE, NR) THERE ARE TWO BASIC TYPES OF FAULTS. FAULTS AMONG THE ZONES OR FAULTS "ABOVE THE FEEDER" (IN A TRANSFORMER) IF A NODE IS FAULTED THIS IS ISOLATED BY FIRST FINDING THE FIRST SWITCH ABOVE THE NODE. CLOSE IT AND THEN CLOSING ALL FIRST SWITCHES BELOW THE FAULTED NODE. IT CAN BE MORE THAN ONE SWITCH DUE TO BRANCH POINTS. WHEN A TRANSFORMER IS FAULTED THE FAULT IS ISOLATED BY THE FAULT ITSELF. ALL THE CIRCUIT BREAKERS ARE TRIPED BY THE FAULT. AFTER THE FAULT HAS BEEN ISOLATED THE ROOTS OF THE ISOLATED TREES ARE WRITTEN ON A FILE AND THEN READ BACK. STARTING FROM THESE ROOTS ALL POSSIBLE CONNECTIONS ARE EXAMINED AND THE ONE WITH THE LEAST CHANCE OF CAUSING ANY VIOLATION IS SELECTED REENERGIZING THE TREE.

SEE ALSO

FILES ROUTINES COMMENTS STEP2 STEP2 SECOND STEP IN THE ALGORITHM STEP2 STEP3 THIRD AND FOURTH STEP IN THE ALGORITHM STEP1 ABOVEF LOGICAL FUNCTION. IS THE FAULT LOCATED ABOVE ANY FEEDER? MOVING MOVE TAKES ONE STEP UP OR DOWN MOVING DIREC COMPUTES NEW DIRECTIONS MOVING STEPUP TAKES ONE STEP UP MOVING STEPDO TAKES ONE STEP DOWN COMP CON CONNECTS DEENERGIZED ZONES CONSTANTS ALL CONSTANTS COMMON ALL COMMON VARIABELS ARE DECLARED

VARIABLES

FAULT IS THE FAULTED POINT (TYPE, NR)

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```
INTEGER FAULT(1:2)
INTEGER TRAF0,FEEDER,ISOSWI
INTEGER CURPOI(2),P(2),P1(2),I
INTEGER DIR,BRA,OLDP(2),Y,J
LOGICAL READY,WR,CHG,ABOVEF,OK
```

```
PARAMETER(TRASH=1)
```

```
EXTERNAL STEPUP,STEPDO
EXTERNAL MOVE,DIREC
EXTERNAL CON,ABOVEF
```

```
OPEN(TRASH,STATUS='SCRATCH',ACCESS='SEQUENTIAL')
```

```
CURPOI(1)=FAULT(1)
```

```
CURPOI(2)=FAULT(2)
```

```
WR=.FALSE.
```

```
FEEDER=0
```

```
ISOSWI=0
```

```
ISOLAT(1)=0.0
```

```
ISOLAT(2)=0.0
```

```
IF (.NOT. ABOVEF(CURPOI)) THEN
```

```
C      FAULT AMONG SECTIONS (BELOW FEEDER)
```

```
      READY=CURPOI(1).EQ.TRATYP
```

```
10     IF (.NOT. READY) THEN
```

```
C      FIND FIRST SWITCH
```

```
      IF (CURPOI(1).EQ.CABTYP) THEN
```

```
        CALL STEPDO(CURPOI,BRA,WR)
```

```
      ENDIF
```

```
      CALL STEPUP(CURPOI,WR)
```

```
      READY=(CURPOI(1).EQ.SWITYP .OR. CURPOI(1).EQ.TRATYP)
```

```
      GOTO 10
```

```
    ENDIF
```

```
    DIR=DOWN
```

```
    BRA=3
```

```
    P(1)=CURPOI(1)
```

```
    P(2)=CURPOI(2)
```

```
    READY=.FALSE.
```

```
    CHG=.FALSE.
```

```
20     IF (.NOT. READY) THEN
```

```
C      FIND ALL SWITCHES BELOW THE FAULTED ZONE
```

```
C      THESE SWITCHES ARE OPENED
```

```
C      THEY MUST BE OPENED IN THIS PERHAPS ODD WAY BECAUSE OTHERWISE WILL
```

```
C      NEITHER STEPUP NOR STEPDO WORK CORRECTLY
```

```
      IF (CHG) THEN
```

```
        SWITCH(Y,1)=CLOSED
```

```
      ENDIF
```

```
      CALL MOVE(P,OLDP,DIR,BRA,WR)
```

```
      IF (CHG) THEN
```

```
        SWITCH(Y,1)=OPEN
```

```
      ENDIF
```

```
      READY=(P(2).EQ.CURPOI(2) .AND. P(1).EQ.CURPOI(1))
```

```
      IF (READY) THEN
```

```
        GOTO 20
```

```
      ENDIF
```

```
      CHG=.FALSE.
```

```
      IF (P(1).EQ.SWITYP) THEN
```

```
        I=P(2)
```

```
        IF (SWITCH(I,1).EQ.CLOSED) THEN
```

```
          SWITCH(I,1)=OPEN
```

```
          SWITCH(I,5)=0
```

```

SWINR(I)=1
CHGSWI(I)=.NOT. CHGSWI(I)
CHG=.TRUE.
Y=I
IF (SWITCH(I,2).EQ.3) THEN
    J=SWITCH(I,4)
ELSE
    J=SWITCH(I,3)
ENDIF

```

```

C     SAVE THE POINT BELOW THE SWITCH THAT HAS ISOLATED THE FAULT
C     THIS POINT IS THEN USED AS ROOT IN A TREE IN ORDER TO DECIDE
C     THE BEST WAY TO REENERGIZE THE NON FAULTED ZONES

```

```

IF (CABTOP(J,1).EQ.SWITYP .AND. CABTOP(J,2).EQ.1) THEN
    WRITE(TRASH,*) CABTOP(J,3),CABTOP(J,4)
ELSEIF (CABTOP(J,3).EQ.SWITYP .AND. CABTOP(J,4).EQ.1) THEN
    WRITE(TRASH,*) CABTOP(J,1),CABTOP(J,2)
ELSE
    WRITE(*,*)'          WRONG IN STEP 1'
    STOP
ENDIF
ELSEIF (SWITCH(I,1).EQ.OPEN) THEN
    SWITCH(I,5)=0
ELSE
    WRITE(*,*)'          WRONG IN STEP1',SWITCH(I,1),OPEN,CLOSED
    STOP
ENDIF
ENDIF
CALL DIREC(P,OLDP,DIR,BRA)
GOTO 20

```

```

ENDIF
IF (CURPOI(1).EQ.SWITYP) THEN
    SWITCH(CURPOI(2),1)=OPEN
    SWITCH(CURPOI(2),5)=0
    SWINR(CURPOI(2))=1
    CHGSWI(CURPOI(2))=.NOT. CHGSWI(CURPOI(2))
ENDIF

```

```

ENDFILE TRASH
REWIND TRASH
ELSE

```

```

C     THE FAULT WAS IN A TRANSFORMER (ABOVE THE FEEDER)
C     START FROM THE TRANSFORMER MOVING UPWARDS
C     OPENING ALL SWITCHES ON THE WAY UNTIL YOU FIND A FEEDER. STOP.
C     WRITE THE NUMBER OF THE FEEDER ON A FILE
C     DO THIS FOR ALL POSSIBLE BRANCHES

```

```

OPEN(3,STATUS='SCRATCH',ACCESS='SEQUENTIAL')
READY=CURPOI(1).EQ.TRATYP
100 IF (.NOT. READY) THEN
    IF (CURPOI(1).EQ.CABTYP) THEN
        CALL STEPDO(CURPOI,BRA,WR)
    ENDIF
    CALL STEPUP(CURPOI,WR)
    READY=(CURPOI(1).EQ.SWITYP .OR. CURPOI(1).EQ.TRATYP)
    GOTO 100

```

```

ENDIF
IF (CURPOI(1).EQ.SWITYP) THEN
    WRITE(3,*)CURPOI(2)
ENDIF
DIR=DOWN
BRA=3
P(1)=CURPOI(1)
P(2)=CURPOI(2)
READY=.FALSE.

```



```

200 IF (.NOT. READY) THEN
    CALL MOVE(P,OLDP,DIR,BRA,.FALSE.)
    READY=(P(2).EQ.CURPOI(2) .AND. P(1).EQ.CURPOI(1))
    IF (READY) THEN
        GOTO 200
    ENDIF
    IF (P(1).EQ.FEETYP) THEN
        TRAF0=MSPEED(P(2),1)

C      ISOLATED LOAD

        ISOLAT(1)=TRAF0S(TRAFO,3)
        ISOLAT(2)=TRAF0S(TRAFO,4)
        DIR=UP
        P1(1)=FEETYP
        P1(2)=P(2)
        CALL STEPDO(P1,3,.FALSE.)
        IF (P1(1).EQ.SWITYP) THEN
            I=P1(2)
            IF (SWITCH(I,1).EQ.CLOSED) THEN
                SWITCH(I,1)=OPEN
                SWITCH(I,5)=0
                SWINR(I)=1
                CHGSWI(I)=.NOT. CHGSWI(I)
                IF (SWITCH(I,2).EQ.3) THEN
                    J=SWITCH(I,4)
                ELSE
                    J=SWITCH(I,3)
                ENDIF
                IF (CABTOP(J,1).EQ.SWITYP .AND. CABTOP(J,2).EQ.1) THEN
                    WRITE(TRASH,*) CABTOP(J,3),CABTOP(J,4)
                ELSEIF (CABTOP(J,3).EQ.SWITYP .AND. CABTOP(J,4).EQ.1) THEN
                    WRITE(TRASH,*) CABTOP(J,1),CABTOP(J,2)
                ELSE
                    WRITE(*,*)'      WRONG IN STEP 1'
                    STOP
                ENDIF
                ELSEIF (SWITCH(I,1).EQ.OPEN) THEN
                    SWITCH(I,5)=0
                ENDIF
            ENDIF
        ELSEIF (P(1).EQ.SWITYP .AND. DIR.EQ.DOWN) THEN
            IF (SWITCH(P(2),1).EQ.CLOSED) THEN
                WRITE(3,*)P(2)
            ENDIF
        ENDIF
        CALL DIREC(P,OLDP,DIR,BRA)
        GOTO 200
    ENDIF
    IF (CURPOI(1).EQ.SWITYP) THEN
        SWITCH(CURPOI(2),1)=OPEN
        SWITCH(CURPOI(2),5)=0
        SWINR(CURPOI(2))=1
        CHGSWI(CURPOI(2))=.NOT. CHGSWI(CURPOI(2))
    ENDIF
    ENDFILE TRASH
    REWIND TRASH
    ENDFILE 3
    REWIND 3
205 READ(3,*,ERR=210) I

C      READ THE NUMBER OF THE SWITCHES THAT WILL BE OPENED TO ISOLATE THE
C      FAULT

        SWITCH(I,1)=OPEN
        SWITCH(I,5)=0

```

```

        SWINR(I)=1
        CHGSWI(I)=.NOT. CHGSWI(I)
        GOTO 205
210  ENDIF
30   READ(TRASH,*,END=40) CURPOI(1),CURPOI(2)
C    READS AND CONNECTS THE DEENERGIZED NODES

        CALL CON(CURPOI,OK)
        IF (.NOT. OK) THEN
            RESTOR(1)=0.0
            RESTOR(2)=0.0
        ENDIF
        GOTO 30
40   CLOSE(TRASH,STATUS='DELETE')
        IF (OK) THEN
            RESTOR(1)=ISOLAT(1)
            RESTOR(2)=ISOLAT(2)
        ENDIF
        RETURN
    END

```

LOGICAL FUNCTION ABOVEF(POI)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE
 C INDICATES IF A FAULT HAS OCCURRED ABOVE A FEEDER

C DESCRIPTION
 C MOVES UPWARDS UNTIL A FEEDER OR A TRANSFORMER IS FOUND.
 C IF A TRANSFORMER IS FOUND THEN TRUE

C SEE ALSO

C FILES	ROUTINES	COMMENTS
C STEP1	STEP1	ISOLATES THE FAULT AND CONNECTS THE NON FAULTED ZONES.
C		CALLS ABOVEF
C MOVING	STEPUP	TAKES ONE STEP UP
C CONSTANTS		ALL CONSTANTS DECLARED
C COMMON		ALL COMMON VARIABLES DECLARED

C VARIABLES
 C POI IS THE FAULTED POINT (TYPE,NR)

C CCC

INCLUDE 'CONSTANTS.FOR'
 INCLUDE 'COMMON.FOR'

INTEGER POI(2),P(2)
 LOGICAL READY

P(1)=POI(1)
 P(2)=POI(2)
 READY=(P(1).EQ.FEETYP .OR. P(1).EQ.TRATYP)
 IF (.NOT. READY) THEN
 10 CALL STEPUP(P,.FALSE.)

```
    READY=(P(1).EQ.FEETYP .OR. P(1).EQ.TRATYP)
    GOTO 10
ENDIF
ABOVEF=(P(1).EQ.TRATYP)
RETURN
END
```

C CCC

C MOVING.FOR

C CCC

SUBROUTINE STEPUP(POINT,WR)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE

C TAKES ONE STEP UP

C DESCRIPTION

C A CENTRAL SUBROUTINE.

C NOTICE THAT POINT CAN BE OF ANY TYPE EXCEPT CABLE.

C IN MSFEED,NODETO,BRAPOI AND SWITCH A POINTER UP IS FOUND.

C IT POINTS OUT AN INDEX TO THE CABLE LEADING TOWARDS THE TRANSFORMER

C IN CABTOP (CABLE TOPOLOGY) THE TWO END POINTS OF THE CABLE ARE STORED .

C THIS SUBROUTINE MOVES 'POINT' TO THE OTHER END POINT IN DIRECTION UP

C SEE ALSO

FILES	ROUTINES	COMMENTS
MOVING	STEPDO	TAKES ONE STEP DOWN
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
CONSTANT		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED

C VARIABLES

C POINT IS THE CURRENT POINT (TYPE,NR)

C IF WR=TRUE THE CURRENT POINT WILL BE PRINTED ON THE SCREEN

C CCC

INCLUDE 'CONSTANTS.FOR'

INCLUDE 'COMMON.FOR'

INTEGER POINT(1:2),I

LOGICAL WR

IF (WR) THEN

WRITE(*,*)' POINT IN IS',POINT(1),POINT(2)

ENDIF

IF (POINT(1) .EQ. BRATYP) THEN

I=BRAPOI(POINT(2),2)

I=BRAPOI(POINT(2),1)

IF (CABTOP(I,1).EQ.POINT(1) .AND. CABTOP(I,2).EQ.POINT(2)) THEN

POINT(1)=CABTOP(I,3)

POINT(2)=CABTOP(I,4)

ELSE

POINT(1)=CABTOP(I,1)

POINT(2)=CABTOP(I,2)

ENDIF

ELSEIF (POINT(1).EQ.NODTYP) THEN

I=NODETO(POINT(2),1)

```

IF (WR) THEN
  WRITE(*,*)'      CABUP IN NODE IS',I
ENDIF
I=NODETO(POINT(2),I)
IF (WR) THEN
  WRITE(*,*)'      CABLE UP IS',I
ENDIF
IF (CABTOP(I,1).EQ.POINT(1) .AND. CABTOP(I,2).EQ.POINT(2)) THEN
  POINT(1)=CABTOP(I,3)
  POINT(2)=CABTOP(I,4)
ELSE
  POINT(1)=CABTOP(I,1)
  POINT(2)=CABTOP(I,2)
ENDIF
ELSEIF (POINT(1).EQ.SWITYP) THEN
  I=SWITCH(POINT(2),2)
  I=SWITCH(POINT(2),I)
  IF (CABTOP(I,1).EQ.POINT(1) .AND. CABTOP(I,2).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I,3)
    POINT(2)=CABTOP(I,4)
  ELSE
    POINT(1)=CABTOP(I,1)
    POINT(2)=CABTOP(I,2)
  ENDIF
ELSEIF (POINT(1).EQ.FEETYP) THEN
  I=MSFEED(POINT(2),2)
  IF (CABTOP(I,1).EQ.FEETYP .AND. CABTOP(I,2).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I,3)
    POINT(2)=CABTOP(I,4)
  ELSE
    POINT(1)=CABTOP(I,1)
    POINT(2)=CABTOP(I,2)
  ENDIF
ELSE
  WRITE(*,*)'      WRONG TYPE IN (SUBROUTINE STEPUP)'
  STOP
ENDIF
IF (WR) THEN
  WRITE(*,*)'      POINT OUT FROM STEPUP IS',POINT(1),POINT(2)
ENDIF
RETURN
END

```

SUBROUTINE STEPDO(POINT,BRANMR,WR)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE

C TAKES ONE STEP DOWN

C DESCRIPTION

C THIS PROCEDURE TAKES A STEP IN THE DIRECTION "NOT DOWN"

C IT WORKS FOR ALL TYPES INCLUDING CABLES.

C WHEN POINT IS A BRANCH POINT STEPDO TAKES A STEP DOWN IN THE

C BRANCH POINTED OUT BY 'BRANKR'

C NOTICE THAT NO CHECK IS MADE THAT THIS BRANCH LEADS DOWNWARDS

```

C     SEE ALSO
C     FILES   ROUTINES   COMMENTS
C     MOVING  STEPUP    TAKES ONE STEP UP
C     MOVING  MOVE      TAKES ONE STEP UP OR DOWN
C     MOVING  DIREC     COMPUTES NEW DIRECTION
C     CONSTANT ALL CONSTANTS DECLARED
C     COMMON  ALL COMMON VARIABLES DECLARED

C     VARIABLES
C     POINT IS THE CURRENT POINT (TYPE, NR)
C     BRANMR IS THE NUMBER OF THE DOWNWARDS LEADING BRANCH THAT WILL BE VISITED
C     IF WR=TRUE THEN THE CURRENT POINT WILL BE WRITTEN ON THE SCREEN

C     CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER BRANMR, POINT(2)
LOGICAL WR
INTEGER I, J, K, L
IF (WR) THEN
  WRITE(*,*)'      POINT IN IS ', POINT(1), POINT(2)
ENDIF
IF (POINT(1).EQ.FEETYP) THEN
  I=MSFEED(POINT(2), 3)
  IF (CABTOP(I, 1).EQ.FEETYP .AND. CABTOP(I, 2).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 3)
    POINT(2)=CABTOP(I, 4)
  ELSEIF (CABTOP(I, 3).EQ.FEETYP .AND. CABTOP(I, 4).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 1)
    POINT(2)=CABTOP(I, 2)
  ELSE
    WRITE(*,*)'      WRONG IN STEPDO.FEEDER'
    STOP
  ENDIF
ELSEIF (POINT(1).EQ.BRATYP) THEN
  I=BRAPOI(POINT(2), BRANMR)
  IF (CABTOP(I, 1).EQ.BRATYP .AND. CABTOP(I, 2).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 3)
    POINT(2)=CABTOP(I, 4)
  ELSEIF (CABTOP(I, 3).EQ.BRATYP .AND. CABTOP(I, 4).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 1)
    POINT(2)=CABTOP(I, 2)
  ELSE
    WRITE(*,*)'      WRONG IN STEPDO.BRA'
    STOP
  ENDIF
ELSEIF (POINT(1).EQ.TRATYP) THEN
  I=TRANEI(POINT(2))
  IF (CABTOP(I, 1).EQ.TRATYP .AND. CABTOP(I, 2).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 3)
    POINT(2)=CABTOP(I, 4)
  ELSEIF (CABTOP(I, 3).EQ.TRATYP .AND. CABTOP(I, 4).EQ.POINT(2)) THEN
    POINT(1)=CABTOP(I, 1)
    POINT(2)=CABTOP(I, 2)
  ELSE
    WRITE(*,*)'      WRONG IN STEPDO.TRAFO'
    WRITE(*,*)'TRAFO', POINT(2)
    WRITE(*,*)CABTOP(I, 1), CABTOP(I, 2), CABTOP(I, 3), CABTOP(I, 4)
    STOP
  ENDIF
ELSEIF (POINT(1).EQ.CABTYP) THEN

```

C THE USED METHOD FOR CABLES

C LOOK AT ONE NEIGHBOUR
C IF CABLE UP IS IDENTICAL WITH THE CURRENT CABLE GO TO THAT NEIGHBOUR
C OTHERWISE GO TO THE OTHER NEIGHBOUR

```
K=POINT(2)
IF (CABTOP(K,1).EQ.SWITYP .OR. CABTOP(K,3).EQ.SWITYP) THEN
  IF (CABTOP(POINT(2),1).EQ.SWITYP) THEN
    I=1
  ELSE
    I=3
  ENDIF
  J=MOD((I+2),4)
  IF (CABTOP(POINT(2),J).EQ.TRATYP) THEN
    POINT(1)=CABTOP(POINT(2),I)
    POINT(2)=CABTOP(POINT(2),I+1)
  ELSEIF (CABTOP(POINT(2),J).EQ.FEETYP) THEN
    K=CABTOP(POINT(2),J+1)
    IF (POINT(2).EQ.MSFEED(K,2)) THEN
      POINT(1)=FEETYP
      POINT(2)=CABTOP(POINT(2),J+1)
    ELSE
      POINT(1)=SWITYP
      POINT(2)=CABTOP(POINT(2),I+1)
    ENDIF
  ELSEIF (CABTOP(POINT(2),J).EQ.BRATYP) THEN
    K=CABTOP(POINT(2),J+1)
    L=BRAPOI(K,2)
    IF (POINT(2).EQ.BRAPOI(K,L)) THEN
      POINT(1)=BRATYP
      POINT(2)=CABTOP(POINT(2),J+1)
    ELSE
      POINT(1)=SWITYP
      POINT(2)=CABTOP(POINT(2),I+1)
    ENDIF
  ELSEIF (CABTOP(POINT(2),J).EQ.NODTYP) THEN
    K=CABTOP(POINT(2),J+1)
    L=NODETO(K,1)
    IF (POINT(2).EQ.NODETO(K,L)) THEN
      POINT(1)=NODTYP
      POINT(2)=CABTOP(POINT(2),J+1)
    ELSE
      POINT(1)=SWITYP
      POINT(2)=CABTOP(POINT(2),I+1)
    ENDIF
  ENDIF
  ENDIF
ELSEIF (CABTOP(POINT(2),1).EQ.TRATYP) THEN
  POINT(1)=CABTOP(POINT(2),3)
  POINT(2)=CABTOP(POINT(2),4)
ELSEIF (CABTOP(POINT(2),1).EQ.FEETYP) THEN
  K=CABTOP(POINT(2),2)
  IF (POINT(2).EQ.MSFEED(K,2)) THEN
    POINT(1)=FEETYP
    POINT(2)=CABTOP(POINT(2),2)
  ELSE
    POINT(1)=CABTOP(POINT(2),3)
    POINT(2)=CABTOP(POINT(2),4)
  ENDIF
ELSEIF (CABTOP(POINT(2),1).EQ.CABTYP) THEN
  WRITE(*,*) '          DEGENERATED NETWORK'
  STOP
ELSEIF (CABTOP(POINT(2),1).EQ.BRATYP) THEN
  K=CABTOP(POINT(2),2)
  L=BRAPOI(K,2)
  IF (POINT(2).EQ.BRAPOI(K,L)) THEN
    POINT(1)=BRATYP
    POINT(2)=CABTOP(POINT(2),2)
```

```

ELSE
  POINT(1)=CABTOP(POINT(2),3)
  POINT(2)=CABTOP(POINT(2),4)
ENDIF
ELSEIF (CABTOP(POINT(2),1).EQ.NODTYP) THEN
  K=CABTOP(POINT(2),2)
  L=NODETO(K,1)
  IF (POINT(2).EQ.NODETO(K,L)) THEN
    POINT(1)=NODTYP
    POINT(2)=CABTOP(POINT(2),2)
  ELSE
    POINT(1)=CABTOP(POINT(2),3)
    POINT(2)=CABTOP(POINT(2),4)
  ENDIF
ENDIF
ENDIF

```

C HERE THE SECTION CONCERNING CABLES ENDS

```

ELSEIF (POINT(1).EQ.NODTYP) THEN
  IF (NODETO(POINT(2),1).EQ.2) THEN
    I=3
  ELSEIF (NODETO(POINT(2),1).EQ.3) THEN
    I=2
  ELSE
    WRITE(*,*)'      WRONG IN STEPDO.NODE'
    STOP
  ENDIF
  J=NODETO(POINT(2),1)
  IF (CABTOP(J,1).EQ.NODTYP .AND. POINT(2).EQ.CABTOP(J,2)) THEN
    POINT(1)=CABTOP(J,3)
    POINT(2)=CABTOP(J,4)
  ELSEIF (CABTOP(J,3).EQ.NODTYP .AND. POINT(2).EQ.CABTOP(J,4)) THEN
    POINT(1)=CABTOP(J,1)
    POINT(2)=CABTOP(J,2)
  ELSE
    WRITE(*,*)'      WRONG IN STEPDO.NODE'
    STOP
  ENDIF
ELSEIF (POINT(1).EQ.SWITYP) THEN
  IF (SWITCH(POINT(2),1).EQ.CLOSED) THEN
    IF (SWITCH(POINT(2),2).EQ.3) THEN
      I=4
    ELSE
      I=3
    ENDIF
    J=SWITCH(POINT(2),1)
    IF (CABTOP(J,1).EQ.SWITYP .AND. CABTOP(J,2).EQ.POINT(2)) THEN
      POINT(1)=CABTOP(J,3)
      POINT(2)=CABTOP(J,4)
    ELSEIF (CABTOP(J,3).EQ.6 .AND. CABTOP(J,4).EQ.POINT(2)) THEN
      POINT(1)=CABTOP(J,1)
      POINT(2)=CABTOP(J,2)
    ELSE
      WRITE(*,*)'      WRONG IN STEPDO.SWITCH'
      WRITE(*,*) SWITYP,POINT(2)
      WRITE(*,*) CABTOP(J,1),CABTOP(J,2),CABTOP(J,3),CABTOP(J,4)
      STOP
    ENDIF
  ENDIF
ENDIF
ENDIF
IF (WR) THEN
  WRITE(*,*)'      POINT OUT FROM STEP DOWN IS',POINT(1),POINT(2)
ENDIF
RETURN
END

```



```

SUBROUTINE MOVE(PPOINT,OLD,DIR,BRA,WR)

C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C AUTHOUR MAGNUS BERGSTRAND
C
C VERSION 1
C
C DATE 1989-10-07
C
C PURPOSE
C TAKES ONE STEP UP OR DOWN. USED IN COOPERATION WITH NEWDIR IT CAN BE
C USED TO TRAVERSE A TREE.
C
C DESCRIPTION
C DIR INDICATES THE DIRECTION (UP,DOWN)
C IF DIRECTION IS DOWN AND PPOINT IS BRANCH POINT THE VARIABLE BRA IS VALID.
C IT INDICATES THE BRANCH THAT WILL BE 'VISITED'
C THIS SUBROUTINE WORKS IF THE NETWORK TOPOLOGY IS ALREADY INITIATED.
C NOTICE THAT THIS PROCEDURE ASSIGNS A NEW VALUE TO 'CABLE UP' IF A
C CLOSED SWITCH IS FOUND. UP WILL BE ASSIGNED THE DIRECTION TOWARDS OLD.
C
C SEE ALSO
C FILES  ROUTINES  COMMENTS
C MOVING  STEPUP   TAKES ONE STEP UP
C MOVING  STEPDO   TAKES ONE STEP DOWN
C MOVING  DIREC    COMPUTES NEW DIRECTION
C CONSTANT  ALL CONSTANTS DECLARED
C COMMON    ALL COMMON VARIABLES DECLARED
C TESTS    TKESTP  TAKES ONE STEP UP OR DOWN. WORKS IN AN UNINITIATED
C NETWORK.
C
C VARIABLES
C PPOINT IS THE CURRENT POINT (TYPE,NR)
C OLD IS THE FORMER POINT (TYPE,NR)
C DIR INDICATES THE DIRECTION (UP,DOWN)
C BRA IS THE NUMBER OF THE BRANCH
C IF WR=TRUE THEN THE CURRENT POINT WILL BE WRITTEN ON THE SCREEN
C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C INCLUDE 'CONSTANTS.FOR'
C INCLUDE 'COMMON.FOR'
C
C INTEGER PPOINT(2),OLD(2)
C INTEGER DIR,BRA
C INTEGER TYP,NR,P(2)
C LOGICAL WR
C
C TYP=PPOINT(1)
C NR=PPOINT(2)
C OLD(1)=PPOINT(1)
C OLD(2)=PPOINT(2)
C IF (DIR.EQ.DOWN) THEN
C   IF (TYP.EQ.SWITYP) THEN
C     IF (SWITCH(NR,1).EQ.CLOSED) THEN
C
C       IT IS ONLY POSSIBLE TO TAKE ONE STEP DOWN IF THE SWITCH IS CLOSED
C
C       CALL STEPDO(PPOINT,BRA,WR)
C       IF (PPOINT(1).EQ.SWITYP) THEN
C         IF (SWITCH(PPOINT(2),1).EQ.OPEN) THEN
C
C           THE RESULT FROM THIS STEP WAS AN OPEN SWITCH

```

```

C      IN THIS SWITCH ASSIGN UP A NEW VALUE TOWARDS OLD.

      I=SWITCH(POINT(2),3)
      IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
        SWITCH(POINT(2),2)=3
      ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
        SWITCH(POINT(2),2)=3
      ELSE
        SWITCH(POINT(2),2)=4
      ENDIF
    ENDIF
  ENDIF
ENDIF
ELSE
  CALL STEPDO(POINT,BRA,WR)
  IF (POINT(1).EQ.SWITYP) THEN
    IF(SWITCH(POINT(2),1).EQ.OPEN) THEN

```

C THE RESULT FROM THIS STEP WAS AN OPEN SWITCH
C IN THIS SWITCH ASSIGN UP A NEW VALUE TOWARDS OLD.

```

      I=SWITCH(POINT(2),3)
      IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
        SWITCH(POINT(2),2)=3
      ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
        SWITCH(POINT(2),2)=3
      ELSE
        SWITCH(POINT(2),2)=4
      ENDIF
    ENDIF
  ENDIF
ENDIF
ELSEIF (DIR.EQ.UP) THEN
  CALL STEPUP(POINT,WR)
ENDIF
RETURN
END

```

SUBROUTINE DIREC(POINT,OLDP,DIR,BRA)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-07

C PURPOSE
C COMPUTES A NEW DIRECTION

C DESCRIPTION
C IS INTENDED TO COOPERATE WITH MOVE. THEN IT CAN BE USED TO TRAVERSE
C A TREE WITH EITHER AN OPEN SWITCH OR A 'LEAF NODE' (ONLY ONE FEEDING
C CABLE ATTACHED) AS LEAFS.

C SEE ALSO

C FILES	ROUTINES	COMMENTS
C MOVING	STEPUP	TAKES ONE STEP UP
C MOVING	STEPDO	TAKES ONE STEP DOWN
C MOVING	MOVE	TAKES ONE STEP UP OR DOWN
C TESTS	NEWDIR	COMPUTES A NEW DIRECTION IN A NON INITIATED NETWORK
C CONSTANT		ALL CONSTANTS DECLARED
C COMMON		ALL COMMON VARIABLES DECLARED

```

C   VARIABLES
C   POINT IS THE CURRENT POINT (TYPE,NR)
C   OLDP IS THE FORMER POINT (TYPE,NR)
C   DIR IS THE DIRECTION
C   BRA IS THE NUMBER OF THE BRANCH
C
C   CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C   INCLUDE 'CONSTANTS.FOR'
C   INCLUDE 'COMMON.FOR'
C
C   INTEGER POINT(2),OLDP(2)
C   INTEGER DIR,BRA,NXTBRA
C   INTEGER I,C
C   LOGICAL READY
C
C   IF (DIR.EQ.DOWN) THEN
C     IF (POINT(1).EQ.NODTYP) THEN
C       IF (NODETO(POINT(2),2).EQ.0 .OR. NODETO(POINT(2),3).EQ.0) THEN
C
C         IF THE NODE IS A LEAF NODE CHANGE DIRECTIONS!
C
C         DIR=UP
C       ENDIF
C     ELSEIF (POINT(1).EQ.SWITYP) THEN
C       IF (SWITCH(POINT(2),1).EQ.OPEN) THEN
C
C         IF AN OPEN SWITCH IS FOUND CHANGE DIRECTION!
C
C         DIR=UP
C       ENDIF
C     ELSEIF (POINT(1).EQ.BRATYP) THEN
C
C       IF CABLE UP=4 THEN BRA=4 ELSE BRA=3
C
C       IF (BRAPOI(POINT(2),2).EQ.3) THEN
C         BRA=4
C       ELSE
C         BRA=3
C       ENDIF
C     ENDIF
C   ELSEIF (DIR.EQ.UP) THEN
C
C     IN THIS CASE THE ONLY TIME WHEN THE DIRECTION IS CHANGED IS WHEN
C     A BRANCH POINT IS REACHED.
C     THEN THE NEXT BRANCH VISITED WILL BE BRA+1 IF THIS BRANCH EXISTS
C     AND IF IT IS NOT THE BRANCH UP
C
C     IF (POINT(1).EQ.BRATYP) THEN
C       BRA=2
C       I=POINT(2)
C       READY=.FALSE.
C     10 IF (.NOT. READY) THEN
C       BRA=BRA+1
C       C=BRAPOI(I,BRA)
C       IF (CABTOP(C,1).EQ.OLDP(1) .AND. CABTOP(C,2).EQ.OLDP(2)) THEN
C         READY=.TRUE.
C       ELSEIF (CABTOP(C,3).EQ.OLDP(1) .AND. CABTOP(C,4).EQ.OLDP(2)) THEN
C         READY=.TRUE.
C       ENDIF
C       GOTO 10
C     ENDIF
C     NXTBRA=BRA+1
C     IF (NXTBRA.LE.(2+BRAPOI(I,1)) .AND. NXTBRA.NE.BRAPOI(I,2)) THEN
C       DIR=DOWN
C       BRA=NXTBRA

```

```
ELSEIF (NXTBRA.LT.BRAPOI(I,1)+2 .AND. (NXTBRA+1).NE.BRAPOI(I,2)) THEN  
  BRA=BRA+2  
  DIR=DOWN  
ENDIF  
ENDIF  
ENDIF  
RETURN  
END
```

C CCC

C TESTS.FOR

C CCC

SUBROUTINE CABFOR(J)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-06

C PURPOSE
C PRODUCES A MENUE

C SEE ALSO

C	FILES	ROUTINES	COMMENTS
C	CONSTANTS		ALL CONSTANTS DECLARED
C	COMMON		ALL COMMON VARIABLES DECLARED
C	TESTS	TRTEST	TEST OF TOPOLOGY (TRANSFORMERS)
C	TESTS	FETEST	TEST OF TOPOLOGY (FEEDERS)
C	TESTS	BRTEST	TEST OF TOPOLOGY (BRANCH POINTS)
C	TESTS	NOTEST	TEST OF TOPOLOGY (NODES)
C	TESTS	SWTEST	TEST OF TOPOLOGY (SWITCHES)
C	THE ABOVE CALLS CABFOR		

C CCC

C INCLUDE 'CONSTANTS.FOR'
C INCLUDE 'COMMON.FOR'

C INTEGER J

```

WRITE(*,200) 'CABLE NR',J,' IS CONNECTED TO '
WRITE(*,300) 'TYPE=',CABTOP(J,1),' NR=',CABTOP(J,2)
WRITE(*,*) ' AND '
WRITE(*,300) 'TYPE=',CABTOP(J,3),' NR=',CABTOP(J,4)
200 FORMAT(' ',TR6,A,I5,A)
300 FORMAT(' ',TR6,A,I2,A,I6)
RETURN
END

```

SUBROUTINE WRICAB(INDEX)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-06

C PURPOSE
C PRODUCES A MENUE

C SEE ALSO

C	FILES	ROUTINES	COMMENTS
---	-------	----------	----------

```

C      CONSTANTS          ALL CONSTANTS DECLARED
C      COMMON             ALL COMMON VARIABLES DECLARED
C      TESTS   TRTEST    TEST OF TOPOLOGY (TRANSFORMERS)
C      TESTS   FETEST    TEST OF TOPOLOGY (FEEDERS)
C      TESTS   BRTEST    TEST OF TOPOLOGY (BRANCH POINTS)
C      TESTS   NOTEST    TEST OF TOPOLOGY (NODES)
C      TESTS   SWTEST    TEST OF TOPOLOGY (SWITCHES)
C      THE ABOVE CALLS WRICAB

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'
INTEGER INDEX

```

```

WRITE(*,*)'          CABLE IS CONNECTED TO:      TYPE?  NR?'
READ(*,*) CABTOP(INDEX,1),CABTOP(INDEX,2)
WRITE(*,*)'          AND TO:                      TYPE?  NR?'
READ(*,*) CABTOP(INDEX,3),CABTOP(INDEX,4)
WRITE(*,*)'          R? X? IMAX?'
READ(*,*) CABPHY(J,1),CABPHY(J,2),CABPHY(J,3)
RETURN
END

```

```

SUBROUTINE TRTEST(OK)

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C      AUTHOUR MAGNUS BERGSTRAND

```

```

C      VERSION 1

```

```

C      DATE 1989-10-06

```

```

C      PURPOSE
C      TESTS THE TOPOLOGY REGARDING TRANSFORMERS

```

```

C      DESCRIPTION
C      IF A POINTS NEIGHBOURS NEIGHBOUR IS NOT THE POINT ITSELF
C      A FAULT IN THE TOPOLOGY IS DETECTED

```

```

C      SEE ALSO
C      FILES   ROUTINES   COMMENTS
C      CONSTANTS          ALL CONSTANTS DECLARED
C      COMMON             ALL COMMON VARIABLES DECLARED
C      TESTS   TRTEST    TEST OF TOPOLOGY (TRANSFORMERS)
C      TESTS   FETEST    TEST OF TOPOLOGY (FEEDERS)
C      TESTS   BRTEST    TEST OF TOPOLOGY (BRANCH POINTS)
C      TESTS   NOTEST    TEST OF TOPOLOGY (NODES)
C      TESTS   SWTEST    TEST OF TOPOLOGY (SWITCHES)
C      TESTS   NETTES    TEST OF NETWORK TOPOLOGY

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER I,J,K,TYPE
LOGICAL OK

```

```

EXTERNAL LINES

```

```

CALL LINES(3)
DO 10 I=1,MAXNRT

```

```

J=TRANEI(I)
K=I
20 IF (CABTOP(J,1).NE.TRATYP .AND. CABTOP(J,2).NE.K) THEN
    IF (CABTOP(J,3).NE.TRATYP .AND. CABTOP(J,4).NE.K) THEN
        WRITE(*,*) '          FAULT IN NETWORK'
        CALL LINES(2)
        OK=.FALSE.
        WRITE(*,100) 'TRAFO NR',I,' IS CONNECTED TO CABLE NR',J
        CALL CABFOR(J)
30 WRITE(*,*) '          WHICH ONE IS THE WRONG ONE? TYPE?'
        READ(*,*) TYPE
        IF (TYPE.EQ.TRATYP) THEN
            WRITE(*,*) '          TRAFO CONNECTED TO WHICH CABLE?'
            READ(*,*) TRANEI(I)
        ELSEIF (TYPE.EQ.CABTYP) THEN
            CALL WRICAB(J)
        ELSE
            WRITE(*,*) '          WRONG TYPE'
            GOTO 30
        ENDIF
        GOTO 20
    ENDIF
ENDIF
ENDIF
10 CONTINUE
100 FORMAT(' ',TR4,A,13,A,15)
WRITE(*,*) '          TRAFOS IS OK'
RETURN
END

```

SUBROUTINE FETEST(OK)

```

C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C AUTHOUR MAGNUS BERGSTRAND
C
C VERSION 1
C
C DATE 1989-10-06
C
C PURPOSE
C TESTS THE TOPOLOGY REGARDING FEEDERS
C
C DESCRIPTION
C IF A POINTS NEIGHBOURS NEIGHBOUR IS NOT THE POINT ITSELF
C A FAULT IN THE TOPOLOGY IS DETECTED
C
C SEE ALSO
C FILES   ROUTINES   COMMENTS
C CONSTANTS          ALL CONSTANTS DECLARED
C COMMON            ALL COMMON VARIABLES DECLARED
C TESTS   TRTEST    TEST OF TOPOLOGY (TRANSFORMERS)
C TESTS   FETEST    TEST OF TOPOLOGY (FEEDERS)
C TESTS   BRTEST    TEST OF TOPOLOGY (BRANCH POINTS)
C TESTS   NOTEST    TEST OF TOPOLOGY (NODES)
C TESTS   SWTEST    TEST OF TOPOLOGY (SWITCHES)
C TESTS   NETTES    TEST OF NETWORK TOPOLOGY
C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C INCLUDE 'CONSTANTS.FOR'
C INCLUDE 'COMMON.FOR'
C
C LOGICAL OK
C INTEGER I,J,K,TY1,TY2,NR1,NR2,TYPE

```

EXTERNAL LINES

```

DO 10 I=1,MAXNRF
  DO 20 J=2,3
    K=MSFEED(I,J)
30    IF (CABTOP(K,1).NE.FEETYP .AND. CABTOP(K,2).NE.I) THEN
      IF (CABTOP(K,3).NE.FEETYP .AND. CABTOP(K,4).NE.I) THEN
        TY1=CABTOP(K,1)
        TY2=CABTOP(K,3)
        NR1=CABTOP(K,2)
        NR2=CABTOP(K,4)
        OK=.FALSE.
        CALL LINES(3)
40        WRITE(*,*) '          FAULTED NETWORK'
        WRITE(*,200)'FEEDER NR',I,'IS CONNECTED TO CABLE NR',K
        WRITE(*,*)
        CALL CABFOR(K)
        WRITE(*,*)
        WRITE(*,*) '          FAULT IN (TYPE) ?'
        READ(*,*) TYPE
        IF (TYPE.EQ.FEETYP) THEN
          WRITE(*,*) '          TO WHICH CABLE IS THE FEEDER CONNECTED?'
          READ(*,*) MSFEED(I,J)
        ELSEIF (TYPE.EQ.CABTYP) THEN
          CALL WRICAB(K)
        ELSE
          WRITE(*,*) '          WRONG TYPE!'
          GOTO 40
        ENDIF
        GOTO 30
      ENDIF
    ENDIF
20    CONTINUE
10    CONTINUE
200  FORMAT(' ',TR6,A,I3,A,I4)
    WRITE(*,*) '          FEEDER IS OK'
    RETURN
  END

```

SUBROUTINE BRTEST(OK)

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-06
C
C      PURPOSE
C      TESTS THE TOPOLOGY REGARDING BRANCH POINTS
C
C      DESCRIPTION
C      IF A POINTS NEIGHBOURS NEIGHBOUR IS NOT THE POINT ITSELF
C      A FAULT IN THE TOPOLOGY IS DETECTED
C
C      SEE ALSO
C      FILES   ROUTINES   COMMENTS
C      CONSTANTS   ALL CONSTANTS DECLARED
C      COMMON     ALL COMMON VARIABLES DECLARED
C      TESTS     TRTEST     TEST OF TOPOLOGY (TRANSFORMERS)

```



```

C TESTS FETEST TEST OF TOPOLOGY (FEEDERS)
C TESTS BRTEST TEST OF TOPOLOGY (BRANCH POINTS)
C TESTS NOTEST TEST OF TOPOLOGY (NODES)
C TESTS SWTEST TEST OF TOPOLOGY (SWITCHES)
C TESTS NETTES TEST OF NETWORK TOPOLOGY

```

```

C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

LOGICAL OK
INTEGER I,J,K,L,TYPE

```

```

EXTERNAL LINES

```

```

DO 10 I=1,MAXNRB
  DO 20 J=3,(BRAPOI(I,1)+2)
30    K=BRAPOI(I,J)
      IF (CABTOP(K,1).NE.BRATYP .AND. CABTOP(K,2).NE.I) THEN
        IF (CABTOP(K,3).NE.BRATYP .AND. CABTOP(K,4).NE.I) THEN
          OK=.FALSE.
          CALL LINES(3)
40          WRITE(*,*)'          FAULTED NETWORK'
            WRITE(*,200)'BRANCH-POINT NR',I,' IS CONNECTED TO CABLE NR',K
            WRITE(*,*)
            CALL CABFOR(K)
            WRITE(*,*)'          FAULT IN BRANCH-POINT(4) OR IN CABLE(3) ?'
            READ(*,*) TYPE
            IF (TYPE.EQ.BRATYP) THEN
              WRITE(*,*)'          THE FOLLOWING BRANCHES'
              WRITE(*,*)
              DO 50 L=3,(2+BRAPOI(I,1))
                WRITE(*,*)'          ',BRAPOI(I,L)
50              CONTINUE
              WRITE(*,*)
              WRITE(*,*)'          THE FAULTING ONE IS (NOT) CONNECTED TO',BRAPOI
              WRITE(*,*)'          THE BRANCH IS CONNECTED TO CABLE NR ?'
              READ(*,*) BRAPOI(I,J)
            ELSEIF(TYPE.EQ.CABTYP) THEN
              CALL WRICAB(K)
            ELSE
              WRITE(*,*)'          WRONG TYPE!'
              GOTO 40
            ENDIF
          GOTO 30
        ENDIF
      ENDIF
    CONTINUE
  CONTINUE
20  CONTINUE
10  CONTINUE
200 FORMAT(' ',TR6,A,13,A,14)
WRITE(*,*)'          BRANCHPOINTS OK'
RETURN
END

```

```

SUBROUTINE NOTEST(OK)

```

```

C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C AUTHOUR MAGNUS BERGSTRAND

```

```

C VERSION 1

```

```

C      DATE 1989-10-06

C      PURPOSE
C      TESTS THE TOPOLOGY REGARDING NODES

C      DESCRIPTION
C      IF A POINTS NEIGHBOURS NEIGHBOUR IS NOT THE POINT ITSELF
C      A FAULT IN THE TOPOLOGY IS DETECTED

C      SEE ALSO
C      FILES      ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON      ALL COMMON VARIABLES DECLARED
C      TESTS      TRTEST      TEST OF TOPOLOGY (TRANSFORMERS)
C      TESTS      FETEST      TEST OF TOPOLOGY (FEEDERS)
C      TESTS      BRTEST      TEST OF TOPOLOGY (BRANCH POINTS)
C      TESTS      NOTEST      TEST OF TOPOLOGY (NODES)
C      TESTS      SWTEST      TEST OF TOPOLOGY (SWITCHES)
C      TESTS      NETTES      TEST OF NETWORK TOPOLOGY

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

      INCLUDE 'CONSTANTS.FOR'
      INCLUDE 'COMMON.FOR'

      LOGICAL OK
      INTEGER I,J,K,TYPE

      EXTERNAL LINES

      DO 10 I=1,MAXNRN
        J=NODETO(I,2)
        K=NODETO(I,3)
        IF (.NOT. (J.EQ.0)) THEN
20          IF (CABTOP(J,1).NE.NODTYP .AND. CABTOP(J,2).NE.I) THEN
              IF (CABTOP(J,3).NE.NODTYP .AND. CABTOP(J,4).NE.I) THEN
                OK=.FALSE.
                CALL LINES(3)
                WRITE(*,*)'      FAULTED NETWORK'
                WRITE(*,*)
                WRITE(*,200)'NODE NR',I,' IS CONNECTED TO CABLE NR',J
                WRITE(*,*)
                CALL CABFOR(J)
                WRITE(*,*)
                CALL CABFOR(K)
30          WRITE(*,*)
                WRITE(*,*)'      FAULT IN NODE (5) OR IN CABLE (3)'
                READ(*,*) TYPE
                IF (TYPE.EQ.NODTYP) THEN
                  WRITE(*,*)'      TO WHICH CABLE IS THE NODE CONNECTED?'
                  READ(*,*) NODETO(I,2)
                ELSEIF (TYPE.EQ.CABTYP) THEN
                  CALL WRICAB(J)
                ELSE
                  WRITE(*,*)'      WRONG TYPE!'
                  GOTO 30
                ENDIF
                GOTO 20
              ENDIF
            ENDIF
          ENDIF
        IF (.NOT. (K.EQ.0)) THEN
40          IF (CABTOP(K,1).NE.NODTYP .AND. CABTOP(K,2).NE.I) THEN
              IF (CABTOP(K,3).NE.NODTYP .AND. CABTOP(K,4).NE.I) THEN
                OK=.FALSE.
                CALL LINES(3)

```

```

WRITE(*,*)'          FAULTED NETWORK'
WRITE(*,*)
WRITE(*,200)'NODE NR',I,' IS CONNECTED TO CABLE NR',K
WRITE(*,*)
CALL CABFOR(J)
WRITE(*,*)
CALL CABFOR(K)
50  WRITE(*,*)
WRITE(*,*)'          FAULT IN NODE (5) OR IN CABLE (3) '
READ(*,*) TYPE
IF (TYPE.EQ.NODTYP) THEN
  WRITE(*,*)'          TO WHICH CABLE IS THE NODE CONNECTED?'
  READ(*,*) NODETO(I,3)
ELSEIF (TYPE.EQ.CABTYP) THEN
  CALL WRICAB(K)
ELSE
  WRITE(*,*)'          WRONG TYPE!'
  GOTO 50
ENDIF
GOTO 40
ENDIF
ENDIF
ENDIF
10  CONTINUE
100  FORMAT('1')
200  FORMAT(' ',TR6,A,14,A,14)
WRITE(*,*)'          NODES OK'
RETURN
END

```

SUBROUTINE SWICAB(A,OK)

```

C  CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C  AUTHOUR MAGNUS BERGSTRAND
C  VERSION 1
C  DATE 1989-10-06
C  PURPOSE
C  TESTS THE TOPOLOGY REGARDING 'A HALF SWITCH'
C  DESCRIPTION
C  IF A POINTS NEIGHBOURS NEIGHBOUR IS NOT THE POINT ITSELF
C  A FAULT IN THE TOPOLOGY IS DETECTED
C  SEE ALSO
C  FILES  ROUTINES  COMMENTS
C  CONSTANTS  ALL CONSTANTS DECLARED
C  COMMON  ALL COMMON VARIABLES DECLARED
C  TESTS  TRTEST  TEST OF TOPOLOGY (TRANSFORMERS)
C  TESTS  FETEST  TEST OF TOPOLOGY (FEEDERS)
C  TESTS  BRTEST  TEST OF TOPOLOGY (BRANCH POINTS)
C  TESTS  NOTEST  TEST OF TOPOLOGY (NODES)
C  TESTS  SWTEST  TEST OF TOPOLOGY (SWITCHES)
C  TESTS  NETTES  TEST OF NETWORK TOPOLOGY
C  CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C  INCLUDE 'CONSTANTS.FOR'
C  INCLUDE 'COMMON.FOR'
C  INTEGER A,I,J,TYPE

```



```

C     AUTHOUR MAGNUS BERGSTRAND
C
C     VERSION 1
C
C     DATE 1989-10-06
C
C     PURPOSE
C     CAN BE USED TO TRAVERSE A NETWORK IN COOPERATION WITH NEWDIR
C
C     DESCRIPTION
C     THIS SUBROUTINE IS ONLY TO BE CALLED FROM INIT
C     TKESTP MANAGES TO TRAVERSE AN UNINITIATED NETWORK
C     IT INITIATES THE NETWORK AT THE SAME TIME
C
C     SEE ALSO
C     FILES      ROUTINES      COMMENTS
C     CONSTANTS          ALL CONSTANTS DECLARED
C     COMMON            ALL COMMON VARIABLES DECLARED
C     MOVING  STEPUP      TAKES ONE STEP UP
C     MOVING  STEPDO      TAKES ONE STEP DOWN
C     MOVING  MOVE        TAKES ONE STEP UP OR DOWN (IN AN INITIATED NETWORK)
C     MOVING  DIREC       COMPUTES NEW DIRECTION
C     TESTS   NEWDIR      COMPUTES NEW DIRECTION. IS INTENDED TO COOPERATE
C                       WITH TKESTP
C
C     VARIABLES
C     POINT IS THE CURRENT POINT. (TYPE,NR)
C     OLDP IS THE FORMER POINT. (TYPE,NR)
C     DIR IS THE DIRECTION (UP OR DOWN)
C     BRA IS THE NUMBER OF BRANCH IN CASE OF A BRANCH POINT
C     IF WR=TRUE THEN TEXT IS WRITTEN ON THE SCREEN INDICATING THE CURRENT
C     POINT
C
C     CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C     INCLUDE 'CONSTANTS.FOR'
C     INCLUDE 'COMMON.FOR'
C
C     INTEGER POINT(2),OLDP(2),DIR,BRA
C     INTEGER IN,I,J
C     LOGICAL WR
C
C     IF (WR) THEN
C       WRITE(*,*)
C       WRITE(*,*)'          INDATA TKESTP (POINT,OLDP,DIR,BRA)'
C       WRITE(*,*)POINT(1),POINT(2),OLDP(1),OLDP(2),DIR,BRA
C     ENDIF
C     IF (DIR.EQ.UP) THEN
C       OLDP(1)=POINT(1)
C       OLDP(2)=POINT(2)
C       CALL STEPUP(POINT,WR)
C
C     ONE STEP UP WAS TAKEN
C
C     IF (POINT(1).EQ.BRATYP) THEN
C
C     FIND OUT IN WHICH BRANCH OLDP IS.
C
C     DO 20 J=3,(2+BRAPO1(POINT(2),1))
C       I=BRAPO1(POINT(2),J)
C       IF (CABTOP(I,1).EQ.OLDP(1) .AND. CABTOP(I,2).EQ.OLDP(2)) THEN
C         BRA=J
C       ELSEIF (CABTOP(I,3).EQ.OLDP(1) .AND. CABTOP(I,4).EQ.OLDP(2)) THEN
C         BRA=J
C       ENDIF

```

```

20      CONTINUE
      ENDIF
ELSEIF (DIR.EQ.DOWN) THEN
  IF (POINT(1).EQ.BRATYP) THEN
    I=POINT(2)
    IF (BRAPOI(I,2).EQ.0) THEN
C      'UP' IS NOT KNOWN
C      UP IS IN THE DIRECTION OF OLDP!

      DO 30 J=3,BRAPOI(I,1)+2
        K=BRAPOI(I,J)
        IF (CABTOP(K,1).EQ.OLDP(1) .AND. CABTOP(K,2).EQ.OLDP(2)) THEN
          BRAPOI(I,2)=J
        ELSEIF (CABTOP(K,3).EQ.OLDP(1) .AND. CABTOP(K,4).EQ.OLDP(2)) THEN
          BRAPOI(I,2)=J
        ENDIF
30      CONTINUE
      IF (BRAPOI(I,2).EQ.3) THEN
C          SINCE THE DIRECTION IS DOWN

          BRA=4
        ELSE
          BRA=3
        ENDIF
      ENDIF
    ENDIF
    OLDP(1)=POINT(1)
    OLDP(2)=POINT(2)
    CALL STEPDO(POINT,BRA,WR)
C      ONE STEP DOWN WAS TAKEN
C      HEREAFTER UP WILL ASSIGNED A VALUE IN CASE OF SWITCH OR NODE

    IF (POINT(1).EQ.SWITYP) THEN
      J=POINT(2)
      I=SWITCH(J,3)
      IF (CABTOP(I,1).EQ.OLDP(1) .AND. CABTOP(I,2).EQ.OLDP(2)) THEN
        SWITCH(J,2)=3
      ELSEIF (CABTOP(I,3).EQ.OLDP(1) .AND. CABTOP(I,4).EQ.OLDP(2)) THEN
        SWITCH(J,2)=3
      ELSE
        SWITCH(J,2)=4
      ENDIF
    ELSEIF (POINT(1).EQ.NODTYP) THEN
      I=NODETO(POINT(2),2)
      IF (CABTOP(I,1).EQ.OLDP(1) .AND. CABTOP(I,2).EQ.OLDP(2)) THEN
        NODETO(POINT(2),1)=2
      ELSEIF (CABTOP(I,3).EQ.OLDP(1) .AND. CABTOP(I,4).EQ.OLDP(2)) THEN
        NODETO(POINT(2),1)=2
      ELSE
        NODETO(POINT(2),1)=3
      ENDIF
    ENDIF
  ENDIF
ENDIF
IF (WR) THEN
  WRITE(*,*)
  WRITE(*,*)'          OUTDATA TKESTP'
  WRITE(*,*)POINT(1),POINT(2),OLDP(1),OLDP(2),DIR,BRA
ENDIF
RETURN
END

SUBROUTINE NEWDIR(POINT,DIR,BRA,WR)

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-06
C
C      PURPOSE
C      COMPUTES THE NEW DIRECTION.
C
C      DESCRIPTION
C      IS USED IN COOPERATION WITH TKESTP IN ORDER TO TRAVERSE AN
C      UNINITIALIZED NETWORK
C
C      SEE ALSO
C      FILES      ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON      ALL COMMON VARIABLES DECLARED
C      MOVING  STEPUP      TAKES ONE STEP UP
C      MOVING  STEPDO      TAKES ONE STEP DOWN
C      MOVING  MOVE        TAKES ONE STEP UP OR DOWN (IN AN INITIATED NETWORK)
C      MOVING  DIREC      COMPUTES NEW DIRECTION
C      TESTS   TKESTP      TAKES ONE STEP UP OR DOWN. IS INTENDED TO COOPERATE
C                          WITH NEWDIR
C
C      VARIABLES
C      POINT IS THE CURRENT POINT. (TYPE,NR)
C      DIR IS THE DIRECTION (UP OR DOWN)
C      BRA IS THE NUMBER OF BRANCH IN CASE OF A BRANCH POINT
C      IF WR=TRUE THEN TEXT IS WRITTEN ON THE SCREEN INDICATING THE CURRENT
C      POINT
C
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      INCLUDE 'CONSTANTS.FOR'
C      INCLUDE 'COMMON.FOR'
C
C      INTEGER POINT(2)
C      INTEGER DIR,BRA
C      INTEGER I
C      LOGICAL WR
C
C      IF (WR) THEN
C        WRITE(*,*)
C        WRITE(*,*) '      INDATA NEWDIR (POINT,DIR,BRA)'
C        WRITE(*,*) POINT(1),POINT(2),DIR,BRA
C      ENDIF
C      IF (POINT(1).EQ.BRATYP) THEN
C        IF (DIR.EQ.UP) THEN
C
C          IN THIS CASE THE ONLY POSSIBLE CHANGE OF DIRECTION IS WHENEVER
C          A BRANCH POINT IS REACHED. THEN THE NEW DIRECTION MAY BE DOWN
C          IN THE NEXT BRANCH
C
C          I=POINT(2)
C          IF (BRAPOI(I,1).GT.(BRA-2) .AND. BRAPOI(I,2).NE.(BRA+1)) THEN
C            BRA=BRA+1
C            DIR=DOWN
C          ELSEIF (BRAPOI(I,2).EQ.(BRA+1)) THEN
C            IF (BRA+2.LE.2+BRAPOI(I,1)) THEN
C              BRA=BRA+2
C              DIR=DOWN
C            ENDIF
C          ENDIF
C        ENDIF
C      ENDIF

```



```
INCLUDE 'CONSTANTS.FOR'  
INCLUDE 'COMMON.FOR'
```

```
INTEGER POINT(2),P(2)  
INTEGER DIR,BRA,TRA,FEE  
INTEGER OLD(2)  
REAL CUR  
LOGICAL FIRST
```

```
IF (DIR.EQ.DOWN) THEN  
  IF (POINT(1).EQ.FEETYP) THEN  
    FEE=POINT(2)  
    MSPEED(POINT(2),1)=TRA  
  ELSEIF (POINT(1).EQ.SWITYP) THEN  
    SWITCH(POINT(2),5)=1  
  ELSEIF (POINT(1).EQ.NODTYP) THEN  
    I=POINT(2)  
    NODETO(I,4)=FEE  
    CUR=(NODEPH(I,1)**2+NODEPH(I,2)**2)/(NODEPH(I,4)**2+NODEPH(I,5)**2)  
    FIRST=.TRUE.  
    P(1)=POINT(1)  
    P(2)=POINT(2)  
    OLD(1)=0  
    OLD(2)=0  
10    IF (P(1).NE.TRATYP) THEN  
      CALL UPDATE(P,OLD,FIRST,CUR,NODEPH(POINT(2),4),NODEPH(POINT(2),5))  
      OLD(1)=P(1)  
      OLD(2)=P(2)  
      CALL STEPUP(P,.FALSE.)  
      GOTO 10  
    ENDIF  
    CALL UPDATE(P,OLD,FIRST,CUR,1.0,1.0)  
  ENDIF  
ENDIF  
RETURN  
END
```

```
SUBROUTINE UPDATE(POINT,OLD,FIRST,CUR,NODER,NODEX)
```

```
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

```
C VERSION 1
```

```
C DATE 1989-10-06
```

```
C PURPOSE  
C COMPUTES IMPEDANCES
```

```
C DESCRIPTION  
C COMPUTES IMPEDANCES WHILE MOVING FROM A NODE TO A TRANSFORMER  
C COMPUTATIONS ARE BASED ON THE FACT THAT THE NETWORK HAS A RADIAL  
C STRUCTURE.  
C THE RESULT  $Z=(RES,j*REA)$  ARE SAVED BETWEEN CALLS  
C THE PARALLELL IMPEDANCE BETWEEN THE OLD VALUE AND A NEW NODE IS COMPUTED  
C THEN IS THE CABLE IMPEDANCE ADDED
```

```
C SEE ALSO  
C FILES  ROUTINES  COMMENTS  
C CONSTANTS  ALL CONSTANTS DECLARED  
C COMMON  ALL COMMON VARIABLES DECLARED  
C MOVING  STEPUP  TAKES ONE STEP UP  
C COMP  COMPJM  COMPUTES IMPEDANCE. INITIATED NETWORK
```

```

C      COMP      BRAPAR      COMPUTES EQUIVALENT PARALLELL IMPEDANCE FOR A BRANCH
C      POINT
C      COMP      PARALL      COMPUTES EQUIVALENT PARALLELL IMPEDANCE

```

```

C      VARIABLES
C      POINT IS THE CURRENT POINT (TYPE,NR)
C      OLD IS THE FORMER POINT (TYPE,NR)
C      FIRST INDICATES FIRST CALL.
C      CUR IS THE CURRENT (AMPERE)
C      NODER IS NODE RESISTANCE (OHM)
C      NODEX IS NODE REACTANCE (OHM)

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER POINT(2),OLD(2)
INTEGER I,J,K
LOGICAL FIRST
REAL NODER,NODEX,CUR,F(2),G(2)
REAL RES,REA,SUMR,SUMX

```

```

SAVE RES,REA

```

```

IF (FIRST) THEN
  RES=NODEPH(POINT(2),4)
  REA=NODEPH(POINT(2),5)
  FIRST=.FALSE.
ELSEIF (.NOT. FIRST) THEN
  IF (POINT(1).EQ. TRATYP) THEN
    TRAFOS(POINT(2),6)=RES
    TRAFOS(POINT(2),7)=REA
  ELSEIF (POINT(1).EQ.FEETYP) THEN
    I=MSFEED(POINT(2),3)
    RES=RES+CABPHY(I,1)
    REA=REA+CABPHY(I,2)
  ELSEIF (POINT(1).EQ.BRATYP) THEN
    I=POINT(2)
    DO 30 J=3,(BRAPOI(I,1)+2)
      K=BRAPOI(I,J)
      IF (CABTOP(K,1).EQ.OLD(1) .AND. CABTOP(K,2).EQ.OLD(2)) THEN
        BRAPH(I,J,1)=RES
        BRAPH(I,J,2)=REA
      ELSEIF (CABTOP(K,3).EQ.OLD(1) .AND. CABTOP(K,4).EQ.OLD(2)) THEN
        BRAPH(I,J,1)=RES
        BRAPH(I,J,2)=REA
      ENDIF
    CONTINUE
    CALL BRAPAR(POINT(2),RES,REA)
    K=BRAPOI(POINT(2),2)
    K=BRAPOI(POINT(2),K)
    RES=CABPHY(K,1)+RES
    REA=CABPHY(K,2)+REA
  ELSEIF (POINT(1).EQ.NODTYP) THEN
    F(1)=RES
    F(2)=REA
    G(1)=NODEPH(POINT(2),4)
    G(2)=NODEPH(POINT(2),5)
    CALL PARALL(F,G,RES,REA)
    I=NODETO(POINT(2),1)
    I=NODETO(POINT(2),I)
    RES=CABPHY(I,1)+RES
    REA=CABPHY(I,1)+REA
  ELSEIF (POINT(1).EQ.SWITYP) THEN
    I=SWITCH(POINT(2),2)

```

```

        I=SWITCH(POINT(2),I)
        RES=RES+CABPHY(I,1)
        REA=REA+CABPHY(I,2)
    ELSE
        WRITE(*,*)'        WRONG TYPE IN UPDATE!'
        WRITE(*,*)'        THE POINT IS',POINT(1),POINT(2)
        STOP
    ENDIF
ENDIF
RETURN
END

```

```

SUBROUTINE INIT(WR)

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C      AUTHOUR MAGNUS BERGSTRAND

```

```

C      VERSION 1

```

```

C      DATE 1989-10-06

```

```

C      PURPOSE
C      INITIATES THE NETWORK

```

```

C      DESCRIPTION
C      FIRST ALL PARALLELL IMPEDANCES IN BRANCH POINTS ARE ASSIGNED A
C      LARGE VALUE
C      THEN FOR EACH TRANSFORMER THE TOPOLOGY IS INITIATED AND IMPEDANES COMPUTED
C      BY A CALL OF FILLIN

```

```

C      SEE ALSO
C      FILES      ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON      ALL COMMON VARIABLES DECLARED
C      TESTS      FILLIN      INITIATE THE NETWORK AND COMPUTES IMPEDANCES
C      TESTS      TKESTP      TAKES A STEP UP OR DOWN IN A NONINITIATED NETWORK
C      TESTS      NEWDIR      COMPUTES NEW DIRECTION IN A NONINITIATED NETWORK
C      MOVING      STEPDO      TAKES ONE STEP DOWN

```

```

C      VARIABLES
C      IF 'WR'=TRUE THEN TEXT SHOWING THE CURRENT POINT WILL BE WRITTEN
C      ON THE SCREEN

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER POINT(2),OLDP(2),DIR,BRA
INTEGER TRA,FEE
INTEGER I,J,K
LOGICAL WR

```

```

DO 10 I=1,MAXNRB
    BRAPOI(I,2)=0
    DO 20 J=3,(2+BRAPOI(I,1))
        BRAPH(I,J,1)=1.0E+20
        BRAPH(I,J,2)=0.0
20    CONTINUE
10    CONTINUE
DO 30 I=1,MAXNRS
    CHGSWI(I)=.FALSE.
30    CONTINUE
DO 50 I=1,MAXNRT

```

```

TRA=I
TRAFOS(TRA,6)=1.001E+06
TRAFOS(TRA,7)=1.001E+06
POINT(1)=TRATYP
POINT(2)=TRA
OLDP(1)=TRATYP
OLDP(2)=TRA
DIR=DOWN
BRA=4
CALL STEPDO(POINT,BRA,WR)
IF (POINT(1).EQ.SWITYP) THEN
  K=SWITCH(POINT(2),3)
  IF (CABTOP(K,1).EQ.OLDP(1) .AND. CABTOP(K,2).EQ.OLDP(2)) THEN
    SWITCH(POINT(2),2)=3
  ELSEIF (CABTOP(K,3).EQ.OLDP(1) .AND. CABTOP(K,4).EQ.OLDP(2)) THEN
    SWITCH(POINT(2),2)=3
  ELSE
    SWITCH(POINT(2),2)=4
  ENDIF
ELSE
  WRITE(*,*)'      DEGENERATED NET.SWITCH EXPECTED'
  STOP
ENDIF
CALL FILLIN(POINT,DIR,BRA,TRA,FEE)
CALL NEWDIR(POINT,DIR,BRA,WR)
60 IF (.NOT. (POINT(1).EQ.TRATYP .AND. POINT(2).EQ.TRA)) THEN
  CALL TKESTP(POINT,OLDP,DIR,BRA,WR)
  CALL FILLIN(POINT,DIR,BRA,TRA,FEE)
  CALL NEWDIR(POINT,DIR,BRA,WR)
  GOTO 60
ENDIF
50 CONTINUE
WRITE(*,*)'      INIT PASSED'
RETURN
END

```

C CCC

C COMP.FOR

C CCC

SUBROUTINE CON(POINT,OK)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
C CONNECTS A GROUP OF SECTIONS TO THE FEEDER
C WITH THE LARGEST MARGIN OF GETTING VIOLATED

C DESCRIPTION:
C STARTS A SEARCH FROM THE ROOT OF A TREE AND
C EXAMINES ALL POSSIBLE LOAD TRANSFERS. WHEN
C THE BEST SWITCH FOR LOAD TRANSFER IS FOUND
C THIS SWITCH IS OPENED
C IF HOWEVER NO POSSIBLE CONNECTION IS AVAILABLE
C OK EQUALS FALSE.

C VARIABLES:
C POINT IS THE ROOT OF THE TREE (TYPE,NR)

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
FETCH	GETPOI	GETS A POINT ON THE OTHER SIDE OF SWITCH
FETCH	GETVOL	GETS THE VOLTAGE ON THE OTHER SIDE
FETCH	GETFEE	GETS THE NUMBER OF THE FEEDER ON THE OTHER SIDE OF THE SWITCH
FETCH	GETPLI	GETS THE TRANSFORMERS MARGIN OF GETTING VIOLATED REGARDING ACTIVE POWER
FETCH	GETQLI	GETS THE TRANSFORMERS MARGIN OF GETTING VIOLATED REGARDING REACTIVE POWER
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
COMP	FORBAC	CHANGES DIRECTIONS IN THE TOPOLOGY DUE TO THE LOAD TRANSFER
COMP	COMPIM	COMPUTES NEW IMPEDANCE FOR A FEEDER
COMP	FILPOW	COMPUTES NEW POWER FOR A TRANSFORMER
COMP	FILVOL	COMPUTES ALL NODE VOLTAGES AND CABLE CURRENTS FOR A FEEDER

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER POINT(2),P1(2),OLD(2)
LOGICAL READY,NEG
INTEGER FEEDER,GETFEE,FEE

```
INTEGER TRAF0,DIR,BRA,SWI
LOGICAL VIO
REAL BEST,TEMP,VOL,P,Q
REAL ABS,GETVOL
REAL GETPLI,GETQLI
LOGICAL FOUND,VAL,OK
```

```
EXTERNAL GETPOI,GETVOL,GETFEE
EXTERNAL GETPLI,GETQLI
EXTERNAL MOVE,DIREC
```

```
VAL=.FALSE.
P1(1)=POINT(1)
P1(2)=POINT(2)
READY=.FALSE.
DIR=DOWN
IF (P1(1).EQ.BRATYP) THEN
  IF (BRAPOI(P1(1),2).EQ.3) THEN
    BRA=4
  ELSE
    BRA=3
  ENDIF
ENDIF
BEST=-1.0E38
IF (P1(1).EQ.NODTYP) THEN
  IF (NODETO(P1(2),2).EQ.0 .OR. NODETO(P1(2),3).EQ.0) THEN
    GOTO 20
  ENDIF
ENDIF
10 IF (.NOT. READY) THEN
  IF (P1(1).EQ.NODTYP .AND. DIR.EQ.DOWN) THEN
```

```
C ADDS UP THE MAGNITUDE OF ISOLATED LOAD
```

```
ISOLAT(1)=ISOLAT(1)+NODEPH(P1(2),1)
ISOLAT(2)=ISOLAT(2)+NODEPH(P1(2),2)
ENDIF
CALL MOVE(P1,OLD,DIR,BRA,.FALSE.)
CALL DIREC(P1,OLD,DIR,BRA)
IF (P1(1).EQ.SWITYP) THEN
  IF (SWITCH(P1(2),1).EQ.OPEN) THEN
    CALL REALFE(P1(2),FEE,FOUND)
    IF (FOUND) THEN
```

```
C A FEEDER THAT CAN TAKE UP LOAD IS FOUND
C IT IS POSSIBLE TO MAKE A CONNECTION
C BELOW FOLLOWS COMPUTATIONS IN ORDER TO
C DECIDE THE BEST CONNECTION
```

```
VAL=.TRUE.
VOL=GETVOL(P1(2))
P=GETPLI(P1(2))
Q=GETQLI(P1(2))
NEG=(VOL.LT.0.0 .OR. P.LT.0.0)
NEG=(NEG .OR. Q.LT.0.0)
IF (NEG) THEN
  TEMP=- (ABS(VOL)+ABS(P)+ABS(Q))
ELSE
  TEMP=VOL+P+Q
ENDIF
IF (TEMP.GT.BEST) THEN
```

```
C THE BEST CONNECTION IS STORED FOR LATER USE
```

```
BEST=TEMP
FEEDER=FEE
```



```
INCLUDE 'CONSTANTS.FOR'  
INCLUDE 'COMMON.FOR'
```

```
INTEGER P(2),OLD(2)  
LOGICAL READY  
INTEGER DIR,BRA
```

```
P(1)=SWITYP  
P(2)=SWI  
READY=.FALSE.  
DIR=DOWN  
BRA=3
```

```
10 IF (.NOT. READY) THEN  
    CALL MOVE(P,OLD,DIR,BRA,.FALSE.)  
    CALL DIREC(P,OLD,DIR,BRA)  
    IF (DIR.EQ.UP) THEN  
        IF (P(1).EQ.NODTYP) THEN  
            NODETO(P(2),4)=FEEDER  
        ELSEIF (P(1).EQ.FEETYP) THEN  
            MSFEED(P(2),1)=TRAFO  
        ENDIF  
    ENDIF  
    READY=(P(1).EQ.SWITYP .AND. P(2).EQ.SWI)  
    GOTO 10  
ENDIF  
RETURN  
END
```

```
SUBROUTINE FORBAC(SWI,TRAFO,FEEDER)
```

```
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

```
C VERSION 1
```

```
C DATE 1989-09-30
```

```
C PURPOSE:  
C CHANGES DIRECTIONS.  
C WHEN A LOAD TRANSFER IS MADE THE FORMER DIRECTION 'DOWN'  
C WILL BECOME 'UP'
```

```
C DESCRIPTION:  
C STARTING AT THE SWITCH THAT WILL BE OPENED THIS ROUTINE  
C MOVES TOWARDS THE SWITCH THAT WILL BE CLOSED CHANGING  
C DIRECTIONS. THE ACTUAL CHANGE IN POINTERS IS DONE BY A  
C CALL OF SUBROUTINE CHGDIR
```

```
C VARIABLES:  
C SWI IS THE NUMBER OF THE SWITCH THAT WILL BE OPENED  
C TRAFO IS THE NUMBER OF THE 'NEW' TRANSFORMER THAT A  
C FEEDER POSSIBLY WILL BE ATTACHED TO  
C FEEDER IS THE NUMBER OF THE 'NEW' FEEDER THAT A NODE  
C WILL BE ATTACHED TO
```

```
C SEE ALSO:
```

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
COMP	UDAT	ASSIGNS NEW VALUES OF TRAFO,FEEDER
COMP	CHGDIR	CHANGES POINTERS DUE TO A LOAD TRANSFER
MOVING	STEPUP	TAKES ONE STEP UP
MOVING	STEPDO	TAKES ONE STEP DOWN

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER TRAF0,FEEDER,I
INTEGER SWI,DIR,BRA
INTEGER POI(2),OLD(2)
LOGICAL READY

POI(1)=SWITYP
POI(2)=SWI
OLD(1)=SWITYP
OLD(2)=SWI
READY=.FALSE.
CALL STEPUP(POI,.FALSE.)
I=SWITCH(SWI,3)
IF (CABTOP(I,1).EQ.POI(1) .AND. CABTOP(I,2).EQ.POI(2)) THEN
 SWITCH(SWI,2)=4
ELSEIF (CABTOP(I,3).EQ.POI(1) .AND. CABTOP(I,4).EQ.POI(2)) THEN
 SWITCH(SWI,2)=4
ELSE
 SWITCH(SWI,2)=3
ENDIF

10 IF (.NOT. READY) THEN
 CALL CHGDIR(POI,OLD,BRA,TRAF0,FEEDER)
 OLD(1)=POI(1)
 OLD(2)=POI(2)
 CALL STEPDO(POI,BRA,.FALSE.)
 IF (POI(1).EQ.SWITYP) THEN
 IF (SWITCH(POI(2),1).EQ.OPEN) THEN
 READY=.TRUE.
 I=SWITCH(POI(2),3)
 IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
 SWITCH(POI(2),2)=3
 ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
 SWITCH(POI(2),2)=3
 ELSE
 SWITCH(POI(2),2)=4
 ENDIF
 I=SWITCH(POI(2),2)
 I=SWITCH(POI(2),I)
 ENDIF
 ENDIF
 GOTO 10
ENDIF
CALL UDAT(SWI,TRAF0,FEEDER)
RETURN
END

SUBROUTINE CHGDIR(POI,OLD,BRA,TRAF0,FEEDER)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
C WHEN A LOAD TRANSFER IS PERFORMED CHANGES MUST BE MADE IN THE
C TOPOLOGY. THIS SUBROUTINE CHANGES POINTER ACCORDING TO THE
C NEW TOPOLOGY

C DESCRIPTION:
C IT IS ONLY TO BE CALLED BY SUBROUTINE FORBAC
C STANDING AT THE POINT POI CHANGES IS MADE ACCORDING TO THE TYPE
C OF THIS POINT.

C VARIABLES:
C POI IS THE CURRENT POINT IN THE NETWORK
C OLD IS THE PRVIOUSLY VISITED POINT IN THE NETWORK
C BRA IS AN OUTPARAMETER. IT INDICATES THE OLD VALUE OF UP.
C TRAF0 IS THE NUMBER OF THE TRANSFORMER WHICH ALL OTHER
C ELEMENTS ARE ATTACHED TO.
C FEEDER IS THE NUMBER OF THE FEEDER ENERGIZING THE TRANSFERED LOAD

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
COMP	FORBAC	CHANGES THE NETWORK TOPOLOGY

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER TRAF0,FEEDER,BRA
INTEGER POI(2),OLD(2)
INTEGER I,J

IF (POI(1).EQ.SWITYP) THEN
 I=SWITCH(POI(2),3)
 IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
 SWITCH(POI(2),2)=3
 ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
 SWITCH(POI(2),2)=3
 ELSE
 SWITCH(POI(2),2)=4
 ENDIF

ELSEIF (POI(1).EQ.NODTYP) THEN
 I=NODETO(POI(2),2)
 IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
 NODETO(POI(2),1)=2
 ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
 NODETO(POI(2),1)=2
 ELSE
 NODETO(POI(2),1)=3
 ENDIF

 I=NODETO(POI(2),1)
 I=NODETO(POI(2),1)
 NODETO(POI(2),4)=FEEDER
ELSEIF (POI(1).EQ.BRATYP) THEN
 BRA=BRAPOI(POI(2),2)
 DO 10 J=3,(2+BRAPOI(POI(2),1))
 I=BRAPOI(POI(2),J)
 IF (CABTOP(I,1).EQ.OLD(1) .AND. CABTOP(I,2).EQ.OLD(2)) THEN
 BRAPOI(POI(2),2)=J
 ELSEIF (CABTOP(I,3).EQ.OLD(1) .AND. CABTOP(I,4).EQ.OLD(2)) THEN
 BRAPOI(POI(2),2)=J
 ENDIF

10 CONTINUE

ELSEIF (POI(1).EQ.FEETYP) THEN
 I=MSFEED(POI(2),2)
 MSFEED(POI(2),2)=MSFEED(POI(2),3)
 MSFEED(POI(2),3)=I
 MSFEED(POI(2),1)=TRAF0

```

ELSE
  WRITE(*,*) '      WRONG TYPE IN CHGDIR',POI(1),POI(2)
  STOP
ENDIF
OLD(1)=POI(1)
OLD(2)=POI(2)
RETURN
END

```

SUBROUTINE COMPIM(FEEDER)

CC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
 C COMPUTES THE IMPEDANCE FOR FEEDER 'FEEDER'

C METHOD:
 C COMPUTATIONS IS BASED ON THE FACT THAT THE DISTRIBUTION NETWORK
 C HAS A RADIAL STRUCTURE. THEREFORE PARALLELL IMPEDANCES WITH CABLE
 C IMPEDANCES WILL BE COMPUTED. SEE ALSO 'SERVICE RESTORATION IN
 C ELECTRIC DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND (SECTION
 C COMPUTATIONS)

C DESCRIPTION:
 C TRAVERSES THE TREE WITH 'FEEDER' AS ROOT.
 C STORES PARALLEL IMPEDANCES IN EVERY BRANCH POINT
 C COMPUTATIONS IS ONLY MADE WHEN MOVING TOWARDS THE TOP
 C WHEN THE ROOT IS REACHED THE SUBROUTINE CONTINUES COMPUTING
 C IMPEDANCES UNTIL THE TOP IS REACHED

C VARIABLES:
 C FEEDER IS THE NUMBER OF THE FEEDER FOR WICH THE IMPEDANCE IS COMPUTED
 C IMPEDANCES IN OHM

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
COMP	PARALL	THE ABOVE ARE USED TO TRAVERSE THE TREE COMPUTES THE EQUIVALENT TO TWO PARALLELL IMPEDANCES
COMP	BRAPAR	COMPUTES THE EQUIVALENT IMPEDANCE FOR A BRANCH POINT

CC

INCLUDE 'CONSTANTS.FOR'
 INCLUDE 'COMMON.FOR'

INTEGER POI(2),OLD(2)
 LOGICAL READY,FIRST
 INTEGER A,DIR,BRA,I,J
 REAL SUMR,SUMX,SUM,RES,REA,F(2),G(2)

POI(1)=FEETYP
 POI(2)=FEEDER

```

READY=.FALSE.
OLD(1)=POI(1)
OLD(2)=POI(2)
DIR=DOWN
BRA=3
10 IF (.NOT. READY) THEN
    CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
    CALL DIREC(POI,OLD,DIR,BRA)
    IF (DIR.EQ.UP) THEN
        IF (POI(1).EQ.NODTYP) THEN
            IF (FIRST) THEN
C      A LEAF IS FOUND. COMPUTATIONS START

                RES=NODEPH(POI(2),4)
                REA=NODEPH(POI(2),5)
                FIRST=.FALSE.
            ELSE
C      THE NODE IS NOT A LEAF. PARALLELL IMPEDANCE IS COMPUTED INCLUDING
C      CABLE IMPEDANCE

                G(1)=RES
                G(2)=REA
                F(1)=NODEPH(POI(2),4)
                F(2)=NODEPH(POI(2),5)
                CALL PARALL(F,G,RES,REA)
                I=NODETO(POI(2),1)
                I=NODETO(POI(2),I)
                RES=CABPHY(I,1)+RES
                REA=CABPHY(I,2)+REA
            ENDIF
            ELSEIF (POI(1).EQ.SWITYP) THEN
                IF (.NOT. FIRST) THEN
C      ADDS UP CABLE IMPEDANCE

                    I=SWITCH(POI(2),2)
                    I=SWITCH(POI(2),I)
                    RES=RES+CABPHY(I,1)
                    REA=REA+CABPHY(I,2)
                ELSE
                    RES=1.0E+30
                    REA=1.0E+30
                ENDIF

                ELSEIF (POI(1).EQ.TRATYP) THEN
C      STORE IMPEDANCE FOR LATER USE

                    TRAFOS(POI(2),6)=RES
                    TRAFOS(POI(2),7)=REA

                ELSEIF (POI(1).EQ.BRATYP) THEN
C      COMPUTES THE PARALLEL IMPEDANCE AND SAVES IT IN BRAPH FOR LATER USE

                    DO 20 I=3,(2+BRAPOI(POI(2),1))
                        J=BRAPOI(POI(2),I)
                        IF (CABTOP(J,1).EQ.OLD(1) .AND. CABTOP(J,2).EQ.OLD(2)) THEN
                            BRAPH(POI(2),I,1)=RES
                            BRAPH(POI(2),I,2)=REA
                        ELSEIF (CABTOP(J,3).EQ.OLD(1) .AND. CABTOP(J,4).EQ.OLD(2)) THEN
                            BRAPH(POI(2),I,1)=RES
                            BRAPH(POI(2),I,2)=REA
                        ENDIF
                    ENDIF
                ENDIF
            ENDIF
        ENDIF
    ENDIF
END IF

```

```

20      CONTINUE
      CALL BRAPAR(POI(2),RES,REA)
      I=BRAPOI(POI(2),2)
      I=BRAPOI(POI(2),1)
      RES=RES+CABPHY(I,1)
      REA=REA+CABPHY(I,2)

      ELSEIF (POI(1).EQ.FEETYP) THEN
      I=MSFEED(POI(2),2)
      RES=RES+CABPHY(I,1)
      REA=REA+CABPHY(I,2)
      ELSE
      WRITE(*,*)'      WRONG TYPE IN COMPIM',POI(1),POI(2)
      STOP
      ENDIF
      ELSEIF (DIR.EQ.DOWN) THEN

C      IF A BRANCH POINT IS REACHED AND A NEW DIRECTION DOWN HAS BEEN
C      COMPUTED RES AND REA CONTAINING THE IMPEDANCE OF THE BRANCH THAT
C      HAS BEEN VISITED MUST BE SAVED FOR LATER USE (Z=(RES,j*REA))

      IF (POI(1).EQ.BRATYP) THEN
      DO 30 I=3,(2+BRAPOI(POI(2),1))
      A=BRAPOI(POI(2),I)
      IF (CABTOP(A,1).EQ.OLD(1) .AND. CABTOP(A,2).EQ.OLD(2)) THEN
      J=I
      ELSEIF (CABTOP(A,3).EQ.OLD(1) .AND. CABTOP(A,4).EQ.OLD(2)) THEN
      J=I
      ENDIF
30      CONTINUE
      IF (J.NE.BRAPOI(POI(2),2)) THEN
      BRAPH(POI(2),J,1)=RES
      BRAPH(POI(2),J,2)=REA
      ENDIF
      ENDIF
      FIRST=.TRUE.
      ENDIF
      READY=POI(1).EQ.FEETYP
      GOTO 10
      ENDIF
      READY=.FALSE.
35      IF (.NOT. READY) THEN

C      HERE THE SAME COMPUTATIONS AS THE ABOVE IS PERFORMED UNTIL
C      THE TRANSFORMER IS REACHED

      OLD(1)=POI(1)
      OLD(2)=POI(2)
      CALL STEPUP(POI,.FALSE.)
      IF (POI(1).EQ.SWITYP) THEN
      IF (.NOT. FIRST) THEN
      I=SWITCH(POI(2),2)
      I=SWITCH(POI(2),I)
      RES=RES+CABPHY(I,1)
      REA=REA+CABPHY(I,2)
      ENDIF
      ELSEIF (POI(1).EQ.TRATYP) THEN
      TRAFOS(POI(2),6)=RES
      TRAFOS(POI(2),7)=REA
      ELSEIF (POI(1).EQ.BRATYP) THEN
      DO 40 I=3,(2+BRAPOI(POI(2),1))
      J=BRAPOI(POI(2),I)
      IF (CABTOP(J,1).EQ.OLD(1) .AND. CABTOP(J,2).EQ.OLD(2)) THEN
      BRAPH(POI(2),I,1)=RES
      BRAPH(POI(2),I,2)=REA
      ELSEIF (CABTOP(J,3).EQ.OLD(1) .AND. CABTOP(J,4).EQ.OLD(2)) THEN

```

```

        BRAPH(POI(2),I,1)=RES
        BRAPH(POI(2),I,2)=REA
    ENDIF
40    CONTINUE
    CALL BRAPAR(POI(2),RES,REA)
    I=BRAPOI(POI(2),2)
    I=BRAPOI(POI(2),1)
    RES=RES+CABPHY(I,1)
    REA=REA+CABPHY(I,2)
    ENDIF
    READY=POI(1).EQ. TRATYP
    GOTO 35
ENDIF
RETURN
END

```

SUBROUTINE FILPOW(TRAFO)

CC

AUTHOUR MAGNUS BERGSTRAND

VERSION 1

DATE 1989-09-30

PURPOSE:
 COMPUTES AND STORES THE POWER OF THE TRANSFORMER 'TRAFO'

DESCRIPTION:
 FIRST THE IMPEDANCE FOR THE TRANSFORMER MUST BE COMPUTED
 USING COMPIM
 THEN THE POWER IS COMPUTED USING $I=I/Z$, $P=Z*I**2$

VARIABLES:
 TRAFO IS THE NUMBER OF THE TRANSFORMER
 POWER IN VA
 VOLTAGE IN VOLT
 IMPEDANCE IN AMPERE

SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
COMP	COMPIM	COMPUTES THE IMPEDANCE FOR ONE FEEDER

CC

INCLUDE 'CONSTANTS.FOR'
 INCLUDE 'COMMON.FOR'

INTEGER TRAFO
 REAL CUR

CUR=TRAMAX/SQRT(TRAFOS(TRAFO,6)**2+TRAFOS(TRAFO,7)**2)
 TRAFOS(TRAFO,3)=TRAFOS(TRAFO,6)*CUR**2
 TRAFOS(TRAFO,4)=TRAFOS(TRAFO,7)*CUR**2
 RETURN
 END

```

SUBROUTINE COMPVO(R,X,CUR,CURFI,VOLT,FI)
C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C AUTHOUR MAGNUS BERGSTRAND
C
C VERSION 1
C
C DATE 1989-09-30
C
C PURPOSE:
C COMPUTES NODE VOLTAGES
C
C DESCRIPTION:
C USES THE FORMULA  $V_2=V_1-Z*I$ 
C
C VARIABLES:
C THE IMPEDANCE Z IS REPRESENTED BY (R,j*X) BOTH IN OHMS
C THE CURRENT I IS EXPONENTIAL FORM  $I=CUR*EXP(CURFI)$ 
C CUR IN AMPERE,CURFI IN RADIANS
C THE VOLTAGE ALSO IN EXPONENTIAL FORM (BOTH IN AND OUT)
C VOLT IN VOLTS AND FI IN RADIANS
C
C SEE ALSO:
C
C FILE          ROUTINE          COMMENT
C CONSTANTS          ALL CONSTANTS DECLARED
C COMMON          ALL COMMON VARIABLES DECLARED
C COMP          COMPCU          COMPUTES THE CABLE CURRENT
C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C REAL R,X,CUR,CURFI,VOLT,FI
C REAL ZABS,ZFI
C REAL REA,IM
C
C ZABS=SQRT(R**2+X**2)
C ZFI=ATAN(X/R)
C REA=VOLT*COS(FI)-CUR*ZABS*COS(CURFI+ZFI)
C IM=VOLT*SIN(FI)-CUR*ZABS*SIN(CURFI+ZFI)
C VOLT=SQRT(REA**2+IM**2)
C FI=ATAN(IM/REA)
C RETURN
C END
C
C
C SUBROUTINE COMPCU(BRAFI,BRACUR,CURREN,FI,POI,BRA)
C
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C AUTHOUR MAGNUS BERGSTRAND
C
C VERSION 1
C
C DATE 1989-09-30
C
C PURPOSE:
C COMPUTES THE CABLE CURRENT IN TWO CASES:
C 1. AFTER A NODE
C 2. THE CURRENT DOWN IN THE BRANCH 'BRA' IF 'POI' IS A BRANCH POINT
C
C DESCRIPTION:
C 1. SINCE THE VOLTAGE IS KNOWN THE CURRENT IS DECREMENTED BY U/Z

```


C 2. ALL IMPEDANCES FOR THE DIFFERENT BRANCHES ARE KNOWN AS WELL
C AS THE VOLTAGE THE CURRENT IS EASILY COMPUTED USING THE LAWS
C OF KIRCHHOFF

C VARIABLES:
C BRACUR IN AMPERE (CURRENT FROM FEEDER TO BRANCHPOINT)
C BRAFI IN RADIAN (PHASE)
C CURREN IN AMPERE (1. IN AND OUT 2. ONLY OUT)
C FI IN RADIAN (1. IN AND OUT 2. ONLY OUT)

C SEE ALSO:

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
COMP	COMPIM	COMPUTES IMPEDANCE OF A FEEDER
COMP	COMPVO	COMPUTES VOLTAGE
COMP	FILVOL	COMPUTES AND FILLS IN NODE VOLTAGES AND CABLE CURRENTS FOR ONE FEEDER

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

REAL BRAFI, BRACUR
REAL CURREN, FI
INTEGER POI(2), BRA
COMPLEX ZP, Z, CUR
REAL RES, REA, VO
REAL ALFA1, ZABS

IF (POI(1).EQ.BRATYP) THEN
CALL BRAPAR(POI(2), RES, REA)
ZP=CMPLX(RES, REA)
Z=CMPLX(BRAPH(POI(2), BRA, 1), BRAPH(POI(2), BRA, 2))
CUR=CMPLX(BRACUR*COS(BRAFI), BRACUR*SIN(BRAFI))
CUR=CUR*ZP/Z
CURREN=ABS(CUR)
FI=ATAN((AIMAG(CUR))/(REAL(CUR)))
ELSEIF (POI(1).EQ.NODTYP) THEN
VO=NODEPH(POI(2), 3)
RES=NODEPH(POI(2), 4)
REA=NODEPH(POI(2), 5)
ALFA1=ATAN(REA/RES)-FI
ZABS=SQRT(RES**2+REA**2)
CURREN=CURREN-COS(ALFA1)*VO/ZABS
FI=FI+ATAN(VO*SIN(ALFA1)/CURREN/ZABS)
ELSE
WRITE(*,*)'WRONG TYPE IN COMPCU'
STOP
ENDIF
RETURN
END

SUBROUTINE FILVOL(FEEDER, VIO)

C CCC

C AUTHOUR MAGNUS BERGSTRAND
C VERSION 1
C DATE 1989-09-30

```

C      PURPOSE:
C      COMPUTES AND FILLS IN NODE VOLTAGES AND CABLE CURRENTS FOR ONE FEEDER

C      DESCRIPTION:
C      TRAVERSES THE TREE WITH 'FEEDER' AS ROOT.
C      WHEN MOVING DOWNWARDS COMPUTING NEW CABLE CURRENTS AND NODE VOLTAGES
C      IF ANY VIOLATIONS OCCUR 'VIO' BECOMES TRUE AND THE COMMON DECLARED
C      VIOVEC (VIOLATION VECTOR) DIFFERS FROM ZERO
C      ONLY THE BIGGEST LINE CAPACITY VIOLATION IS STORED IN THE FIRST ELEMENT
C      THE SUM OF VOLTAGE DROP CONSTRAINTS ARE STORED IN SECOND ELEMENT

C      VARIABLES:
C      FEEDER IS THE NUMBER OF THE FEEDER
C      VIO INDICATES IF A VIOLATION OCCURED

C      SEE ALSO:

C      FILE          ROUTINE          COMMENT
C      CONSTANTS          ALL CONSTANTS DECLARED
C      COMMON            ALL COMMON VARIABLES DECLARED
C      MOVING            MOVE            TAKES ONE STEP UP OR DOWN
C      MOVING            DIREC          COMPUTES NEW DIRECTION
C      MOVING            STEPUP         TAKES ONE STEP UP
C      MOVING            STEPDOWN      TAKES ONE STEP DOWN
C      COMP              COMPV          COMPUTES NODE VOLTAGE
C      COMP              COMPC          COMPUTES CURRENT

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER FEEDER
REAL R,X,VOL,BRAVOL(MAXNRB),RBSUM,XSUM,SUM,CURREN
INTEGER I,J,K,A,B
INTEGER WAY(MAXNRB)
INTEGER BRA,DIR,POI(2),OLD(2)
LOGICAL READY,VIO,ENDNOD
REAL CUR,BRACUR(MAXNRB)
REAL SQRT
REAL NODFI(MAXNRN),BRAFI(MAXNRB)
REAL CURFI,VOLFI,BCURFI(MAXNRB)

INTRINSIC SQRT

VIO=.FALSE.
VIOVEC(FEEDER,1)=0.0
VIOVEC(FEEDER,2)=0.0
POI(1)=FEETYP
POI(2)=FEEDER
READY=.FALSE.
I=1
10  IF (.NOT. READY) THEN

C      MOVES TO TRANSFORMER MEMORIZING PATH

      OLD(1)=POI(1)
      OLD(2)=POI(2)
      CALL STEPUP(POI,.FALSE.)
      IF (POI(1).EQ.BRATYP) THEN
        DO 20 J=3,(2+BRAPOI(POI(2),1))
          K=BRAPOI(POI(2),J)
          IF (CABTOP(K,1).EQ.OLD(1) .AND. CABTOP(K,2).EQ.OLD(2)) THEN
            WAY(I)=J
            I=I+1

```

```

        ELSEIF (CABTOP(K,3).EQ.OLD(1) .AND. CABTOP(K,4).EQ.OLD(2)) THEN
            WAY(I)=J
            I=I+1
        ENDIF
20    CONTINUE
    ENDIF
    READY=(POI(1).EQ. TRATYP)
    GOTO 10
ENDIF
READY=.FALSE.

C    INITIATE

R=0.0
X=0.0
VOL=TRAMAX
VOLFI=0.0
DO 25 J=1,MAXNRB
    BRAVOL(J)=0.0
    BRACUR(J)=0.0
    BRAFI(J)=0.0
    BCURFI(J)=0.0
25    CONTINUE

C    STARTS COMPUTING ON THE WAY BACK TO FEEDER

CUR=TRAMAX/SQRT( TRAFOS(POI(2),6)**2+TRAFOS(POI(2),7)**2)
CURFI=ATAN( TRAFOS(POI(2),7)/TRAFOS(POI(2),6))
BRA=4
30    IF (.NOT. READY) THEN
        CALL STEPDO(POI,BRA,.FALSE.)
        IF (POI(1).EQ. BRATYP) THEN
            I=I-1
            BRA=WAY(I)
            J=BRAPOI(POI(2),2)
            J=BRAPOI(POI(2),J)
            R=R+CABPHY(J,1)
            X=X+CABPHY(J,2)
            IF (CUR.GT. CABPHY(J,3)) THEN
                VIO=.TRUE.
                IF ((CUR-CABPHY(J,3)).GT. VIOVEC(FEEDER,1)) THEN
                    VIOVEC(FEEDER,1)=VIOVEC(FEEDER,1)+CUR-CABPHY(J,3)
                ENDIF
            ENDIF
            CABPHY(J,4)=CUR
            BRACUR(POI(2))=CUR
            BCURFI(POI(2))=CURFI
            CALL COMPVO(R,X,CUR,CURFI,VOL,VOLFI)
            BRAVOL(POI(2))=VOL
            BRAFI(POI(2))=VOLFI
            IF (VOL.LT. VOLLIM) THEN
                VIO=.TRUE.
            ENDIF
            R=0.0
            X=0.0
            CALL COMPCU(BCURFI(POI(2)),BRACUR(POI(2)),CUR,FI,POI,BRA)
        ELSEIF (POI(1).EQ. SWITYP) THEN
            J=SWITCH(POI(2),2)
            J=SWITCH(POI(2),J)
            R=R+CABPHY(J,1)
            X=X+CABPHY(J,2)
            CABPHY(J,4)=CUR
            IF (CUR.GT. CABPHY(J,3)) THEN
                VIO=.TRUE.
                IF ((CUR-CABPHY(J,3)).GT. VIOVEC(FEEDER,1)) THEN
                    VIOVEC(FEEDER,1)=VIOVEC(FEEDER,1)+CUR-CABPHY(J,3)
                ENDIF
            ENDIF
        ENDIF
    ENDIF

```

```

        ENDIF
    ENDIF
ELSEIF (POI(1).EQ.FEETYP) THEN
    J=MSFEED(POI(2),2)
    R=R+CABPHY(POI(2),1)
    X=X+CABPHY(POI(2),2)
ELSE
    WRITE(*,*)'      DEGENERATED NETWORK.FILVOL'
    STOP
ENDIF
READY=(POI(1).EQ.FEETYP .AND. POI(2).EQ.FEEDER)
GOTO 30
ENDIF

```

```

C      NOW THE FEEDER IS REACHED
C      FROM HERE ON THE TREE IS TRAVERSED COMPUTING CURRENTS AND VOLTAGES

```

```

READY=.FALSE.
DIR=DOWN
40  IF (.NOT. READY) THEN
    CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
    IF (DIR.EQ.DOWN) THEN
        ENDNOD=.TRUE.
        IF (POI(1).EQ.SWITYP) THEN
            IF (SWITCH(POI(2),1).EQ.CLOSED) THEN
                J=SWITCH(POI(2),2)
                J=SWITCH(POI(2),J)
                R=R+CABPHY(J,1)
                X=X+CABPHY(J,2)
                IF (CUR.GT.CABPHY(J,3)) THEN
                    VIO=.TRUE.
                    IF ((CUR-CABPHY(J,3)).GT.VIOVEC(FEEDER,1)) THEN
                        VIOVEC(FEEDER,1)=VIOVEC(FEEDER,1)+CUR-CABPHY(J,3)
                    ENDIF
                ENDIF
                CABPHY(J,4)=CUR
            ENDIF
        ELSEIF (POI(1).EQ.NODTYP) THEN
            J=NODETO(POI(2),1)
            J=NODETO(POI(2),J)
            R=R+CABPHY(J,1)
            X=X+CABPHY(J,2)
            CABPHY(J,4)=CUR
            CALL COMPVO(R,X,CUR,CURFI,VOL,VOLFI)
            NODEPH(POI(2),3)=VOL
            IF (VOL.LT.VOLLIM) THEN
                VIO=.TRUE.
            ENDIF
            IF (CUR.GT.CABPHY(J,3)) THEN
                VIO=.TRUE.
                IF ((CUR-CABPHY(J,3)).GT.VIOVEC(FEEDER,1)) THEN
                    VIOVEC(FEEDER,1)=VIOVEC(FEEDER,1)+CUR-CABPHY(J,3)
                ENDIF
            ENDIF
            CURREN=VOL/SQRT(NODEPH(POI(2),4)**2+NODEPH(POI(2),5)**2)
            NODEPH(POI(2),1)=NODEPH(POI(2),4)*CURREN**2
            NODEPH(POI(2),2)=NODEPH(POI(2),5)*CURREN**2
            CALL COMPCU(1.0,1.0,CUR,CURFI,POI,BRA)
            R=0.0
            X=0.0
        ELSEIF (POI(1).EQ.BRATYP) THEN
            J=BRAPOI(POI(2),2)
            J=BRAPOI(POI(2),J)
            R=R+CABPHY(J,1)
            X=X+CABPHY(J,2)
            CALL COMPVO(R,X,CUR,CURFI,VOL,VOLFI)

```

```

BRAVOL(POI(2))=VOL
BRAFI(POI(2))=VOLFI
IF (CUR.GT.CABPHY(J,3)) THEN
  VIO=.TRUE.
  IF ((CUR-CABPHY(J,3)).GT.VIOVEC(FEEDER,1)) THEN
    VIOVEC(FEEDER,1)=VIOVEC(FEEDER,1)+CUR-CABPHY(J,3)
  ENDIF
ENDIF
CABPHY(J,4)=CUR
BCURFI(POI(2))=CURFI
BRACUR(POI(2))=CUR
ELSE
  WRITE(*,*)'      WRONG IN NETWORK.FILVOL'
  WRITE(*,*)'      POINT IS',POI(1),POI(1)
  STOP
ENDIF
ENDIF
CALL DIREC(POI,OLD,DIR,BRA)
IF (DIR.EQ.DOWN .AND. POI(1).EQ.BRATYP) THEN
  CALL COMPCU(BCURFI(POI(2)),BRACUR(POI(2)),CUR,CURFI,POI,BRA)
  VOL=BRAVOL(POI(2))
  VOLFI=BRAFI(POI(2))
  R=0.0
  X=0.0
ENDIF
IF (DIR.EQ.UP .AND. ENDNOD) THEN
  IF (POI(1).EQ.NODTYP) THEN

```

C THIS IS THE ONLY PLACE NECESSARY TO CHECK THE VOLTAGE DROP CONSTRAINT
C ANY VIOLATION WILL ALWAYS OCCUR IN A LEAF NODE

```

  IF (NODEPH(POI(2),3).LT.VOLLIM) THEN
    VIOVEC(FEEDER,2)=VIOVEC(FEEDER,2)+VOLLIM-NODEPH(POI(2),3)
  ENDIF
  ENDNOD=.FALSE.
ENDIF
ENDIF
READY=(POI(1).EQ.FEETYP)
GOTO 40

```

ENDIF
RETURN
END

SUBROUTINE PARALL(A,B,RES,REA)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-09-30

C PURPOSE:
C COMPUTES THE IMPEDANCE (RES, j*REA)=A//B

C VARIABLES:
C A AND B ARE VECTORS (OHM, j*OHM)
C RES IS RESISTANCE (OUT) IN OHM
C REA IS REACTANCE (OUT) IN OHM

C SEE ALSO:

C FILE ROUTINE COMMENT
C CONSTANTS ALL CONSTANTS DECLARED

```

C      COMMON          ALL COMMON VARIABLES DECLARED
C      COMP           BRAPAR      COMPUTES PARALLELL IMPEDANCE WHEN POINT
C                               IS A BRANCH POINT

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

REAL A(2),B(2),RES,REA
COMPLEX A1,B1,C1,C2

```

```

A1=CMPLX(A(1),A(2))
B1=CMPLX(B(1),B(2))
C1=CMPLX(1,0)
C2=C1/A1+C1/B1
C2=C1/C2
RES=REAL(C2)
REA=AIMAG(C2)
RETURN
END

```

```

SUBROUTINE BRAPAR(NR,RES,REA)

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C      AUTHOUR MAGNUS BERGSTRAND

```

```

C      VERSION 1

```

```

C      DATE 1989-09-30

```

```

C      PURPOSE:
C      COMPUTES THE PARALLELL IMPEDANCE FOR BRANCH POINT 'NR'

```

```

C      VARIABLES:
C      THE IMPEDANCE (RES,j*REA) IN OHMS (OUT)

```

```

C      SEE ALSO:

```

FILE	ROUTINE	COMMENT
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
COMP	PARALL	COMPUTES THE PARALLELL IMPEDANCE BETWEEN TWO VECTORS

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER NR
REAL RES,REA
INTEGER I,J
COMPLEX A,B,C
LOGICAL READY

```

```

IF (BRAPOI(NR,2).EQ.3) THEN
  I=4
ELSE
  I=3
ENDIF
A=CMPLX(BRAPH(NR,I,1),BRAPH(NR,I,2))
B=CMPLX(1,0)
A=B/A
10 IF (1.LT.(BRAPOI(NR,1)+2)) THEN

```

```
I=I+1
IF (I.NE.BRAPOI(NR,2)) THEN
  C=CMPLX(BRAPH(NR,I,1),BRAPH(NR,I,2))
  A=A+B/C
ENDIF
GOTO 10
ENDIF
A=B/A
RES=REAL(A)
REA=AIMAG(A)
RETURN
END
```

C CCC

C STEP2.FOR

C CCC

SUBROUTINE STAGE(FEEDER,OK)

C CCC

AUTHOR MAGNUS BERGSTRAND

VERSION 1

DATE 1989-10-05

PURPOSE

WHEN A LOAD TRANSFER IS TO BE PERFORMED THIS SUBROUTINE EXAMINE DIFFERENT POSSIBILITIES AND MAKES A FILE 'STAGE' CONTAINING DATA ABOUT POSSIBLE SWITCHES

DESCRIPTION

TRAVERSES THE TREE WITH THE FEEDER 'FEEDER' AS ROOT. WHEN A CLOSED SWITCH IS FOUND THE ESTIMATED REDUCTION OF THE VIOLATION VECTOR IS COMPUTED AND STORED IN THE FILE 'STAGE'. ALSO THE NUMBERS OF THE SWITCHES INVOLVED, THE NODES PRIORITY AS WELL AS THE NUMBER OF THE FEEDER THAT WILL TAKE UP ANY LOAD WILL BE STORED. THIS SUBROUTINE IS ONLY TO BE CALLED FROM 'STAGE1'

SEE ALSO

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
STEP2	STAGE1	FIRST STAGE SUPPORT. TRANSFERS LOAD TO AN ADJACENT FEEDER.
STEP2	STAGE2	SECOND STAGE SUPPORT. TRANSFERS LOAD TO AN ADJACENT FEEDER WHICH IN ITS TURN TRANSFERS LOAD TO A SECONDARY SUPPORT FEEDER.
STEP2	IMP	COMPUTES THE ESTIMATED REDUCTION OF THE VIOLATION VECTOR
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
FETCH	GETFEE	GETS THE NUMBER OF THE FEEDER ON THE OTHER SIDE OF AN OPEN SWITCH

VARIABLES

FEEDER IS THE NUMBER OF THE FEEDER THAT WILL TRANSFER LOAD
OK INDICATES IF ANY POSSIBLE LOAD TRANSFERS WERE FOUND

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER FEEDER
LOGICAL OK
INTEGER I
INTEGER POI(2),OLD(2)
INTEGER DIR,BRA,SWI2,FEE


```
REAL SUM,LOAD,ALPHA
INTEGER GETFEE
LOGICAL VIOLTE
```

```
EXTERNAL GETFEE
EXTERNAL VIOLTE
```

```
WRITE(*,*)'ENTERS STAGE*****'
```

```
OPEN(1,FILE='STAGE',STATUS='NEW',ACCESS='SEQUENTIAL')
```

```
POI(1)=FEETYP
```

```
POI(2)=FEEDER
```

```
DIR=DOWN
```

```
BRA=3
```

```
CALL STEPDO(POI,BRA,.FALSE.)
```

```
OK=.FALSE.
```

```
10 IF (POI(1).NE.FEETYP) THEN
```

```
C TRAVERSES THE TREE
```

```
CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
```

```
CALL DIREC(POI,OLD,DIR,BRA)
```

```
IF (POI(1).EQ.SWITYP) THEN
```

```
IF (SWITCH(POI(2),1).EQ.OPEN .AND. SWITCH(POI(2),5).EQ.1) THEN
```

```
C A POSSIBLE LOAD TRANSFER IS FOUND
```

```
FEE=GETFEE(POI(2))
```

```
IF (.NOT. VIOLTE(FEE)) THEN
```

```
C THE FEEDER THAT WILL TAKE UP LOAD IS NOT VIOLATED
```

```
C COMPUTE THE IMPROVEMENT AND STORE DATA
```

```
OK=.TRUE.
```

```
CALL IMP(POI(2),SWI2,SUM,LOAD,ALPHA)
```

```
WRITE(1,*) POI(2),SWI2,SUM,FEE,LOAD,ALPHA
```

```
ENDIF
```

```
ENDIF
```

```
ENDIF
```

```
GOTO 10
```

```
ENDIF
```

```
CLOSE(1,STATUS='KEEP')
```

```
RETURN
```

```
END
```

```
SUBROUTINE IMP(SWI,SWI2,SUM,LOAD,ALPHA)
```

```
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

```
C VERSION 1
```

```
C DATE 1989-10-05
```

```
C PURPOSE
```

```
C COMPUTES THE ESTIMATED REDUCTION OF THE VIOLATION VECTOR IF  
C THE SWITCH 'SWI' WOULD BE CLOSED AND 'SWI2' OPENED.
```

```
C DESCRIPTION
```

```
C STARTING AT SWITCH 'SWI' THIS SUBROUTINE MOVES TO SWITCH 'SWI2'  
C CHECKING VIOLATIONS ON THE WAY.
```

```
C SEE ALSO
```

```
C FILES ROUTINES COMMENTS
```

```

C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON         ALL COMMON VARIABLES DECLARED
C      COMP           FILVOL      COMPUTES NODE VOLTAGES AND CABLE CURRENTS
C                                     ASSIGNS THE VIOLATION VECTOR
C      STEP2         MKESUM      COMPUTES THE IMPROVEMENT MEASURE SUM
C      STEP2         EVAL        EVALUATES THE MEASURE IN THE LOAD CURTAILMENT
C      STEP2         IMPROV      COMPUTES THE ESTIMATED REDUCTION IN THE
C                                     LOAD CURTAILMENT
C      MOVING        STEPUP      TAKES ONE STEP UP
C      MOVING        STEPDO      TAKES ONE STEP DOWN
C      MOVING        MOVE        TAKES ONE STEP UP OR DOWN
C      MOVING        DIREC       COMPUTES NEW DIRECTION

```

```

C      VARIABLES
C      SW1 IS THE OPEN SWITCH THAT WILL BE CLOSED DUE TO A LOAD TRANSFER
C      SW12 IS THE CLOSED SWITCH THAT WILL BE OPENED DUE TO A LOAD TRANSFER
C      SUM (OUT) IS THE ESTIMATED IMPROVEMENT
C      LOAD (OUT) IS THE POWER CONSUMPTION FOR THE THE NODES TRANSFERED VA.
C      ALPHA (OUT) IS THE NODE PRIORITY FOR TRANSFERED NODES.

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER SW1,SW12
REAL SUM,CUR,LOAD,ALPHA
INTEGER POI(2),OLD(2),P(2)
INTEGER DIR,BRA
REAL SQRT
LOGICAL ENDNOD,FIRST

```

```

INTRINSIC SQRT

```

```

ALPHA=100
POI(1)=SW1TYP
POI(2)=SW1
CALL STEPUP(POI,.FALSE.)
10 IF (POI(1).NE.SW1TYP) THEN
    CALL STEPUP(POI,.FALSE.)
    GOTO 10
ENDIF
DIR=DOWN
SW12=POI(2)

```

```

C      SW12 IS FOUND
C      ONE HAVE TO START FROM THIS SWITCH BECAUSE OF THE REDUCTION
C      DEPENDS ON POSSIBLE BRANCH POINTS. ONE SWITCHING OPERATION
C      CAN TRANSFER LARGE LOADS IN DIFFERENT BRANCHES

```

```

SUM=0.0
BRA=3
FIRST=.TRUE.
P(1)=POI(1)
P(2)=POI(2)
OLD(1)=0
OLD(2)=0
ENDNOD=.TRUE.
CALL STEPDO(POI,BRA,.FALSE.)
20 IF (POI(1).NE.P(1) .OR. POI(2).NE.P(2)) THEN

```

```

C      THIS LITTLE TREE WITH SW12 AS ROOT IS TRAVERSED COMPUTING
C      THE ESTIMATED REDUCTION OF THE VIOLATION VECTOR

```

```

IF (DIR.EQ.UP) THEN
    CALL MKESUM(POI,SUM,LOAD,ENDNOD,FIRST,ALPHA)

```

```

ELSE
  ENDNOD=.TRUE.
ENDIF
CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
CALL DIREC(POI,OLD,DIR,BRA)
GOTO 20
ENDIF
ENDNOD=.FALSE.
30 IF (POI(1).NE.FEETYP) THEN

C THE BIGGEST LINE CAPACITY VIOLATION CAN BE IN THIS CABLE

  CALL STEPUP(POI,.FALSE.)
  CALL MKESUM(POI,SUM,LOAD,ENDNOD,FIRST,ALPHA)
  GOTO 30
ENDIF
RETURN
END

```

```

SUBROUTINE MKESUM(POI,SUM,LOAD,ENDNOD,FIRST,ALPHA)
CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C AUTHOUR MAGNUS BERGSTRAND
C VERSION 1
C DATE 1989-10-05

```

```

C PURPOSE
C COMPUTES THE MEASURE SUM

```

```

C DESCRIPTION
C IF THE CURRENT POINT IS A NODE A CHECK IS DONE TO SEE IF
C THE NODE VOLTAGE IS BELOW THE MAXIMUM VOLTAGE DROP
C THEN A CHECK OF THE LINE CAPACITY CONSTRAINTS IS PERFORMED.
C IF ANY VIOLATIONS IS FOUND SUM WILL BE INCREASED.

```

```

C SEE ALSO

```

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
STEP2	IMP	COMPUTES THE ESTIMATED REDUCTION OF THE VIOLATION IN CASE OF FIRST OR SECOND STAGE SUPPORT BY A CALL OF MKESUM
STEP2	EVAL	COMPUTES THE ESTIMATED REDUCTION OF THE VIOLATION IN CASE OF LOAD CURTAILMENT BY A CALL OF IMPROV.
STEP2	IMPROV	COMPUTES THE REDUCTION OF THE VIOLATION VECTOR IN CASE OF A LOAD CURTAILMENT

```

C VARIABLES
C POI IS THE CURRENT POINT (TYPE,NR)
C SUM (OUT) IS THE ESTIMATED REDUCTION
C LOAD IS THE MAGNITUDE OF TRANSFERED LOAD IN VA.
C ENDNOD INDICATES THAT A NODE IS A LEAF
C FIRST INDICATES FIRST CALL FOR A FEEDER
C ALPHA IS THE PRIORITY

```

```

C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```
REAL SUM,CUR,SQRT,MIN,VOLSUM,LOAD,LOAD1,ALPHA
INTEGER POI(2),I
LOGICAL ENDNOD,FIRST
```

```
INTRINSIC SQRT,MIN
```

```
C CUR (REDUCTION IN LINE CAPACITY VIOLATIONS),VOLSUM (REDUCTION IN
C VOLTAGE DROP VIOLATIONS) AND LOAD1 (MAGNITUDE OF TRANSFERED LOAD)
C ARE SAVED BETWEEN CALLS
```

```
SAVE CUR
SAVE VOLSUM
SAVE LOAD1
```

```
IF (FIRST) THEN
```

```
C INITIATE
```

```
  CUR=0.0
  VOLSUM=0.0
  FIRST=.FALSE.
  LOAD1=0.0
ENDIF
```

```
C THE INDEX OF THE CABLE LEADING TOWARDS THE FEEDER IS STORED IN 'I'
```

```
IF (POI(1).EQ.SWITYP) THEN
  I=SWITCH(POI(2),2)
  I=SWITCH(POI(2),1)
ELSEIF (POI(1).EQ.NODTYP) THEN
  I=NODETO(POI(2),1)
  I=NODETO(POI(2),1)
  IF (ENDNOD) THEN
    LOAD1=LOAD1+SQRT(NODEPH(POI(2),1)**2+NODEPH(POI(2),2)**2)
    IF (VOLLIM.GT.NODEPH(POI(2),3)) THEN
```

```
C VOLTAGE DROP VIOLATION
```

```
      VOLSUM=VOLSUM+VOLLIM-NODEPH(POI(2),3)
      ALPHA=MIN(ALFA(POI(2)),ALPHA)
    ENDIF
  ENDNOD=.FALSE.
ENDIF
ELSEIF (POI(1).EQ.BRATYP) THEN
  I=BRAPOI(POI(2),2)
  I=BRAPOI(POI(2),1)
ENDIF
```

```
IF (CABPHY(I,4).GT.CABPHY(I,3)) THEN
```

```
C LINE CAPACITY VIOLATION
```

```
  IF ((CABPHY(I,4)-CABPHY(I,3)).GT.CUR) THEN
    CUR=CABPHY(I,4)-CABPHY(I,3)
  ENDIF
ENDIF
LOAD=LOAD1
SUM=SQRT((CUR**2)+(VOLSUM**2))
RETURN
END
```

```
LOGICAL FUNCTION VIOLTE(FEE)
```

```
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

C VERSION 1
 C DATE 1989-10-05
 C PURPOSE
 C CHECKS IF THE FEEDER 'FEE' IS VIOLATED
 C DESCRIPTION
 C IF THE FEEDER 'FEE' IS VIOLATED THEN
 C THIS LOGICAL FUNCTION EQUALS TRUE
 C VARIABLES
 C FEE IS THE NUMBER OF THE FEEDER
 C
 C CCC

INCLUDE 'CONSTANTS.FOR'
 INCLUDE 'COMMON.FOR'

VIOLTE=(VIOVEC(FEE,1).GT.0.0 .OR. VIOVEC(FEE,2).GT.0.0)
 RETURN
 END

SUBROUTINE STAGE1(FEE,OK,FIRST,OUTFEE)
 C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1
 C DATE 1989-10-05

C PURPOSE
 C TRANSFERS LOAD TO AN ADJACENT NON VIOLATED FEEDER.

C METHODE
 C THE METHODE IS DESCRIBED IN 'SERVICE RESTORATION IN ELECTRIC
 C DISTRIBUTIONAL NETWORKS' BY MAGNUS BERGSTRAND
 C THE METHOD WAS FIRST DESCRIBED BY AOKI IN
 C 'A NEW ALGORITHM FOR SERVICE RESRORATION IN DISTRIBUTION SYSTEMS'
 C IEEE/PES WINTER MEETING, NEW YORK 1989 PAPER 89 WM 085-2 PWRD

C DESCRIPTION
 C STARTING AT A VIOLATED FEEDER ALL POSSIBLE LOAD TRANSFERS TO
 C A NON VIOLATED FEEDER IS EXAMINED. THE LOAD TRANSFER THAT WILL
 C REDUCE THE VIOLATION VECTOR THE MOST WILL BE TRIED FIRST. IF THIS
 C LOAD TRANSFER CAUSES A VIOLATION IT IS ABANDONED AND THE NEXT LOAD
 C TRANSFER IS TRIED. IF NO LOAD TRANSFER IS POSSIBLE WITHOUT ANY NEW
 C VIOLATIONS OK EQUALS FALSE.
 C STAGE1 HAS ANOTHER FUNCTION. WHEN A FIRST STAGE SUPPORT HAS FAILED
 C A SECOND STAGE SUPPORT IS TRIED. FIRST IN THIS SECOND STAGE A CALL
 C OF STAGE1 IS MADE WITH OK.EQ..FALSE. . THIS CAUSES THE FIRST STAGE
 C TO ALLOW VIOLATIONS IN THE PRIMARY SUPPORT FEEDER.

C SEE ALSO

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
STEP2	STEP2	IS THE SECOND STEP OF THE ALGORITHM
STEP2	STAGE	HANDLES THE CALLS OF STAGE1 AND STAGE2
		MAKES A FILE CONTAINING DATA ABOUT THE
		DIFFERENT LOAD TRANSFERS.

```

C      STEP2          STAGE2          THE SECOND STAGE SUPPORT. TRANSFERS FIRST LOAD
C      TO A PRIMARY SUPPORT FEEDER ALLOWING VIOLATIONS
C      THESE VIOLATIONS ARE THEN REMOVED BY A LOAD
C      TRANSFER TO A SECONDARY SUPPORT FEEDER.

C      VARIABLES
C      FEE IS THE NUMBER OF VIOLATED FEEDER (IN)
C      IF OK IS FALSE A VIOLATION IN THE SUPPORT FEEDER IS ALLOWED
C      FIRST INDICATES IF IT IS THE FIRST TIME A LOAD TRANSFER IS TRIED.
C      IF A PREVIOUS LOAD TRANSFER HAS FAILED THE FILE CONTAINING DATA OF
C      POSSIBLE LOAD TRANSFERS IS KEPT.
C      OUTFEE IS AN OUTPARAMETER INDICATING THE NUMBER OF THE SWITCH NOW
C      VIOLATED DUE TO A LOAD TRANSFER. IT IS ONLY VALID WHEN THE CALL WAS
C      MADE WITH OK=FALSE.

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

LOGICAL FIRST
INTEGER FEE, TRAF0, OUTFEE
LOGICAL VIO, OK
INTEGER SWI, SWI2, FEEDER
INTEGER BSWI, BSWI2, BFEE
REAL SUM, BSUM, SLASK, VIOSUM, LOAD
REAL MAX, MAX2

SAVE MAX, MAX2

EXTERNAL VIOSUM

WRITE(*,*)'ENTERS STAGE 1', FEE
IF (OK) THEN
  CALL STAGE(FEE, OK)
  IF (OK) THEN
    WRITE(*,*)'READS IN STAGE1 *****'
    OPEN(1, FILE='STAGE', STATUS='OLD', ACCESS='SEQUENTIAL')
    BSUM=0.0
    MAX=1.0E+30
    BSWI=MAXNRS+1
10    READ(1, *, ERR=20) SWI, SWI2, SUM, FEEDER, LOAD, ALPHA

C      READS, COMPUTES AND DECIDES WHICH LOAD TRANSFER IS THE BEST
C      THIS LOAD TRANSFER IS SAVED

        SLASK=D.01+VIOSUM(FEE)-SUM
        SLASK=SUM/(SLASK+BETA)/LOAD/ALPHA
        WRITE(*,*)'IMPROVEMENT, FEEDER', SLASK, FEEDER
        IF (SLASK.GT.BSUM .AND. SLASK.LT.MAX) THEN
          BSUM=SLASK
          BSWI=SWI
          BSWI2=SWI2
          BFEE=FEEDER
        ENDIF
        GOTO 10
20    IF (BSWI.LT.MAXNRS+1) THEN

C      THEN A POSSIBLE LOAD TRANSFER WAS FOUND
C      THE NETWORKS STATE IS SAVED

        WRITE(*,*)'SAVENE*****'
        CALL SAVENE(BFEE)
        WRITE(*,*)FEE, BFEE
        WRITE(*,*)'SAVENE*****'
        CALL SAVENE(FEE)

```

C THE LOAD TRANSFER IS PERFORMED

```
SWITCH(BSWI2,1)=OPEN
SWITCH(BSWI,1)=CLOSED
SWINR(BSWI)=3
CHGSWI(BSWI2)=.NOT. CHGSWI(BSWI2)
SWINR(BSWI2)=3
CHGSWI(BSWI)=.NOT. CHGSWI(BSWI)
TRA=MSFEED(BFEE,1)
CALL FORBAC(BSWI,TRA,BFEE)
```

C THE NEW STATE IS COMPUTED

```
CALL COMPIM(BFEE)
IF (SWITCH(BSWI2,2).EQ.3) THEN
  SWITCH(BSWI2,2)=4
ELSE
  SWITCH(BSWI2,2)=3
ENDIF
CALL COMPIM(FEE)
CALL FILPOW(TRA)
TRA=MSFEED(FEE,1)
CALL FILPOW(TRA)
CALL FILVOL(FEE,VIO)
CALL FILVOL(BFEE,VIO)
IF (VIO) THEN
```

C SINCE A VIOLATION OCCURED THE OLD STATE MUST BE READ BACK

```
REWIND 1
WRITE(*,*)'CALLS RESNET*****'
REWIND 1
CALL RESNET
WRITE(*,*)'CALLS RESNET*****'
CALL RESNET
```

C TAKE NEXT LOAD TRANSFER AND TRY AGAIN

```
MAX=BSUM
BSUM=0.0
BSWI=MAXNRS+1
GOTO 10
ELSEIF (.NOT. VIO) THEN
```

C DELETE FILE AFTER A SUCCESSFUL LOAD TRANSFER

```
CLOSE(1,STATUS='DELETE')
WRITE(*,*)'LEAVES STAGE1 AFTER SUCCESSFULL SESSION,OK=T'
ENDIF
ELSE
```

C NO POSSIBLE LOAD TRANSFER WAS FOUND.

C SAVE FILE FOR LATER USE IN THE SECOND STAGE SUPPORT

```
CLOSE(1,STATUS='KEEP')
WRITE(*,*)'OK=.FALSE.'
OK=.FALSE.
ENDIF
ENDIF
ELSEIF (.NOT. OK) THEN
```

C ALLOWS VIOLATION IN THE SUPPORT FEEDER

```
WRITE(*,*)'OK IS FALSE'
IF (FIRST) THEN
```

```
BSUM=0.0
MAX2=1.0E+30
ELSE
```

```
C A PREVIOUS UNSUCCESSFUL SECOND STAGE HAS BEEN MADE.
C THE OLD STATE MUST BE READ BACK
```

```
CALL RESNET
CALL RESNET
ENDIF
BSWI=MAXNRS+1
```

```
C READ, COMPUTE AND IDENTIFY THE BEST LOAD TRANSFER
```

```
25 OPEN(1, FILE='STAGE', STATUS='OLD', ACCESS='SEQUENTIAL')
READ(1, *, ERR=30) SWI, SWI2, SUM, FEEDER, LOAD, ALPHA
SLASK=0.01+VIOSUM(FEE)-SUM
SLASK=SUM/(SLASK+BETA)/LOAD/ALPHA
IF (SLASK.GT.BSUM .AND. SLASK.LT.MAX2) THEN
BSUM=SLASK
BSWI=SWI
BSWI2=SWI2
BFEE=FEEDER
ENDIF
GOTO 25
```

```
30 IF (BSWI.LT.MAXNRS+1) THEN
```

```
C A POSSIBLE LOAD TRANSFER WAS FOUND
C SAVE OLD STATE
```

```
WRITE(*,*) 'SAVENE*****'
CALL SAVENE(BFEE)
WRITE(*,*) FEE, BFEE
WRITE(*,*) 'SAVENE*****'
CALL SAVENE(FEE)
```

```
C PERFORM LOAD TRANSFER
```

```
SWITCH(BSWI2,1)=OPEN
SWITCH(BSWI2,5)=0
SWITCH(BSWI,1)=CLOSED
SWINR(BSWI2)=3
SWINR(BSWI)=3
CHGSWI(BSWI2)=.NOT. CHGSWI(BSWI2)
CHGSWI(BSWI)=.NOT. CHGSWI(BSWI)
TRA=MSFEED(BFEE,1)
CALL FORBAC(BSWI,TRA,BFEE)
```

```
C COMPUTE NEW STATE
```

```
CALL COMPIM(BFEE)
CALL COMPIM(FEE)
CALL FILPOW(TRA)
TRA=MSFEED(FEE,1)
CALL FILPOW(TRA)
CALL FILVOL(FEE,VIO)
CALL FILVOL(BFEE,VIO)
OK=.TRUE.
MAX2=BSUM
CLOSE(1, STATUS='KEEP')
OUTFEE=BFEE
```

```
ELSE
```

```
C NO POSSIBLE LOAD TRANSFER WAS FOUND
```

```
WRITE(*,*) 'STAGE1 WITH OK=F HAD AN UNSUCCESSFULL COMPLETION'
```



```

        OK=.FALSE.
        CLOSE(1,STATUS='DELETE')
    ENDIF
ENDIF
RETURN
END

```

```

SUBROUTINE SAVENE(FEEDER)
CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

AUTHOUR MAGNUS BERGSTRAND

```

```

VERSION 1

```

```

DATE 1989-10-05

```

```

PURPOSE
SAVES THE FORMER STATE IN THE NETWORK WHEN A LOAD TRANSFER IS
TO BE MADE

```

```

DESCRIPTION
TRAVERSES A TREE WITH THE FEEDER 'FEEDER' AS ROOT SAVING EVERYTHING
IN THE NETWORK TOPOLOGY THAT COULD BE CHANGED DUE TO A LOAD TRANSFER
IN THE FILE 'HISTORY.DAT'
A STACK OF BACKUP FILES IS CREATED

```

```

SEE ALSO

```

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
MOVING	STEPUP	TAKES ONE STEP UP
MOVING	STEPDO	TAKES ONE STEP DOWN
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
STEP2	RESNET	READS BACK THE FORMER STATE FROM THE FILE 'HISTORY.DAT'
STEP2	STAGE1	CREATES A STACK OF BACKUP FILES BY CALLS OF SAVENET.

```

VARIABLES
'FEEDER' IS THE NUMBER OF THE FEEDER THAT WILL BE SAVED.

```

```

CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

```

```

INTEGER FEEDER
INTEGER POI(2),OLD(2)
INTEGER I,J
INTEGER DIR,BRA
LOGICAL READY

```

```

OPEN(2,FILE='HISTORY',STATUS='NEW',ACCESS='SEQUENTIAL')
POI(1)=FEETYP
POI(2)=FEEDER
WRITE(2,*)FEEDER,VIOVEC(FEEDER,1),VIOVEC(FEEDER,2)
DIR=DOWN
BRA=3
OLD(1)=POI(1)
OLD(2)=POI(2)
READY=.FALSE.

```

```

10 IF (.NOT. READY) THEN
C   TRAVERSES THE TREE

    CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
    CALL DIREC(POI,OLD,DIR,BRA)
    IF (DIR.EQ.UP) THEN

C   FIRST IS THE POINT WRITTEN.
C   THEN DIFFERENT DATA THAT MAY BE CHANGED ARE WRITTEN DEPENDING
C   ON THE TYPE OF THE POINT

        WRITE(2,*)POI(1),POI(2)
        IF (POI(1).EQ.SWITYP) THEN
            WRITE(2,*)SWITCH(POI(2),1),SWITCH(POI(2),2),SWITCH(POI(2),5)
            I=SWITCH(POI(2),2)
            I=SWITCH(POI(2),I)
        ELSEIF (POI(1).EQ.NODTYP) THEN
            WRITE(2,*)NODETO(POI(2),1),NODETO(POI(2),4)
            WRITE(2,*)NODEPH(POI(2),1),NODEPH(POI(2),2),NODEPH(POI(2),3)
            I=NODETO(POI(2),1)
            I=NODETO(POI(2),I)
        ELSEIF (POI(1).EQ.BRATYP) THEN
            WRITE(2,*) BRAPOI(POI(2),2)
            DO 20 J=3,(2+BRAPOI(POI(2),1))
                WRITE(2,*)BRAPH(POI(2),J,1),BRAPH(POI(2),J,2)
20            CONTINUE
            WRITE(2,*)BRAPOI(POI(2),2)
            I=BRAPOI(POI(2),2)
            I=BRAPOI(POI(2),I)
        ELSEIF (POI(1).EQ.FEETYP) THEN
            I=MSFEED(POI(2),2)
        ELSEIF (POI(1).EQ.TRATYP) THEN
            WRITE(2,*)TRAFOS(POI(2),3),TRAFOS(POI(2),4)
            WRITE(2,*)TRAFOS(POI(2),5),TRAFOS(POI(2),6)
            WRITE(2,*)TRAFOS(POI(2),7)
        ELSE
            WRITE(*,*)'      WRONG TYPE IN SAVENE'
            STOP
        ENDIF
        IF (POI(1).NE.TRATYP) THEN
            WRITE(2,*)CABPHY(I,4)
        ENDIF
    ENDIF
    READY=(POI(1).EQ.FEETYP)
    GOTO 10
ENDIF
DIR=UP
READY=.FALSE.
30 IF (.NOT. READY) THEN
C   EVEN THE TOPOLOGY ABOVE THE FEEDER IS SAVED

    CALL STEPUP(POI,.FALSE.)
    IF (DIR.EQ.UP) THEN
        WRITE(2,*)POI(1),POI(2)
        IF (POI(1).EQ.SWITYP) THEN
            WRITE(2,*)SWITCH(POI(2),1),SWITCH(POI(2),2),SWITCH(POI(2),5)
            I=SWITCH(POI(2),2)
            I=SWITCH(POI(2),I)
        ELSEIF (POI(1).EQ.BRATYP) THEN
            WRITE(2,*) BRAPOI(POI(2),2)
            DO 40 J=3,(2+BRAPOI(POI(2),1))
                WRITE(2,*)BRAPH(POI(2),J,1),BRAPH(POI(2),J,2)
40            CONTINUE
            WRITE(2,*)BRAPOI(POI(2),2)

```

```

        I=BRAPOI(POI(2),2)
        I=BRAPOI(POI(2),I)
    ELSEIF (POI(1).EQ.FEETYP) THEN
        I=MSFEED(POI(2),2)
    ELSEIF (POI(1).EQ.TRATYP) THEN
        WRITE(2,*)TRAFOS(POI(2),3),TRAFOS(POI(2),4)
        WRITE(2,*)TRAFOS(POI(2),5),TRAFOS(POI(2),6)
        WRITE(2,*)TRAFOS(POI(2),7)
    ELSE
        WRITE(*,*)'        WRONG TYPE IN SAVENE'
        STOP
    ENDIF
    IF (POI(1).NE.TRATYP) THEN
        WRITE(2,*)CABPHY(I,4)
    ENDIF
ENDIF
READY=POI(1).EQ.TRATYP
GOTO 30
ENDIF
CLOSE(2,STATUS='KEEP')
RETURN
END

```

SUBROUTINE RESNET

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-05

C PURPOSE

C READS BACK A FORMER STATE WHEN A LOAD TRANSFER HAS FAILED.

C DESCRIPTION

C OPENS AND READS THE FILE 'HISTORY.DAT' (ALWAYS THE TOP OF THE STACK
C OF BACKUP FILES)

C SEE ALSO

C FILES

C ROUTINES

C COMMENTS

C CONSTANTS

C ALL CONSTANTS DECLARED

C COMMON

C ALL COMMON VARIABLES DECLARED

C STEP2

C SAVENE

C SAVES A FORMER STATE IN THE FILE 'HISTORY.DAT'

C CCC

INCLUDE 'CONSTANTS.FOR'

INCLUDE 'COMMON.FOR'

INTEGER POI(2),I,J

OPEN(2,FILE='HISTORY',STATUS='OLD',ACCESS='SEQUENTIAL')

READ(2,*)I,VIOVEC(I,1),VIOVEC(I,2)

10 READ(2,*,ERR=40)POI(1),POI(2)

IF (POI(1).EQ.SWITYP) THEN

 J=SWITCH(POI(2),1)

 READ(2,*)SWITCH(POI(2),1),SWITCH(POI(2),2),SWITCH(POI(2),5)

 I=SWITCH(POI(2),2)

 I=SWITCH(POI(2),I)

 IF (J.NE.SWITCH(POI(2),1)) THEN

 CHGSI(POI(2))=.NOT. CHGSI(POI(2))

 ENDIF

ELSEIF (POI(1).EQ.NODTYP) THEN

```

      READ(2,*)NODETO(POI(2),1),NODETO(POI(2),4)
      READ(2,*)NODEPH(POI(2),1),NODEPH(POI(2),2),NODEPH(POI(2),3)
      I=NODETO(POI(2),1)
      I=NODETO(POI(2),I)
ELSEIF (POI(1).EQ.BRATYP) THEN
      READ(2,*) BRAPOI(POI(2),2)
      J=BRAPOI(POI(2),2)
      J=BRAPOI(POI(2),J)
      I=J
      DO 20 J=3,(2+BRAPOI(POI(2),1))
        READ(2,*)BRAPH(POI(2),J,1),BRAPH(POI(2),J,2)
20      CONTINUE
      READ(2,*)BRAPOI(POI(2),2)
      I=BRAPOI(POI(2),2)
      I=BRAPOI(POI(2),I)
ELSEIF (POI(1).EQ.FEETYP) THEN
      I=MSFEED(POI(2),2)
ELSEIF (POI(1).EQ.TRATYP) THEN
      READ(2,*)TRAFOS(POI(2),3),TRAFOS(POI(2),4)
      READ(2,*)TRAFOS(POI(2),5),TRAFOS(POI(2),6)
      READ(2,*)TRAFOS(POI(2),7)
ELSE
      WRITE(*,*)'      WRONG TYPE IN RESNET'
      STOP
ENDIF
IF (POI(1).NE.TRATYP) THEN
      READ(2,*)CABPHY(I,4)
ENDIF
GOTO 10
40  CLOSE(2,STATUS='DELETE')
      RETURN
      END

```

```

SUBROUTINE STAGE2(FEE,OK,FIRST,OUTFEE,FEE2,I)
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C      AUTHOUR MAGNUS BERGSTRAND
C
C      VERSION 1
C
C      DATE 1989-10-05
C
C      PURPOSE
C      TRANSFERS LOAD TO A SECONDARY SUPPORT FEEDER VIA A PRIMARY SUPPORT
C      FEEDER.
C
C      METHOD
C      THE METHOD IS DESCRIBED IN 'SERVICE RESTORATION IN ELECTRIC
C      DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND.
C      THE METHOD WAS FIRST DESCRIBED BY AOKI IN
C      'A NEW ALGORITHM FOR SERVICE RESRORATION IN DISTRIBUTION SYSTEMS'
C      IEEE/PES WINTER MEETING,NEW YORK 1989 PAPER 89 WM 085-2 PWRD
C
C      DESCRIPTION
C      LOAD IS TRANSFERED TO A PRIMARY SUPPORT FEEDER CAUSING A VIOLATION.
C      THIS VIOLATION IS THEN REMOVED BY TRANSFERED LOAD TO A SECONDARY
C      SUPPORT FEEDER
C
C      SEE ALSO
C
C      FILES      ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON      ALL COMMON VARIABLES DECLARED
C      STEP2      STEP2      PROVIDES THE FRAMEWORK TO STAGE1 AND
C      STAGE2

```

```

C     STEP2          STAGE1      FIRST STAGE SUPPORT. CAN BE FORCED TO ALLOW
C     STEP2          DELHIS      VIOLATION IN THE SUPPORT FEEDER.
C     STEP2          DELHIS      READS BACK THE FORMER STATE

C     VARIABLES
C     FEE IS THE NUMBER OF THE VIOLATED FEEDER
C     OK=FALSE INDICATES THAT VIOLATION IS ALLOWED IN THE SECONDARY SUPPORT
C     FEEDER. (THE ALGORITHM CAN BE GENERALIZED WITH A THIRD,FOURTH...
C     STAGE SUPPORT)
C     FIRST INDICATES FIRST CALL FOR FEEDER 'FEE'
C     OUTFEE (OUTPARAMETER) IS THE NUMBER OF THE VIOLATED PRIMARY
C     SUPPORT FEEDER. IT IS ONLY VALID WHEN OK=F.
C     FEE2 (OUTPARAMETER) IS THE NUMBER OF THE VIOLATED SECONDARY
C     SUPPORT FEEDER. IT IS ONLY VALID WHEN OK=F.
C     'I' IS THE NUMBER OF SUCSESFUL CALLS OF FIRST STAGE SUPPORT.
C     2*I IS THE NUMBER OF BACKUP FILES IN THE STACK.

C     CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
C
C     INCLUDE 'CONSTANTS.FOR'
C     INCLUDE 'COMMON.FOR'

C     INTEGER FEE,OUTFEE,I
C     LOGICAL OK,FIRST
C     INTEGER FEE2,SLASK
C     LOGICAL READY,FIRST1,FIRST2,VIOLTE

C     EXTERNAL VIOLTE

C     WRITE(*,*)'ENTERS STAGE 2'
C     IF (OK) THEN

C     NO VIOLATION IS ALLOWED IN THE SECONDARY SUPPORT FEEDER

C     FIRST1=.TRUE.
C     WRITE(*,*)'STAGE2 CALLS STAGE1 *****'

C     ALLOWS VIOLATION IN THE PRIMARY SUPPORT FEEDER

C     OK=.FALSE.
C     CALL STAGE1(FEE,OK,FIRST1,FEE2)
C     FIRST2=.FALSE.
C     FIRST1=.FALSE.
C     READY=.NOT. OK
C     I=0
10    IF (.NOT. READY) THEN

C     DURING THIS LOOP THE VIOLATION IN THE PRIMARY SUPPORT FEEDER
C     IS TRIED TO BE ELIMINATED
C     THE VARIABEL I IS INCREMENTED FOR EACH SUCSESFUL LOAD TRANSFER
C     TO A SECONDARY SUPPORT FEEDER.

C     WRITE(*,*)'LOOP IN STAGE 2 *****'
C     OK=.TRUE.
C     CALL STAGE1(FEE2,OK,FIRST2,SLASK)
C     IF (.NOT. OK) THEN

C     IT WAS NOT POSSIBLE TO TRANSFER LOAD FROM THE PRIMARY SUPPORT FEEDER
C     TO ANY SECONDARY SUPPORT FEEDER
C     READ BACK THE FORMER STATE
C     TRY A NEW PRIMARY SUPPORT FEEDER

C     CALL DELHIS(I)
C     I=0
C     OPEN(1,FILE='STAGE',STATUS='OLD',ACCESS='SEQUENTIAL')
C     CLOSE(1,STATUS='DELETE')

```

```

        CALL STAGE1(FEE,OK,FIRST1,FEE2)
        READY=(.NOT. OK)
    ELSE
        READY=.NOT. VIOLTE(FEE2)
        I=I+1
    ENDIF
    GOTO 10
ENDIF
ELSEIF (FIRST) THEN
C      OK IS FALSE. VIOLATIONS ARE ALLOWED IN THE SECONDARY SUPPORT FEEDER

        FIRST=.FALSE.
        FIRST1=.TRUE.
        OK=.FALSE.
        CALL STAGE1(FEE,OK,FIRST1,FEE2)
C      VIOLATIONS ARE ALLOWED IN THE PRIMARY SUPPORT FEEDER

        FIRST2=.TRUE.
        FIRST1=.FALSE.
        READY=(.NOT. OK)
        I=0
40     IF (.NOT. READY) THEN
            WRITE(*,*)'LOOP IN STAGE 2 *****'
            OK=.TRUE.
            CALL STAGE1(FEE2,OK,FIRST2,OUTFEE)
C      VIOLATIONS IN THE SECONDARY SUPPORT FEEDER WAS ALLOWED

            IF (.NOT. OK) THEN
C              NO POSSIBLE LOAD TRANSFER TO ANY SECONDARY SUPPORT FEEDER WAS FOUND
C              READ BACK FORMER STATE
C              FIND A NEW POSSIBLE PRIMARY SUPPORT FEEDER

                READY=.TRUE.
                OK=.FALSE.
                CALL STAGE1(FEE2,OK,FIRST2,OUTFEE)
                IF (.NOT. OK) THEN
                    CALL DELHIS(I)
                    I=0
                    OPEN(1,FILE='STAGE',STATUS='OLD',ACCESS='SEQUENTIAL')
                    CLOSE(1,STATUS='DELETE')
                    CALL STAGE1(FEE,OK,FIRST1,FEE2)
C      A NEW VIOLATED SECONDARY SUPPORT FEEDER WAS TRIED

                        READY=(.NOT. OK)
                ELSE
                    I=I+1
                ENDIF
            ELSE
C      THE ROUTINE MANAGED TO TRANSFER LOAD TO A NOW VIOLATED SECONDARY
C      SUPPORT FEEDER

                READY=.NOT. VIOLTE(FEE2)
                I=I+1
            ENDIF
            GOTO 40
        ENDIF
    ELSE
C      THERE HAVE BEEN AN UNSUCCESSFUL THIRD STAGE SUPPORT
C      TRY TO TRANSFER LOAD TO A NEW SECONDARY SUPPORT FEEDER

```

```

FIRST2=.FALSE.
FIRST1=.FALSE.
READY=.FALSE.
I=0
50 IF (.NOT. READY) THEN
    OK=.FALSE.
    CALL STAGE1(FEE,OK,FIRST2,SLASK)

C     LOAD WAS TRANSFERED TO A PRIMARY SUPPORT FEEDER

    FIRST2=.FALSE.
    IF (.NOT. OK) THEN

C         THE LOAD TRANSFER TO ANY PRIMARY SUPPORT FEEDER WAS
C         UNSUCCESSFUL.

        READY=.TRUE.
    ELSE
        CALL STAGE1(SLASK,OK,FIRST1,FEE2)
        IF (OK) THEN

C             LOAD WAS TRANSFERED TO THE SECONDARY SUPPORT FEEDER

            READY=.TRUE.
        ELSE

C             LOAD COULD NOT BE TRANSFERED TO ANY SECONDARY SUPPORT FEEDER
C             READ BACK OLD STATE AND TRY AGAIN

                CALL DELHIS(I)
                I=0
            ENDIF
            GOTO 50
        ENDIF
    ENDIF
ENDIF
WRITE(*,*)'EXIT STAGE 2 OK=',OK
RETURN
END

```

```

SUBROUTINE STAGE3(FEE,OK,FIRST,OUTFEE)
CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C     AUTHOUR MAGNUS BERGSTRAND

```

```

C     VERSION 1

```

```

C     DATE 1989-10-05

```

```

C     PURPOSE
C     TRANSFERS LOAD TO TERTIARY SUPPORT FEEDER

```

```

C     DESCRIPTION
C     FIRST A FIRST STAGE SUPPORT IS TRIED ALLOWING VIOLATIONS
C     IN THE SUPPORT FEEDER. THEN IS SECOND STAGE SUPPORT APPLIED
C     TO THIS VIOLATED SUPPORT FEEDER.

```

```

C     SEE ALSO

```

```

C     FILES          ROUTINES          COMMENTS
C     CONSTANTS     ALL CONSTANTS DECLARED
C     COMMON         ALL COMMON VARIABLES DECLARED

```

```

C      STEP2          STAGE1          FIRST STAGE SUPPORT
C      STEP2          STAGE2          SECOND STAGE SUPPORT

C      VARIABLES
C      FEE IS THE NUMBER OF THE VIOLATED FEEDER
C      OK INDICATES THAT NO VIOLATION IS ALLOWED IN THE TERTIARY SUPPORT
C      FEEDER
C      FIRST INDICATES WETHER A FORMER CALL WITH FEEDER=FEE HAS BEEN MADE
C      OUTFEE IS THE NUMBER OF THE TERTIARY SUPPORT FEEDER
C      THESE LAST VARIABLES ARE WITHOUT MEANING IF NOT A FOURTH STAGE
C      SUPPORT IS IMPLEMENTED

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER FEE,OUTFEE
LOGICAL OK,FIRST

IF (OK) THEN
  FIRST1=.TRUE.
  OK=.FALSE.
  CALL STAGE2(FEE,OK,FIRST1,OUTF,FEE2,NR)
10  IF (.NOT. READY) THEN
    CALL STAGE1(OUTF,OK,.TRUE.,SLASK)
    IF (.NOT. OK) THEN
      CALL DELHIS(NR)
      NR=0
      OPEN(1,FILE='STAGE',STATUS='OLD',ACCESS='SEQUENTIAL')
      CLOSE(1,STATUS='DELETE')
      CALL STAGE2(FEE,FIRST1,OUTF,FEE2,NR)
      READY=(.NOT. OK)
    ENDIF
    GOTO 10
  ENDIF
ENDIF
RETURN
END

SUBROUTINE DELHIS(I)
C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

C      AUTHOUR MAGNUS BERGSTRAND

C      VERSION 1

C      DATE 1989-10-05

C      PURPOSE
C      READS BACK THE STATE OF NETWORK AS IT WAS BEFORE AN UNSUCCESSFULL
C      LOAD TRANSFER

C      SEE ALSO

C      FILES          ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON          ALL COMMON VARIABLES DECLARED
C      STEP2          RESNET         READS THE FILE 'HISTORY.DAT' CONTAINING
C      PARTS OF THE STATE

C      VARIABLES
C      'I' INDICATES THE SIZE OF THE STACK OF BACK UP FILES.
C      2*I IS THE NUMBER OF FILES

```


C CCC

INTEGER I,J

DO 10 J=1,I
CALL RESNET
CALL RESNET

10 CONTINUE
RETURN
END

SUBROUTINE STEP2(VIO)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-05

C PURPOSE
C PROVIDES THE FRAMEWORK FOR THE FIRST AND SECOND STAGE SUPPORT

C METHOD
C THE ALGORITHM IS DESCRIBED IN 'SERVICE RESTORATION IN ELECTRIC
C DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND
C THE METHOD WAS FIRST DESCRIBED BY AOKI IN
C 'A NEW ALGORITHM FOR SERVICE RESRORATION IN DISTRIBUTION SYSTEMS'
C IEEE/PES WINTER MEETING,NEW YORK 1989 PAPER 89 WM 085-2 PWRD

C DESCRIPTION
C ALL FEEDERS ARE CHECKED IN ORDER IF THEY ARE VIOLATED
C FOR ANY VIOLATED FEEDER:
C FIRST STAGE SUPPORT
C IF FAILURE: SECOND STAGE SUPPORT
C IF SUCCES REPEAT IT

C SEE ALSO

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
STEP2	STAGE1	FIRST STAGE SUPPORT
STEP2	STAGE2	SECOND STAGE SUPPORT
STEP2	VIOLTE	LOGICAL FUNCTION. CHECKS IF A FEEDER IS VIOLATED

C VARIABLES
C VIO IS AN OUTPARAMETER. IT INDICATES WHETHER ANY FEEDER IS VIOLATED.

C CCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

LOGICAL OK,VIO
INTEGER FEE,OUTFEE
INTEGER I,NR,J
LOGICAL VIOLTE

EXTERNAL VIOLTE

VIO=.FALSE.
DO 10 I=1,MAXNRF

```
FEE=I
IF (VIOLTE(FEE)) THEN
```

```
C      A VIOLATED FEEDER IS FOUND
```

```
      OK=.TRUE.
20     IF (OK) THEN
          CALL STAGE1(FEE,OK,.TRUE.,OUTFEE)
          IF (.NOT. OK) THEN
```

```
C      FIRST STAGE SUPPORT FAILED
C      TRY A SECOND STAGE SUPPORT
```

```
      OK=.TRUE.
      CALL STAGE2(FEE,OK,.TRUE.,OUTFEE,NR)
      IF (OK) THEN
        DO 30 J=1,2*NR
```

```
C      REMOVE USED FILES
```

```
        OPEN(3,FILE='HISTORY',STATUS='OLD',ACCESS='SEQUENTIAL')
        CLOSE(3,STATUS='DELETE')
```

```
30     CONTINUE
        OPEN(3,FILE='STAGE',STATUS='OLD',ACCESS='SEQUENTIAL')
        CLOSE(3,STATUS='DELETE')
      ENDIF
    ELSE
```

```
C      REMOVE USED FILES
```

```
      OPEN(1,FILE='HISTORY',STATUS='OLD',ACCESS='SEQUENTIAL')
      CLOSE(1,STATUS='DELETE')
      OPEN(1,FILE='HISTORY',STATUS='OLD',ACCESS='SEQUENTIAL')
      CLOSE(1,STATUS='DELETE')
```

```
      ENDIF
      OK=(OK .AND. VIOLTE(FEE))
      GOTO 20
```

```
    ENDIF
```

```
  ENDIF
```

```
  VIO=VIO .OR. VIOLTE(FEE)
```

```
10 CONTINUE
RETURN
END
```

```
SUBROUTINE IMPROV(FEE,SWI)
```

```
C CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

```
C VERSION 1
```

```
C DATE 1989-10-05
```

```
C PURPOSE
```

```
C EXAMINES DIFFERENT POSSIBLE SWITCHES SUITABLE FOR A LOAD TRANSFER
C THE BEST SWITCH, THAT IS THE ONE THAT WILL CAUSE A SMALL REDUCTION
C OF THE VIOLATION VECTOR WILL BE CHOSEN
```

```
C DESCRIPTION
```

```
C TRAVERSES THE TREE WITH FEEDER 'FEE' AS ROOT.
C WHEN A LEAF IS FOUND A CALL OF EVAL IS PERFORMED.
```

```
C SEE ALSO
```

```

C      FILES          ROUTINES      COMMENTS
C      CONSTANTS      ALL CONSTANTS DECLARED
C      COMMON         ALL COMMON VARIABLES DECLARED
C      MOVING         MOVE          TAKES ONE STEP UP OR DOWN
C      MOVING         DIREC        COMPUTES NEW DIRECTION
C      STEP2          EVAL         EVALUATES THE NUMBER OF SWITCH THAT IS
C                                  THE MOST SUITABLE TO PERFORM A LOAD CURTAILMENT

```

```

C      VARIABLES
C      'FEE' IS THE NUMBER OF THE VIOLATED FEEDER CAUSING THE LOAD CURTAILMENT
C      SWI (OUTPARAMETER) IS THE NUMBER OF THE SWITCH TO OPEN FOR THE
C      BEST LOAD CURTAILMENT

```

```

C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

      INCLUDE 'CONSTANTS.FOR'
      INCLUDE 'COMMON.FOR'

```

```

      INTEGER FEE,SWI,P(2),OLD(2)
      INTEGER DIR,BRA
      LOGICAL READY,ENDS,FIRST

```

```

      P(1)=FEETYP
      P(2)=FEE
      DIR=DOWN
      BRA=3
      FIRST=.TRUE.
      READY=.FALSE.

```

```

10     IF (.NOT. READY) THEN
           IF (DIR.EQ.DOWN) THEN
               ENDS=.TRUE.
           ENDIF
           CALL MOVE(P,OLD,DIR,BRA,.FALSE.)
           CALL DIREC(P,OLD,DIR,BRA)
           IF (DIR.EQ.UP .AND. ENDS) THEN
               CALL EVAL(P,SWI,FIRST,FEE)
               ENDS=.FALSE.
           ENDIF
           READY=P(1).EQ.FEETYP
           GOTO 10
       ENDIF
       RETURN
       END

```

```

      SUBROUTINE EVAL(P,SWI,FIRST,FEE)
      C      CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

```

```

C      AUTHOUR MAGNUS BERGSTRAND

```

```

C      VERSION 1

```

```

C      DATE 1989-10-05

```

```

C      PURPOSE
C      COMPUTES THE NUMBER OF THE SWITCH MOST SUITABLE FOR A LOAD
C      CURTAILMENT SO FAR

```

```

C      DESCRIPTION
C      STARTING FROM A LEAF THE SUBROUTINE FINDS THE FIRST SWITCH
C      TOWARDS THE TOP. WHEN A SWITCH IS FOUND THIS IS USED AS A ROOT IN
C      A TREE. THE TREE IS TRAVERSED AND IN THE MEAN TIME THE SUBROUTINE
C      COMPUTES THE MAGNITUDE OF THE WITHDRAWAL OF THE VIOLATION VECTOR.
C      BEGINING AT THE ROOT IS NECESSARY DUE TO THE EXISTANCE OF BRANCH POINTS
C      LOAD IN MORE THAN ONE BRANCH MAY BE CUT OFF
C      THE BEST MAGNITUDE OF VIOLATIONS WITHDRAWAL AND THE

```

C NUMBER OF THE SWITCH CAUSING THE WITHDRAWAL ARE SAVED
C BETWEEN CALLS

C SEE ALSO

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
MOVING	STEPUP	TAKES ONE STEP UP
MOVING	STEPDO	TAKES ONE STEP DOWN
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
STEP2	MKESUM	DECIDES THE NUMBER OF THE SWITCH DUE TO A STAGE SUPPORT

C VARIABLES

C P IS THE LEAF (TYPE, NR)

C SWI (OUTPARAMETER) IS THE NUMBER OF THE SWITCH SO FAR

C FIRST INDICATES FIRST CALL IN NEW STATE

C FEE IS THE NUMBER OF THE VIOLATED FEEDER

C CCC

INCLUDE 'CONSTANTS.FOR'

INCLUDE 'COMMON.FOR'

INTEGER POI(2), P(2), SWI, FEE, TEMP, DIR, BRA, OLD(2)

LOGICAL FIRST, READY

REAL BEST, VOLSUM, SUM, CURSUM, LOAD, SLASK

REAL VIOSUM, SQRT, Z, MIN, ALPHA

REAL CUR1

INTRINSIC SQRT, MIN

EXTERNAL VIOSUM

SAVE TEMP

SAVE BEST

IF (FIRST) THEN

BEST=1.0E+30

FIRST=.FALSE.

TEMP=0

ENDIF

VOLSUM=0.0

CURSUM=0.0

CUR1=0.0

LOAD=0.0

POI(1)=P(1)

POI(2)=P(2)

READY=.FALSE.

5 IF (.NOT. READY) THEN

CALL STEPUP(POI, .FALSE.)

READY=POI(1).EQ.SWITYP

GOTO 5

ENDIF

C THE ROOT IS FOUND

TEMP=POI(2)

DIR=DOWN

BRA=3

READY=.FALSE.

10 IF (.NOT. READY) THEN

IF (DIR.EQ.UP) THEN

C THE NUMBER OF A CABLE IS STORED

```

IF (POI(1).EQ.SWITYP) THEN
  I=SWITCH(POI(2),2)
  I=SWITCH(POI(2),I)
ELSEIF (POI(1).EQ.NODTYP) THEN
  I=NODETO(POI(2),1)
  I=NODETO(POI(2),I)

```

C NODE PRIORITY,CURRENT REDUCTION,LOAD REDUCTION AND
C REDUCTION IN VOLTAGE DROP IS COMPUTED

```

ALPHA=ALFA(POI(2))
Z=SQRT((NODEPH(POI(2),4)**2)+(NODEPH(POI(2),5)**2))
CUR1=CUR1+NODEPH(POI(2),3)/Z
LOAD=LOAD+SQRT(NODEPH(POI(2),1)**2+NODEPH(POI(2),2)**2)
IF (VOLLIM.GT.NODEPH(POI(2),3)) THEN
  VOLSUM=VOLSUM+VOLLIM-NODEPH(POI(2),3)
ENDIF
ELSEIF (POI(1).EQ.BRATYP) THEN
  I=BRAPOI(POI(2),2)
  I=BRAPOI(POI(2),I)
ENDIF
IF (CABPHY(I,4).GT.CABPHY(I,3)) THEN

```

C CHECK OF LINE CAPACITY CONSTRAINTS

```

IF ((CABPHY(I,4)-CABPHY(I,3)).GT.CURSUM) THEN
  CURSUM=MIN(CUR1,(CABPHY(I,4)-CABPHY(I,3)))
ENDIF
ENDIF
ENDIF

```

C TRAVERSES THE TREE

```

CALL MOVE(POI,OLD,DIR,BRA,.FALSE.)
CALL DIREC(POI,OLD,DIR,BRA)
IF (POI(1).EQ.SWITYP .AND. POI(2).EQ.TEMP) THEN
  READY=.TRUE.
ENDIF
GOTO 10
ENDIF
READY=.FALSE.
20 IF (.NOT. READY) THEN

```

C THE SAME TEST IS THEN MADE UP TO THE FEEDER

```

IF (POI(1).EQ.SWITYP) THEN
  I=SWITCH(POI(2),2)
  I=SWITCH(POI(2),I)
ELSEIF (POI(1).EQ.NODTYP) THEN
  I=NODETO(POI(2),1)
  I=NODETO(POI(2),I)
  ALPHA=ALFA(POI(2))
ELSEIF (POI(1).EQ.BRATYP) THEN
  I=BRAPOI(POI(2),2)
  I=BRAPOI(POI(2),I)
ENDIF
IF (CABPHY(I,4).GT.CABPHY(I,3)) THEN
  IF ((CABPHY(I,4)-CABPHY(I,3)).GT.CURSUM) THEN
    CURSUM=MIN(CUR1,(CABPHY(I,4)-CABPHY(I,3)))
  ENDIF
ENDIF
CALL STEPUP(POI,.FALSE.)
IF (POI(1).EQ.FEETYP) THEN
  READY=.TRUE.
ENDIF

```

```
GOTO 20
ENDIF
```

```
C COMPUTE AND DECIDE
```

```
SUM=SQRT((CURSUM**2)+(VOLSUM**2))
SLASK=0.01+VIOSUM(FEE)-SUM
SLASK=LOAD*(SLASK+BETA)*ALPHA
IF (SLASK.LT.BEST .AND. SLASK.GT.0.0) THEN
    BEST=SLASK
    SWI=TEMP
ENDIF
```

```
C FINAL REMARK
C BEST AND SWI IS SAVED
```

```
RETURN
END
```

```
SUBROUTINE STEP3
```

```
CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC
```

```
C AUTHOUR MAGNUS BERGSTRAND
```

```
C VERSION 1
```

```
C DATE 1989-10-05
```

```
C PURPOSE
```

```
C WHEN NO MORE LOAD TRANSFER IS POSSIBLE LOAD MUST BE CUT OFF.
C AFTER A LOAD CURTAILMENT THERE IS A POSSIBILITY TO RESTORE
C LOAD AMONG FORMER VIOLATION FEEDERS.
```

```
C METHODE
```

```
C THE METHODE IS DESCRIBED IN 'SERVICE RESTORATION IN ELECTRIC
C DISTRIBUTION NETWORKS' BY MAGNUS BERGSTRAND.
C THE METHOD WAS FIRST DESCRIBED BY AOKI IN
C 'A NEW ALGORITHM FOR SERVICE RESRORATION IN DISTRIBUTION SYSTEMS'
C IEEE/PES WINTER MEETING,NEW YORK 1989 PAPER 89 WM 085-2 PWRD
```

```
C DESCRIPTION
```

```
C THIS SUBROUTINE HANDLES BOTH STEP3 AND STEP4 IN TME ALGORITHM
C (LOAD CURTAILMENT AND RESTORATION OF CURTAILED LOADS)
C THE SUBROUTINE WILL SEARCH FOR A VIOLATED FEEDER.
C WHEN A VIOLATED FEEDER IS FOUND A CALL OF IMPROV IS MADE, THUS
C PRODUCING THE NUMBER OF THE BEST SWITCH FOR A LOAD CURTAILMENT.
C SWITCHING OPERATION IS PERFORMED AND THE NEW STATE WILL BE COMPUTED
C THE NUMBER OF THE SWITCH IS SAVED ON THE FILE 'CURTAI.DAT' FOR LATER
C USE
C A DESCRIPTION OF THE FOURTH STEP IN THE ALGORITHM IS FOUND FURTHER DOWN
```

```
C SEE ALSO
```

FILES	ROUTINES	COMMENTS
CONSTANTS		ALL CONSTANTS DECLARED
COMMON		ALL COMMON VARIABLES DECLARED
MOVING	STEPUP	TAKES ONE STEP UP
MOVING	STEPDN	TAKES ONE STEP DOWN
MOVING	MOVE	TAKES ONE STEP UP OR DOWN
MOVING	DIREC	COMPUTES NEW DIRECTION
STEP2	IMPROV	COMPUTES THE NUMBER OF THE SWITCH THAT WILL BE OPENED DUE TO A LOAD CURTAILMENT
COMP	COMPIM	COMPUTES IMPEDANCE
COMP	FILVOL	COMPUTES AND FILLS IN NODE VOLTAGES AND

```

C          CABLE CURRENTS
C  COMP          FILPOW          COMPUTES POWER
C  STEP2        SAVENE          SAVES THE STATE OF THE NETWORK
C  STEP2        RESNET          READS BACK FORMER STATE
C  STEP2        STAGE1          FIRST STAGE SUPPORT

C  CCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCCC

INCLUDE 'CONSTANTS.FOR'
INCLUDE 'COMMON.FOR'

INTEGER I,FEE,OUTFEE,SWI,SWI2
LOGICAL VIOLTE
LOGICAL OK,FIRST,READY
LOGICAL VIO
REAL S1,S2,BSUM,SUM,MINI,SLASK,VIOSUM,LOAD
INTEGER BSWI,BSWI2,BFEE,FEEDER,TRA
INTEGER POI(2),P(2),OLD(2),DIR,BRA

EXTERNAL VIOSUM
EXTERNAL VIOLTE

WRITE(*,*)'ENTERS STEP 3'
OPEN(2,FILE='CURTAI',STATUS='NEW',ACCESS='SEQUENTIAL')
DO 10 I=1,MAXNRF
  FEE=I
  OK=.TRUE.
  FIRST=.TRUE.
  READY=.NOT. VIOLTE(FEE)
15  IF (.NOT. READY) THEN

C      A VIOLATED FEEDER IS FOUND

  WRITE(*,*)'STAGE 3 WILL BE CUTTING LOAD IN FEEDER',FEE
30  CALL IMPROV(FEE,SWI)

C      PERFORM SWITCHING OPERATIONS

  SWITCH(SWI,1)=OPEN
  SWINR(SWI)=2
  CHGSWI(SWI)=.NOT. CHGSWI(SWI)

C      SAVE NUMBER OF THE SWITCH ON 'CURTAI.DAT'

  WRITE(2,*) SWI

C      COMPUTE NEW STATE

  CALL COMPIM(FEE)
  TRA=MSFEED(FEE,1)
  CALL FILPOW(TRA)
  CALL FILVOL(FEE,VIO)
  READY=.NOT. VIO
  GOTO 15
  ENDIF
10  CONTINUE
ENDFILE 2
REWIND 2

C  DESCRIPTION
C  HERE STEP 4, THE RESTORATION OF CURTAILED LOADS, BEGINS.
C  ONE AT A TIME THE CURTAILED LOADS ARE REENRGIZED IF IT IS
C  POSSIBLE. ONE CURTAILED SECTION MAY LACK ENERGIZED
C  NEIGHBOURS DUE TO REPEATED LOAD CURTAILMENTS IN THE SAME
C  BRANCH.
C  THEN A FIRST STAGE SUPPORT IS TRIED. THE LOAD

```

```

C      WILL ONCE AGAIN BE CUT OFF IF THIS FAILS

      S1=0.0
      S2=0.0
40     READ(2,*,ERR=60)SWI

C      READ IN 'CURTAL.DAT'

      POI(1)=SWITYP
      POI(2)=SWI

C      THEN A SEARCH FOR THE FIRST SWITCH ABOVE THE READ SWITCH STARTS
C      IF THIS NEW SWITCH IS CLOSED IT IS POSSIBLE TO REENERGIZE THE
C      FORMER DEENERGIZED ZONE

      CALL STEPUP(POI,.FALSE.)
      READY=POI(1).EQ.SWITYP
50     IF (.NOT. READY) THEN
          CALL STEPUP(POI,.FALSE.)
          READY=POI(1).EQ.SWITYP
          GOTO 50
      ENDIF
      IF (SWITCH(POI(2),1).EQ.CLOSED) THEN
          READY=.FALSE.
          DIR=DOWN
          P(1)=POI(1)
          P(2)=POI(2)
65     IF (.NOT. READY) THEN
          CALL MOVE(P,OLD,DIR,BRA,.FALSE.)
          IF (P(1).EQ.NODTYP .AND. DIR.EQ.UP) THEN
              FEEDER=NODETO(P(2),4)
              S1=S1+NODEPH(P(2),1)
              S2=S2+NODEPH(P(2),2)
          ENDIF
          CALL DIREC(P,OLD,DIR,BRA)
          READY=(P(1).EQ.POI(1) .AND. P(2).EQ.POI(2))
          GOTO 65
      ENDIF

C      SAVE OLD STATE

      CALL SAVENE(FEEDER)

C      PERFORM SWITCHING OPERATION

      TRA=MSFEED(FEEDER,1)
      SWITCH(SWI,1)=CLOSED
      CHGSWI(SWI)=.NOT. CHGSWI(SWI)

C      COMPUTE NEW STATE

      CALL COMPIM(FEEDER)
      CALL FILPOW(TRA)
      CALL FILVOL(FEEDER,V10)
      OK=.TRUE.
      FIRST=.TRUE.

C      TRY A FIRST STAGE SUPPORT

      CALL STAGE1(FEEDER,OK,FIRST,OUTFEE)
      IF (.NOT. OK) THEN

C      READ BACK OLD STATE

      CALL RESNET
      RESTOR(1)=RESTOR(1)-S1

```



```
RESTOR(2)=RESTOR(2)-S2
ENDIF
ELSE
```

```
C THE SWITCH WAS CLOSED
C FIND OUT HOW BIG THE LOAD WAS THAT COULD NOT BE RESTORED
```

```
READY=.FALSE.
DIR=DOWN
P(1)=POI(1)
P(2)=POI(2)
58 IF (.NOT. READY) THEN
    CALL MOVE(P,OLD,DIR,BRA,.FALSE.)
    IF (P(1).EQ.NODTYP .AND. DIR.EQ.UP) THEN
        FEEDER=NODETO(P(2),4)
        S1=S1+NODEPH(P(2),1)
        S2=S2+NODEPH(P(2),2)
    ENDIF
    CALL DIREC(P,OLD,DIR,BRA)
    READY=(P(1).EQ.POI(1) .AND. P(2).EQ.POI(2))
    GOTO 58
```

```
ENDIF
RESTOR(1)=RESTOR(1)-S1
RESTOR(2)=RESTOR(2)-S2
```

```
ENDIF
GOTO 40
```

```
60 CLOSE(2,STATUS='DELETE')
RETURN
END
```

REAL FUNCTION VIOSUM(FEE)

C CCC

C AUTHOUR MAGNUS BERGSTRAND

C VERSION 1

C DATE 1989-10-05

C PURPOSE

C THIS FUNCTION COMPUTES THE EFFECTIVE LENGTH OF THE VIOLATION

C FILES

C STEP2

ROUTINES

VIOLTE

COMMENTS

INDICATES WHETHER A FEEDER IS VIOLATED OR NOT

C VARIABLES

C 'FEE' IS THE NUMBER OF THE (PERHAPS) VIOLATED FEEDER

C CCC

INCLUDE 'CONSTANTS.FOR'

INCLUDE 'COMMON.FOR'

INTEGER FEE

REAL SQRT

INTRINSIC SQRT

VIOSUM=SQRT(VIOVEC(FEE,1)**2+VIOVEC(FEE,2)**2)

RETURN

END