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## Plenary Lecture at Chinese Control Conference 2011: Distributed Control Using Positive Quadratic Programming

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# Distributed Control Using Positive Quadratic Programming

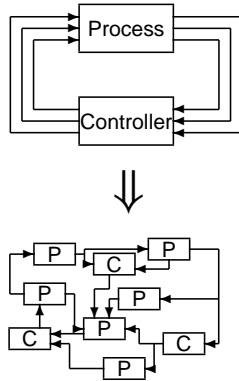
Anders Rantzer

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LCCC — Lund Center for Control of Complex Engineering Systems  
Lund University

## Outline

- Why Distributed Control?
  - Distributed Control of Positive Systems
  - Example: Optimizing Electrical Power Flow
  - Solution using Positive Quadratic Programming
  - Finding Optimum by Distributed Control

## Building theoretical foundations for distributed control



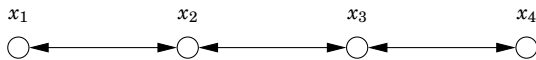
We need methodology for

- ▶ Decentralized specifications
- ▶ Decentralized design
- ▶ Verification of global behavior

## Example 1: A vehicle formation



## Example 1: A vehicle formation



Each vehicle obeys the independent dynamics

$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \\ \dot{x}_3(t) \\ \dot{x}_4(t) \end{bmatrix} = \begin{bmatrix} * & 0 & 0 & 0 \\ 0 & * & 0 & 0 \\ 0 & 0 & * & 0 \\ 0 & 0 & 0 & * \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ x_4(t) \end{bmatrix} + \begin{bmatrix} B_1 u_1(t) + w_1(t) \\ B_2 u_2(t) + w_2(t) \\ B_3 u_3(t) + w_3(t) \\ B_4 u_4(t) + w_4(t) \end{bmatrix}$$

The objective is to make  $\mathbf{E}|Cx_{i+1} - Cx_i|^2$  small for  $i = 1, \dots, 4$ .

## Example 3: A Wind Farm



## Example 2: A supply chain for fresh products



Fresh products degrade with time:

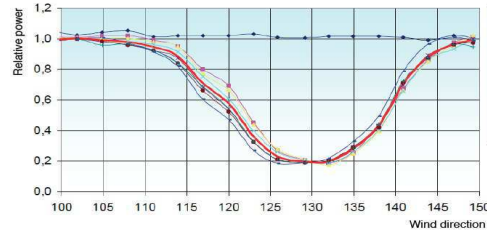
$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \\ \dot{x}_3(t) \\ \dot{x}_4(t) \end{bmatrix} = \begin{bmatrix} * & 0 & 0 & 0 \\ 0 & * & 0 & 0 \\ 0 & 0 & * & 0 \\ 0 & 0 & 0 & * \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ x_4(t) \end{bmatrix} + \begin{bmatrix} -u_1(t) + w_1(t) \\ u_1(t) - u_2(t) \\ u_2(t) - u_3(t) \\ u_3(t) + w_4(t) \end{bmatrix}$$

## Example 3: A Wind Farm



$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \\ \dot{x}_3(t) \\ \dot{x}_4(t) \end{bmatrix} = \begin{bmatrix} * & 0 & 0 & 0 \\ * & * & 0 & 0 \\ 0 & * & * & 0 \\ 0 & 0 & * & * \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ x_4(t) \end{bmatrix} + \begin{bmatrix} B_1 u_1 + w_1 \\ B_2 u_2 + w_2 \\ B_3 u_3 + w_3 \\ B_4 u_4 + w_4 \end{bmatrix}$$

### Example 3: A Wind Farm

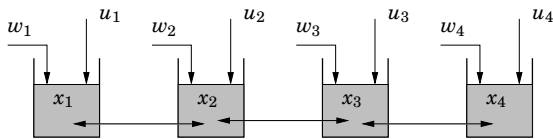


### Outline

- o Why Distributed Control?
- **Distributed Control of Positive Systems**
- o Example: Optimizing Electrical Power Flow
- o Solution using Positive Quadratic Programming
- o Finding Optimum by Distributed Control



### Example 4: Irrigation Channels



$$\begin{bmatrix} \dot{x}_1(t) \\ \dot{x}_2(t) \\ \dot{x}_3(t) \\ \dot{x}_4(t) \end{bmatrix} = \begin{bmatrix} * & * & 0 & 0 \\ * & * & * & 0 \\ 0 & * & * & * \\ 0 & 0 & * & * \end{bmatrix} \begin{bmatrix} x_1(t) \\ x_2(t) \\ x_3(t) \\ x_4(t) \end{bmatrix} + \begin{bmatrix} u_1(t) + w_1(t) \\ u_2(t) + w_2(t) \\ u_3(t) + w_3(t) \\ u_4(t) + w_4(t) \end{bmatrix}$$

**Note:** Off-diagonal elements are typically positive!

### Positive systems have nonnegative impulse response

If the matrices  $A$ ,  $B$  and  $C$  have nonnegative coefficients except possibly for the diagonal of  $A$ , then the system

$$\frac{dx}{dt} = Ax + Bu$$

$$y = Cx$$

has non-negative impulse response  $Ce^{At}B$ .

#### Examples:

- ▶ Ecological system with  $x_k$  the population of species  $k$ .
- ▶ Chemical reaction with  $x_k$  the concentration of reactant  $k$ .
- ▶ Economic system with  $x_k$  the quantity of commodity  $k$ .
- ▶ Probabilistic model with  $x_k$  the probability of state  $k$ .

### Stability of Positive systems

Suppose the matrix  $A$  has nonnegative off-diagonal elements. Then the following conditions are equivalent:

- (i) The system  $\frac{dx}{dt} = Ax$  is exponentially stable.
- (ii) There exists a vector  $x > 0$  such that  $Ax < 0$ . (The vector inequalities are elementwise.)
- (iii) There is a diagonal matrix  $P > 0$  such that  $PA^T + AP < 0$

### Positive Systems and Nonnegative Matrices

#### Classics:

- ▶ Perron (1907) and Frobenius (1912)
- ▶ Leontief (1936)
- ▶ Hirsch (1985)

#### Books:

- ▶ Gantmacher (1959)
- ▶ Berman and Plemmons (1979)
- ▶ Luenberger (1979)

#### Recent control related work:

- ▶ Angeli and Sontag (2003)
- ▶ Moreau (2004)

### Stability can be Tested in a Distributed Way



Stability of  $\dot{x} = Ax$  follows from existence of  $x_k > 0$  such that

$$\underbrace{\begin{bmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{32} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix}}_A \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} < \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

The first node verifies the inequality of the first row.

The second node verifies the inequality of the second row.

and so on ...

## Performance of Positive systems

Suppose the matrices  $A$ ,  $B$  and  $C$  have nonnegative coefficients except for the diagonal of  $A$ . Suppose  $A$  is Hurwitz. Then the following conditions are equivalent:

- (i)  $\max_{\omega} |C(i\omega I - A)^{-1}B| < \gamma$
- (ii)  $|CA^{-1}B| < \gamma$
- (iii) There exists  $x > 0$  such that  $Cx < \gamma$ ,  $Ax + B = 0$ .
- (iv) There is a *diagonal* matrix  $P > 0$  such that

$$PA^T + AP + PC^T CP + \gamma^{-2}BB^T < 0$$

**Note:** The linear inequalities (iii) can be tested row by row.

## Synthesizing Positive Systems

$$A + BL = \begin{bmatrix} a_{11} + \ell_1 & a_{12} & 0 & 0 \\ a_{21} - \ell_1 & a_{22} + \ell_2 & a_{23} & 0 \\ 0 & a_{32} - \ell_2 & a_{33} & a_{32} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix}$$

is stable and nonnegative if and only if  $p_k \geq 0$  and

$$(A + BL)P = \begin{bmatrix} (a_{11} + \ell_1)p_1 & a_{12}p_2 & 0 & 0 \\ (a_{21} - \ell_1)p_1 & (a_{22} + \ell_2)p_2 & a_{23}p_3 & 0 \\ 0 & (a_{32} - \ell_2)p_2 & a_{33}p_3 & a_{32}p_4 \\ 0 & 0 & a_{43}p_3 & a_{44}p_4 \end{bmatrix}$$

make  $(A + BL)P + P(A + BL)^T$  negative definite with nonnegative off-diagonal elements.

Solve using convex optimization in the pair  $(P, PL)$ !

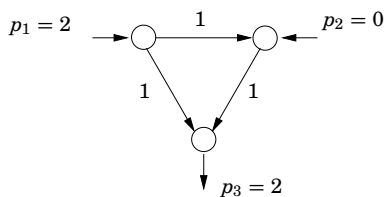
[Tanaka and Langbort, ACC 2010]

## Outline

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## Optimal Allocation for Example A

Both transmission lines serving the load need to be used at full capacity to meet the demand  $p_3 = 2$ .



Optimal profit:  $10p_3 - p_1 = 18$

In real power networks, electrons flow according to Kirchhoff's laws. The allocation above is not feasible when all three lines are identical. Why?

## Distributed Control Synthesis

Suppose the matrix

$$\begin{bmatrix} a_{11} + \ell_1 & a_{12} & 0 & 0 \\ a_{21} - \ell_1 & a_{22} + \ell_2 & a_{23} & 0 \\ 0 & a_{32} - \ell_2 & a_{33} & a_{32} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix}$$

is nonnegative for all  $\ell_1, \ell_2 \in [0, 1]$ . For stabilizing gains  $\ell_1, \ell_2$ , find  $0 \leq u_k \leq x_k$  such that

$$\begin{bmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{32} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} + \begin{bmatrix} 1 & 0 \\ -1 & 1 \\ 0 & -1 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} u_1 \\ u_2 \end{bmatrix} < \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$

and set  $\ell_1 = u_1/x_1$  and  $\ell_2 = u_2/x_2$ . Every row gives a local test.

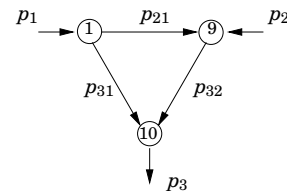
**Note:** Positivity assumed a priori. What if  $\ell_1, \ell_2 \in \mathbf{R}$ ?

## Positivity versus Passivity

- Passivity can be described naturally in frequency domain.
- Positivity can be described naturally in time-domain.
- Negative feedback loops preserve passivity.
- Positive feedback loops preserve positivity.
- Parallel connections preserve both passivity and positivity.
- Series connections preserves positivity, but not passivity.

## Example A: Electrical Power Transmission

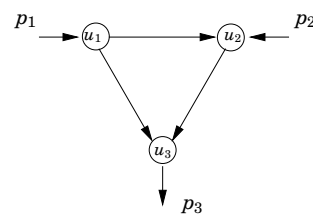
Two generators with generation cost 1 and 9 respectively. One load willing to buy  $p_3 = 2$  at the price 10:



Maximize profit:  $10p_3 - 9p_2 - p_1$   
 subject to capacity constraints:  $|p_{jk}| \leq 1, p_1 \geq 0, p_2 \geq 0, p_3 \geq 2$   
 and conservation laws:  
 $p_1 = p_{21} + p_{31}$   
 $p_{32} = p_{21} + p_2$   
 $p_3 = p_{31} + p_{32}$

## Example B: Optimal Potential Flow

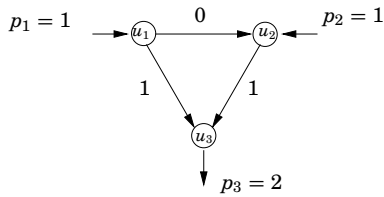
Power flow is driven by potential differences:



Maximize profit:  $10p_3 - 9p_2 - p_1$   
 subject to capacity constraints:  $|u_j - u_k| \leq 1, p_j \geq 0, p_3 \geq 2$   
 and conservation laws:  
 $p_1 = (u_1 - u_2) + (u_1 - u_3)$   
 $p_2 = (u_2 - u_1) + (u_2 - u_3)$   
 $p_3 = (u_1 - u_3) + (u_2 - u_3)$

## Optimal Allocation for Example B

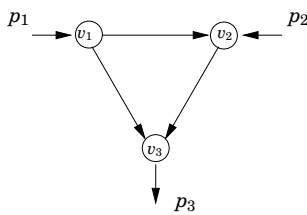
Both transmission lines serving the load need to be used at full capacity to meet the demand  $p_3 = 2$ . Hence  $u_1 = u_2$  and there is no flow between node 1 and node 2!



The optimal profit is much smaller:  $10p_3 - p_1 - 9p_2 = 10$

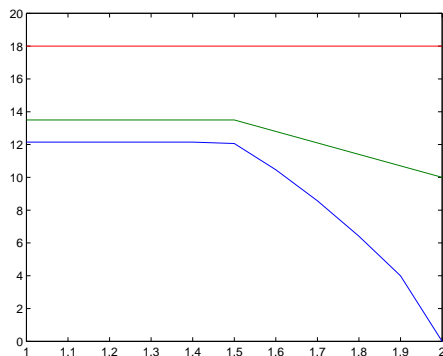
When transmission lines operate near capacity limits, losses are big. Can we take losses into account in the optimization?

## Example C: Optimal Power Flow with Losses



Maximize profit:  $10p_3 - 9p_2 - p_1$   
 subject to capacity constraints:  $0 \leq v_j \leq 2$   
 and conservation laws:  $p_1 = v_1(v_1 - v_2) + v_1(v_1 - v_3)$   
 $p_2 = v_2(v_2 - v_1) + v_2(v_2 - v_3)$   
 $p_3 = v_3(v_1 - v_3) + v_3(v_2 - v_3)$

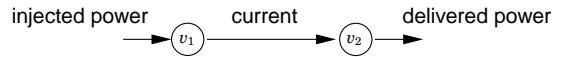
## Profit Versus Power Demand



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## Power Losses in a DC Transmission Line



For a DC transmission line with admittance  $y$ , input voltage  $v_1$  and output voltage  $v_2$ , we have:

Line current:  $i = y(v_1 - v_2)$   
 Injected power:  $p_1 = yv_1(v_1 - v_2)$   
 Delivered power:  $p_2 = yv_2(v_1 - v_2)$   
 Power loss:  $p_1 - p_2 = y(v_1 - v_2)^2$

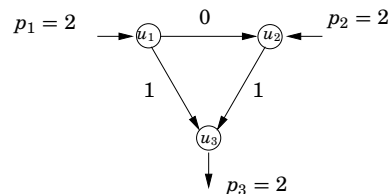
If the voltages are bounded from above by  $\bar{v}$ , there is an upper bound on how much power the transmission line can deliver:

$$p_2 = yv_2(v_1 - v_2) \leq yv_2(\bar{v} - v_2) \leq y\bar{v}^2/4$$

At the capacity limit, the power loss equals the delivered power.

## Optimal Allocation for Example C

Both transmission lines serving the load need to be used at full capacity to meet the demand  $p_3 = 2$ . Hence  $v_1 = v_2 = \bar{v}$  and there is no current between node 1 and node 2!



There is no room for profit:  $10p_3 - p_1 - 9p_2 = 0$

Notice that half of the generated power is lost in transmission!

## Analogies to Electric Power Flow

**Water distribution systems:** Electrical voltage corresponds to water pressure. Differences in pressure creates flow.

**Gas diffusion:** Electrical voltage corresponds to partial pressure. Gradients in partial pressure creates diffusion.

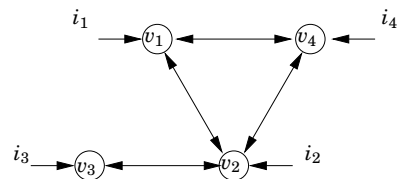
**Exchange economy:** Voltages correspond to *inverse* prices. Price differences drive commodity flows. Delivered electric power corresponds to delivered commodity volume.

**Two kinds of flow of simultaneous interest.**

In power transmission networks, electric current is conserved, but electric power is dissipated due to transmission losses.

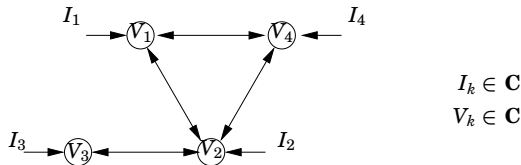
In economic systems the commodity value is conserved, but the commodity volume is dissipated due to transportation losses.

## A General Power Transmission Network



$$\begin{bmatrix} i_1 \\ i_2 \\ i_3 \\ i_4 \end{bmatrix} = \begin{bmatrix} Y_{12} + Y_{14} & -Y_{12} & 0 & -Y_{14} \\ -Y_{21} & Y_{21} + Y_{23} + Y_{24} & -Y_{23} & -Y_{24} \\ 0 & -Y_{32} & Y_{32} & 0 \\ -Y_{41} & -Y_{42} & 0 & Y_{41} + Y_{42} \end{bmatrix} \begin{bmatrix} v_1 \\ v_2 \\ v_3 \\ v_4 \end{bmatrix}$$

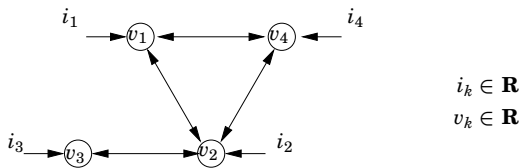
## An Optimal Flow Problem for AC Power



Minimize  $\text{Re} \sum_k I_k^* V_k$   
 subject to  $I = YV$  and  $\underline{P}_k \leq \text{Re}(I_k^* V_k) \leq \bar{P}_k$   
 $\underline{Q}_k \leq \text{Im}(I_k^* V_k) \leq \bar{Q}_k$   
 $\underline{v}_k \leq |V_k| \leq \bar{v}_k$  for  $k = 1, \dots, 4$

(Convex relaxation by Lavaei/Low inspired this talk.)

## Optimizing DC Power Flow



Minimize  $\sum_k i_k v_k$   
 subject to  $i = Yv$  and  $i_k v_k \leq \bar{p}_k$   
 $(v_k - v_j)^2 \leq c_{kj}$   
 $\underline{v}_k \leq v_k \leq \bar{v}_k$  for all  $k, j$

Notice:  $\bar{p}_k$  negative at loads, positive at generators.

## Positive Quadratic Programming

Given  $A_0, \dots, A_K \in \mathbf{R}^{n \times n}$  with nonnegative off-diagonal entries and  $b_1, \dots, b_K \in \mathbf{R}$ , the following equality holds:

max  $x^T A_0 x$  = max  $\text{trace}(A_0 X)$   
 subject to  $x \in \mathbf{R}_+^n$  subject to  $X \geq 0$   
 $x^T A_k x \geq b_k$   $\text{trace}(A_k X) \geq b_k$   
 $k = 1, \dots, K$   $k = 1, \dots, K$

**Proof**

If  $X = \begin{bmatrix} |x_1|^2 & & * \\ & \ddots & \\ * & & |x_n|^2 \end{bmatrix}$  maximizes the right hand side,  
 then  $x = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$  maximizes the left.

**Note:** The problem is convex in  $|v_1|^2, \dots, |v_4|^2$ !

## Dual Positive Quadratic Programming

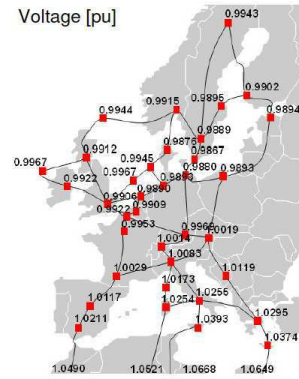
Given  $A_0, \dots, A_K \in \mathbf{R}^{n \times n}$  with nonnegative off-diagonal entries and  $b_1, \dots, b_K \in \mathbf{R}$ , the following equality holds:

max  $x^T A_0 x$  = min  $-\sum_k \lambda_k b_k$   
 subject to  $x \in \mathbf{R}_+^n$  subject to  $\lambda_1, \dots, \lambda_K \geq 0$   
 $x^T A_k x \geq b_k$   $0 \geq A_0 + \sum_k \lambda_k A_k$   
 $k = 1, \dots, K$

**Interpretation:**

In the power flow example,  $\lambda_k$  is the price of power at node  $k$ .

## Future DC Power Transmission Network in Europe?



From Cigré Conference 2010, "Continental Overlay HVDC-Grid" by ABB

## Positive Quadratic Programming

Given  $A_0, \dots, A_K \in \mathbf{R}^{n \times n}$  with nonnegative off-diagonal entries and  $b_1, \dots, b_K \in \mathbf{R}$ , the following equality holds:

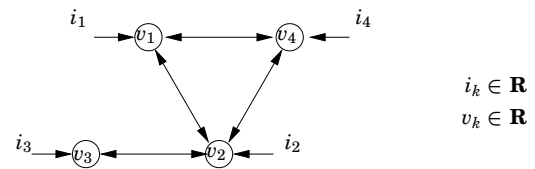
max  $x^T A_0 x$  = max  $\text{trace}(A_0 X)$   
 subject to  $x \in \mathbf{R}_+^n$  subject to  $X \geq 0$   
 $x^T A_k x \geq b_k$   $\text{trace}(A_k X) \geq b_k$   
 $k = 1, \dots, K$   $k = 1, \dots, K$

**Proof**

If  $X = \begin{bmatrix} |x_1|^2 & & * \\ & \ddots & \\ * & & |x_n|^2 \end{bmatrix}$  maximizes the right hand side,  
 then  $x = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$  maximizes the left.

[Goemans/Williamson (1994), Zhang (1999), Kim/Kojima (2003)]

## Optimizing DC Power Flow



Minimize  $\sum_k i_k v_k$   
 subject to  $i = Yv$  and  $i_k v_k \leq \bar{p}_k$   
 $(v_k - v_j)^2 \leq c_{kj}$   
 $\underline{v}_k \leq v_k \leq \bar{v}_k$  for all  $k, j$

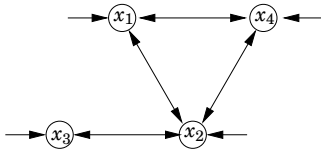
Notice: All mixed terms have the right sign!

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## Primal Decomposition



The convex problem

$$\min_{x_k} [V_1(x_1, x_2, x_4) + V_2(x_1, x_2, x_3, x_4) + V_3(x_2, x_3) + V_4(x_1, x_2, x_4)]$$

can be solved by the following distributed iteration:

$$\begin{cases} x_1^+ = \arg \min_{x_1} [V_1(x_1, x_2, x_4) + V_2(x_1, x_2, x_3, x_4) + V_4(x_1, x_2, x_4)] \\ x_2^+ = \arg \min_{x_2} [V_1(x_1, x_2, x_4) + V_2(x_1, x_2, x_3, x_4) + V_3(x_2, x_3) + V_4(x_1, x_2, x_4)] \\ x_3^+ = \arg \min_{x_3} [V_2(x_1, x_2, x_3, x_4) + V_3(x_2, x_3)] \\ x_4^+ = \arg \min_{x_4} [V_1(x_1, x_2, x_4) + V_2(x_1, x_2, x_3, x_4) + V_4(x_1, x_2, x_4)] \end{cases}$$

## The Distributed Control Law

The dynamics

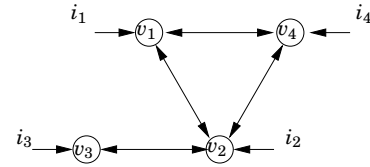
$$v_k^+ = \arg \min_{v_k \leq \bar{v}_k} \sum_j \left[ \underbrace{\lambda_k y_{kj} v_k (v_k - v_j)}_{\text{value into link } jk} - \underbrace{\lambda_j y_{jk} v_j (v_k - v_j)}_{\text{value out from link } jk} \right]$$

has the form

$$v^+ = \min\{\bar{v}, Av\}$$

where  $A$  has nonnegative coefficients.

## Finding Optimum by Distributed Control



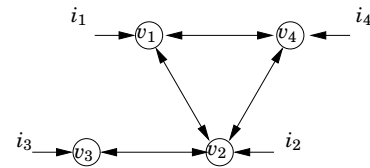
Given power prices  $\lambda_k$  at each node, find the optimal allocation:

$$\text{Minimize } \sum_{j,k} \lambda_k y_{kj} v_k (v_k - v_j) \text{ subject to } v_k \leq \bar{v}_k$$

Primal decomposition gives convergence to optimum:

$$v_k^+ = \arg \min_{v_k \leq \bar{v}_k} \sum_j \left[ \underbrace{\lambda_k y_{kj} v_k (v_k - v_j)}_{\text{value into link } jk} - \underbrace{\lambda_j y_{jk} v_j (v_k - v_j)}_{\text{value out from link } jk} \right]$$

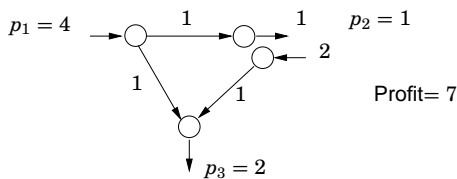
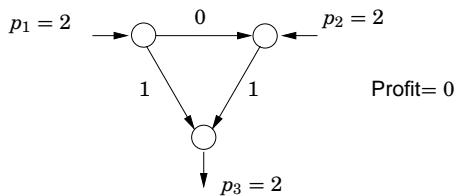
## Can it pay off to disconnect a line?



Given power prices  $\lambda_k$  at each node, find the optimal allocation:

$$\text{Minimize } \sum_{j,k} \lambda_k y_{kj} v_k (v_k - v_j) \text{ subject to } v_k \leq \bar{v}_k, y_{kj} \in [0, \bar{y}_{kj}]$$

## Yes, it can!



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## Summary

- ▶ Why Distributed Control?
- ▶ Optimizing Electrical Power Flow
- ▶ Positive Quadratic Programming
- ▶ Distributed Control of Positive Systems
- ▶ Finding Optimum by Distributed Control

To read:

Slides on [www.control.lth.se/Staff/anders\\_rantzer.html](http://www.control.lth.se/Staff/anders_rantzer.html)  
 Extended abstract in Proceedings of CCC 2011  
 Upcoming paper in CDC 2011